## Solving Problems Recursively

- Recursion is an indispensable tool in a programmer's toolkit
> Allows many complex problems to be solved simply
$>$ Elegance and understanding in code often leads to better programs: easier to modify, extend, verify (and sometimes more efficient!!)
> Sometimes recursion isn't appropriate, when it's bad it can be very bad---every tool requires knowledge and experience in how to use it
- The basic idea is to get help solving a problem from coworkers (clones) who work and act like you do
> Ask clone to solve a simpler but similar problem
> Use clone's result to put together your answer
- Need both concepts: call on the clone and use the result


## Print words entered, but backwards

- Can use an ArrayList, store all the words and print in reverse order
> Probably the best approach, recursion works too

```
public void printReversed(Scanner s){
    if (s.hasNext()){ // reading succeeded?
        String word = s.next(); // store word
        printReversed(s); // print rest
        System.out.println(word);// print the word
    }
}
```

- The function printReversed reads a word, prints the word only after the clones finish printing in reverse order
$>$ Each clone has own version of the code, own word variable
> Who keeps track of the clones?
$>$ How many words are created when reading $N$ words?
- What about when ArrayList<String> used?


## Exponentiation

- Computing $x^{n}$ means multiplying $n$ numbers (or does it?)
$>$ What's the easiest value of $n$ to compute $x^{n}$ ?
> If you want to multiply only once, what can you ask a clone?

```
public static double power(double x, int n){
    if (n == 0){
        return 1.0;
    }
    return x * power(x, n-1);
}
```

- What about an iterative version?


## Faster exponentiation

- How many recursive calls are made to computer $2^{1024}$ ?
> How many multiplies on each call? Is this better?

```
public static double power(double x, int n) {
    if (n == 0) {
            return 1.0;
        }
        double semi = power(x, n/2);
        if (n % 2 == 0) {
            return semi*semi;
        }
        return x * semi * semi;
}
```

- What about an iterative version of this function?


## Back to Recursion

- Recursive functions have two key attributes
> There is a base case, sometimes called the exit case, which does not make a recursive call
- See print reversed, exponentiation
> All other cases make a recursive call, with some parameter or other measure that decreases or moves towards the base case
- Ensure that sequence of calls eventually reaches the base case
- "Measure" can be tricky, but usually it's straightforward
- Example: sequential search in an array
> If first element is search key, done and return
> Otherwise look in the "rest of the array"
> How can we recurse on "rest of array"?


## Thinking recursively

- Problem: find the largest element in an array
> Iteratively: loop, remember largest seen so far
$>$ Recursive: find largest in [1..n), then compare to $0^{\text {th }}$ element

```
public static double max(double[] a) {
    double maxSoFar = a[0];
    for(int k=1; k < a.length; k++) {
            maxSoFar = Math.max(maxSoFar,a[k]);
    }
    return maxSoFar;
}
```

> In a recursive version what is base case, what is measure of problem size that decreases (towards base case)?

## Recursive Max

```
public static double recMax(double[] a, int index){
    if (index == a.length-1){ // last element, done
        return a[index];
    }
    double maxAfter = recMax(a,index+1);
    return Math.max(a[index],maxAfter);
}
```

- What is base case (conceptually)?
> Do we need variable maxAfter?
- We can use recMax to implement arrayMax as follows
return recMax $(a, 0)$;


## Recognizing recursion:

```
public static void change(String[] a, int first, int last){
    if (first < last) {
        String temp = a[first]; // swap a[first], a[last]
        a[first] = a[last];
        a[last] = temp;
        change(a, first+1, last-1);
    }
}
// original call (why?): change(a, 0, a.length-1);
```

- What is base case? (no recursive calls)
- What happens before recursive call made?
- How is recursive call closer to the base case?


## More recursion recognition

```
public static int value(int[] a, int index){
    if (index < a.length) {
        return a[index] + value(a,index+1);
    }
    return 0;
}
// original call: int v = value(a,0);
```

- What is base case, what value is returned?
- How is progress towards base case realized?
- How is recursive value used to return a value?
- What if a is array of doubles, does anything change?


## Analysis: Algorithms and Data Structures

- We need a vocabulary to discuss performance and to reason about alternative algorithms and implementations
> It's faster! It's more elegant! It's safer! It's cooler!
- We need empirical tests and analytical/mathematical tools
> Given two methods, which is better? Run them to check.
- 30 seconds vs. 3 seconds, easy. 5 hours vs. 2 minutes, harder
- What if it takes two weeks to implement the methods?
> Use mathematics to analyze the algorithm,
> The implementation is another matter, cache, compiler optimizations, OS, memory,...


## Jaron Lanier (nttp: ///mm. advenced. ory/jaron)

Jaron Lanier is a computer scientist, composer, visual artist, and author. He coined the term 'Virtual Reality' ... he co-developed the first implementations of virtual reality applications in surgical simulation, vehicle interior prototyping, virtual sets for television production, and assorted other areas
"What's the difference between a bug and a variation or an imperfection? If you think about it, if you make a small change to a program, it can result in an enormous change in what the program does. If nature worked that way, the universe would crash all the time."

Lanier has no academic degrees


## Recursion and recurrences

- Why are some functions written recursively?
> Simpler to understand, but ...
> Mt. Everest reasons
- Are there reasons to prefer iteration?
> Better optimizer: programmer/scientist v. compiler
> Six of one? Or serious differences
- "One person's meat is another person's poison"
- "To each his own", "Chacun a son gout", ...
- Complexity (big-Oh) for iterative and recursive functions
> How to determine, estimate, intuit


## What's the complexity of this code?

```
// first and last are integer indexes, list is List
int lastIndex = first;
Comparable pivot = list.get(first);
for(int k=first+1; k <= last; k++) {
    Comparable ko = list.get(k);
    if (ko.compareTo(pivot) <= 0){
        lastIndex++;
        Collections.swap(list,lastIndex,k);
    }
}
```

- What is big-Oh cost of a loop that visits $n$ elements of a vector?
> Depends on loop body, if body $O$ (1) then ...
> If body is $O(n)$ then ...
> If body is $O(k)$ for $k$ in [0..n) then ...

```
FastFinder.findHelper
private Object findHelper(ArrayList<Comparable> list,
                                    int first, int last, int kindex) {
    int lastIndex = first;
    Comparable pivot = list.get(first);
    for(int k=first+1; k <= last; k++) {
    Comparable ko = list.get(k) ;
    if (ko.compareTo(pivot) <= 0) {
        lastIndex++;
        Collections.swap(list,lastIndex,k);
    }
    }
    Collections.swap(list,lastIndex,first);
    if (lastIndex == kindex) return list.get(lastIndex);
    if (kindex < lastIndex)
    return findHelper(list,first,lastIndex-1,kindex);
    return findHelper(list,lastIndex+1,last,kindex);
}
```


## Different measures of complexity

- Worst case
> Gives a good upper-bound on behavior
$>$ Never get worse than this
> Drawbacks?
- Average case
$>$ What does average mean?
$>$ Averaged over all inputs? Assuming uniformly distributed random data?
> Drawbacks?
- Best case
> Linear search, useful?


## Multiplying and adding big-Oh

- Suppose we do a linear search then we do another one
> What is the complexity?
$>$ If we do 100 linear searches?
$>$ If we do $\boldsymbol{n}$ searches on vector of size $\mathbf{n}$ ?
- What if we do binary search followed by linear search?
$>$ What are big-Oh complexities? Sum?
> What about 50 binary searches? What about $n$ searches?
- What is the number of elements in the list $(1,2,2,3,3,3)$ ?
$>$ What about ( $1,2,2, \ldots, n, n, \ldots, n$ )?
> How can we reason about this?


## Helpful formulae

- We always mean base 2 unless otherwise stated
$>$ What is $\log (1024)$ ?

```
\(>\log (x y) \quad \log \left(x^{y}\right) \quad \log \left(2^{n}\right) \quad 2^{(\log n)}\)
    - \(\log (x)+\log (y)\)
    \(\cdot y \log (x)\)
    -nlog(2) \(=n\)
    \(\cdot 2^{(\log n)}=n\)
```

$$
\begin{aligned}
& \bullet \text { Sums (also, use sigma notation when possible) } \\
& >1+2+4+8+\ldots+2^{k}=2^{k+1}-1=\sum_{2^{i}}^{k} \\
& >1+2+3+\ldots+n=n(n+1) / 2=\sum_{i=0}^{n} \\
& >a+a r+a r^{2}+\ldots+a r^{n-1}=a\left(r^{n}-1\right) /(r-1)=\sum_{i=0}^{n-1} a r^{i}
\end{aligned}
$$

## Recursion Review

- Recursive functions have two key attributes
> There is a base case, sometimes called the exit case, which does not make a recursive call
> All other cases make recursive call(s), the results of these calls are used to return a value when necessary
- Ensure that every sequence of calls reaches base case
- Some measure decreases/moves towards base case
- "Measure" can be tricky, but usually it's straightforward
- Example: sequential search in an ArrayList
> If first element is search key, done and return
> Otherwise look in the "rest of the list"
> How can we recurse on "rest of list"?


## Sequential search revisited

- What is complexity of sequential search? Of code below?

```
boolean search(ArrayList<Object> list,
    int first, Object target) {
    if (first >= list.size()) return false;
    else if (list.get(first).equals(target))
        return true;
    else return search(list,first+1,target);
}
- Why are there three parameters? Same name good idea?
boolean search(ArrayList list, Object target) {
    return search(list,0,target);
}
```


## Why we study recurrences/complexity?

- Tools to analyze algorithms
- Machine-independent measuring methods
- Familiarity with good data structures/algorithms
- What is CS person: programmer, scientist, engineer? scientists build to learn, engineers learn to build
- Mathematics is a notation that helps in thinking, discussion, programming


## Recurrences

- Summing Numbers

```
int sum(int n)
{
        if (0 == n) return 0;
        else return n + sum(n-1);
    }
```

- What is complexity? justification?
- $T(n)=$ time to compute sum for $n$

$$
\begin{aligned}
& T(n)=T(n-1)+1 \\
& T(0)=1
\end{aligned}
$$

- instead of 1 , use $\mathrm{O}(1)$ for constant time
> independent of $\mathbf{n}$, the measure of problem size


## Solving recurrence relations

- plug, simplify, reduce, guess, verify?

$$
\begin{aligned}
& T(n)=T(n-1)+1 \\
& T(0)=1 \\
& T(n-1)=T(n-1-1)+1 \\
& T(n)=[T(n-2)+1]+1=T(n-2)+2 \\
& T(n-2)=T(n-2-1)+1 \\
& T(n)=[(T(n-3)+1)+1]+1=T(n-3)+3 \\
& T(n)=T(n-k)+k \quad \text { find the pattern! }
\end{aligned}
$$

Now, let $k=n$, then $T(n)=T(0)+n=1+n$

- get to base case, solve the recurrence: $O(n)$


## Complexity Practice

- What is complexity of Build? (what does it do?)

```
ArrayList<Integer> build(int n)
{
    if (0 == n) return new ArrayList<Integer>(); // empty
    ArrayList<Integer> list = build(n-1);
    for(int k=0;k < n; k++) {
            list.add(n);
    }
    return list;
}
```

- Write an expression for $T(n)$ and for $T(0)$, solve.


## Recognizing Recurrences

- Solve once, re-use in new contexts
> T must be explicitly identified
> n must be some measure of size of input/parameter
- $T(n)$ is the time for quicksort to run on an $n$-element vector

- Remember the algorithm, re-derive complexity


## Eugene (Gene) Myers

- Lead computer scientist/software engineer at Celera Genomics (now at Berkeley,now at ...?)
- "What really astounds me is the architecture of life. The system is extremely complex. It's like it was designed." ... "There's a huge intelligence there."


