

Homework 2 (due before class, Monday, September 24, 2001)

## 1 Recurrences, Master Method, Sorting [Chapter 4 in CLRS]

1. [CLRS 4.2-5] Solve the recurrence  $T(n) = T(\alpha n) + T((1 - \alpha)n) + n$ , where  $0 < \alpha < 1$  constant. (Hint: Use a recursion tree. Use the `parsetree.sty` package to L<sup>A</sup>T<sub>E</sub>X the tree.)
2. Give asymptotic upper and lower bounds using  $\Theta$  notation for  $T(n)$  at those points where  $T(n)$  is defined, for the following recurrences. Assume  $T(n)$  is constant for  $n \leq 20$ . Make your bounds as tight as possible, and justify your answers.

(a)  $T(n) = 2T(n/3) + n^3$

(b)  $T(n) = 3T(n/4) + \sqrt{n}$

(c)  $T(n) = 9T(n/4) + n^2$

(d)  $T(n) = T(n - 6) + n$

(e)  $T(n) = T(\sqrt{n}) + 1$

(f)  $T(n) = 2T(n/2) + n \lg^2 n$

(g)  $T(n) = 3T(n/3 + 5) + n/2$

(h)  $T(n) = 2T(n/2) + n/\lg n$

(i)  $T(n) = T(n - 2) + 1/n$

(j)  $T(n) = \sqrt{n}T(\sqrt{n}) + n$

## 2 Quicksort, Median, Matrices [Chapter 7, 9, 28 in CLRS]

3. [CLRS 9.3-1] In the algorithm `Select`, the input elements are divided into groups of 5. Will the algorithm work in linear time if they are divided into groups of 7? How about groups of 3?

4. [CLRS 9-2] For  $n$  distinct elements  $x_1, x_2, \dots, x_n$  with positive weights  $w_1, w_2, \dots, w_n$  such that  $\sum_{i=0}^n w_i = 1$ , the weighted (lower) median is the element  $x_k$  satisfying

$$\sum_{x_i < x_k} w_i < \frac{1}{2} \text{ and } \sum_{x_i > x_k} w_i \leq \frac{1}{2}$$

- (a) Argue that the median of  $x_1, x_2, \dots, x_n$  is the weighted median of the  $x_i$  with weights  $w_i = 1/n$  for  $i = 1, 2, \dots, n$ .
- (b) Show how to compute the weighted median of  $n$  elements in  $O(n \lg n)$  worst-case time using sorting.
- (c) Show how to compute the weighted median in  $\Theta(n)$  worst-case time using a linear-time median algorithm.

The post office location problem is defined as follows. We are given  $n$  points  $p_1, p_2, \dots, p_n$  with associated weights  $w_1, w_2, \dots, w_n$ . We wish to find a point  $p$  (not necessarily one of the input points) that minimizes the sum  $\sum_{i=1}^n w_i d(p, p_i)$ , where  $d(a, b)$  is the distance between points  $a$  and  $b$ . Intuitively, the point  $p$  represents a location of the post office that minimizes average distance.

- (d) Argue that the weighted median is a best solution for the 1-dimensional post office location problem, in which points are simply real numbers and the distance between points  $a$  and  $b$  is  $d(a, b) = |a - b|$ .
- (e) Find the best solution for the 2-dimensional post office location problem, in which the points are  $(x, y)$  coordinate pairs and the distance between points  $a = (x_1, y_1)$  and  $b = (x_2, y_2)$  is the Manhattan distance given by  $d(a, b) = |x_1 - x_2| + |y_1 - y_2|$ .

**5.** Show that the smallest and the second smallest of  $n$  distinct elements can be found with  $n + \lceil \lg n \rceil - 2$  comparisons in the worst case.