

## CPS 130: Homework 7

### More Graphs: Spanning Trees, Shortest Paths (CLRS: Chapter 23.2, 24.3, 25.2)

Date on which distributed: Thursday, November 01, 2001

Date on which due: Tuesday, November 13, 2001

Note: Zero credit will be given for homeworks submitted late.

**1. [Exc 23.2-2 CLRS]**

Suppose that the graph  $G = (V, E)$  is represented as an adjacency matrix. Give a simple implementation of Prim's algorithm for this case that runs in  $O(V^2)$  time.

**2. [Exc 23.2-4 CLRS]**

We know that Kruskal's algorithm runs in  $O(E \lg V)$  time on a graph  $G = (V, E)$ . Suppose that all edge weights in  $G$  are integers in the range from 1 to  $|V|$ . How fast can you make Kruskal's algorithm on  $G$ ? What if the edge weights are integers in the range from 1 to  $W$  for some constant  $W$ ? (You can assume  $W$  is an integer too.)

**3. [Exc 24.3-4 CLRS]**

We are given a directed graph  $G = (V, E)$  on which each edge  $(u, v) \in E$  has an associated value  $r(u, v)$  which is a real number in the range  $0 \leq r(u, v) \leq 1$  that represents the reliability of a communication channel from vertex  $u$  to vertex  $v$ . We interpret  $r(u, v)$  as the probability that the channel from  $u$  to  $v$  will not fail, and we assume that these probabilities are independent. Give an efficient algorithm to find the most reliable path between two given vertices. (Hint: The probability of the *success* of any path from  $u$  to  $v$  is the *product* of the probabilities of success of the individual edges that make up this path. You may want to try to convert the problem to a Dijkstra-like situation. If this is your approach, remember that Dijkstra's algorithm finds the *shortest* path between any two given vertices, and that weights of paths are computed as *sums*, not *products* of weights of edges along them.)

**4. [Exc 25.2-6 CLRS]**

How can the output of the Floyd-Warshall algorithm be used to detect the presence of a negative-weight cycle?