

## SQL: Part I

CPS 116  
Introduction to Database Systems

## Announcements (Thu. Sep. 10)

- ❖ Homework #1 due next Tuesday
  - Dongtao will run a help session next Monday 4-5pm, in LSRC D344
    - Bring your questions!
- ❖ Course project description available!
  - Choice of “standard” or “open”
  - One- to three-person team (approval needed beyond 3)
  - Two milestones + demo/report
  - Milestone #1 due in 4 weeks, right after fall break

## SQL

- ❖ SQL: Structured Query Language
  - Pronounced “S-Q-L” or “sequel”
  - The standard query language supported by most commercial DBMS
- ❖ A brief history
  - IBM System R
  - ANSI SQL89
  - ANSI SQL92 (SQL2)
  - ANSI SQL99 (SQL3)
  - ANSI SQL 2003 (added OLAP, XML, etc.)
  - ANSI SQL 2006 (added more XML)

## Creating and dropping tables

- ❖ CREATE TABLE *table\_name* (... , *column\_name*, *column\_type*, ...);
- ❖ DROP TABLE *table\_name*;
- ❖ Examples

```
create table Student (SID integer,
                    name varchar(30), email varchar(30),
                    age integer, GPA float);
create table Course (CID char(10), title varchar(100));
create table Enroll (SID integer, CID char(10));
drop table Student;
drop table Course;
drop table Enroll;
-- everything from -- to the end of the line is ignored.
-- SQL is insensitive to white space.
-- SQL is insensitive to case (e.g., ...Course... is equivalent to
-- ...COURSE...)
```

## Basic queries: SFW statement

- ❖ SELECT  $A_1, A_2, \dots, A_n$   
FROM  $R_1, R_2, \dots, R_m$   
WHERE *condition*;
- ❖ Also called an SPJ (select-project-join) query
- ❖ Equivalent (not really!) to relational algebra query  
 $\pi_{A_1, A_2, \dots, A_n} (\sigma_{condition} (R_1 \times R_2 \times \dots \times R_m))$

## Example: reading a table

- ❖ SELECT \* FROM Student;
  - Single-table query, so no cross product here
  - WHERE clause is optional
  - \* is a short hand for “all columns”

## Example: selection and projection

7

- ❖ Name of students under 18
  - SELECT name FROM Student WHERE age < 18;
- ❖ When was Lisa born?
  - SELECT 2009 - age  
FROM Student  
WHERE name = 'Lisa';
  - SELECT list can contain expressions
    - Can also use built-in functions such as SUBSTR, ABS, etc.
  - String literals (case sensitive) are enclosed in single quotes

## Example: join

8

- ❖ SID's and names of students taking courses with the word "Database" in their titles
  - SELECT Student.SID, Student.name  
FROM Student, Enroll, Course  
WHERE Student.SID = Enroll.SID  
AND Enroll.CID = Course.CID  
AND title LIKE '%Database%';
  - LIKE matches a string against a pattern
    - % matches any sequence of 0 or more characters
  - Okay to omit *table\_name* in *table\_name.column\_name* if *column\_name* is unique

## Example: rename

9

- ❖ SID's of all pairs of classmates
  - Relational algebra query:  
 $\pi_{e1.SID, e2.SID} (\rho_{e1} Enroll \bowtie_{e1.CID = e2.CID \wedge e1.SID > e2.SID} \rho_{e2} Enroll)$
  - SQL:  
SELECT e1.SID AS SID1, e2.SID AS SID2  
FROM Enroll AS e1, Enroll AS e2  
WHERE e1.CID = e2.CID  
AND e1.SID > e2.SID;
  - AS keyword is completely optional

## A more complicated example

10

- ❖ Titles of all courses that Bart and Lisa are taking together
  - SELECT c.title  
FROM Student sb, Student sl, Enroll eb, Enroll el, Course c  
WHERE sb.name = 'Bart' AND sl.name = 'Lisa'  
AND eb.SID = sb.SID AND el.SID = sl.SID  
AND eb.CID = c.CID AND el.CID = c.CID;
  - Tip: Write the FROM clause first, then WHERE, and then SELECT

## Why SFW statements?

11

- ❖ Out of many possible ways of structuring SQL statements, why did the designers choose SELECT-FROM-WHERE?
  - A large number of queries can be written using only selection, projection, and cross product (or join)
  - Any query that uses only these operators can be written in a canonical form:  $\pi_L (\sigma_p (R_1 \times \dots \times R_m))$ 
    - Example:  $\pi_{R.A, S.B} (R \bowtie_{p1} S) \bowtie_{p2} (\pi_{T.C} \sigma_{p3} T) = \pi_{R.A, S.B, T.C} \sigma_{p1 \wedge p2 \wedge p3} (R \times S \times T)$
  - SELECT-FROM-WHERE captures this canonical form

## Set versus bag semantics

12

- ❖ Set
  - No duplicates
  - Relational model and algebra use set semantics
- ❖ Bag
  - Duplicates allowed
  - Number of duplicates is significant
  - SQL uses bag semantics by default

## Set versus bag example

13

Enroll		$\pi_{SID} Enroll$
SID	CID	SID
142	CPS116	142
142	CPS114	123
123	CPS116	857
857	CPS116	456
857	CPS130	...
456	CPS114	...
...	...	...

SELECT SID  
FROM Enroll;

SID
142
142
123
857
857
456
...

## A case for bag semantics

14

- ❖ Efficiency
  - Saves time of eliminating duplicates
- ❖ Which one is more useful?
  - $\pi_{GPA} Student$
  - SELECT GPA FROM Student;
  - The first query just returns all possible GPA's
  - The second query returns the actual GPA distribution
- ❖ Besides, SQL provides the option of set semantics with DISTINCT keyword

## Forcing set semantics

15

- ❖ SID's of all pairs of classmates
  - SELECT e1.SID AS SID1, e2.SID AS SID2  
FROM Enroll AS e1, Enroll AS e2  
WHERE e1.CID = e2.CID  
AND e1.SID > e2.SID;
    - Say Bart and Lisa both take CPS116 and CPS114
  - SELECT DISTINCT e1.SID AS SID1, e2.SID AS SID2  
...
  - With DISTINCT, all duplicate (SID1, SID2) pairs are removed from the output

## Operational semantics of SFW

16

- ❖ SELECT {DISTINCT}  $E_1, E_2, \dots, E_n$   
FROM  $R_1, R_2, \dots, R_m$   
WHERE *condition*;
- ❖ For each  $t_1$  in  $R_1$ :  
For each  $t_2$  in  $R_2$ : ... ..  
For each  $t_m$  in  $R_m$ :  
If *condition* is true over  $t_1, t_2, \dots, t_m$ :  
Compute and output  $E_1, E_2, \dots, E_n$  as a row  
If DISTINCT is present  
Eliminate duplicate rows in output
- ❖  $t_1, t_2, \dots, t_m$  are often called tuple variables

## SQL set and bag operations

17

- ❖ UNION, EXCEPT, INTERSECT
  - Set semantics
    - Duplicates in input tables, if any, are first eliminated
  - Exactly like set  $\cup$ ,  $-$ , and  $\cap$  in relational algebra
- ❖ UNION ALL, EXCEPT ALL, INTERSECT ALL
  - Bag semantics
    - Think of each row as having an implicit count (the number of times it appears in the table)
    - Bag union: sum up the counts from two tables
    - Bag difference: proper-subtract the two counts
    - Bag intersection: take the minimum of the two counts

## Examples of bag operations

18

Bag1	Bag2
fruit	fruit
apple	apple
apple	orange
orange	orange

Bag1 UNION ALL Bag2

fruit
apple
apple
orange
apple
orange
orange

Bag1 INTERSECT ALL Bag2

fruit
apple
orange

Bag1 EXCEPT ALL Bag2

fruit
apple

## Examples of set versus bag operations <sup>19</sup>

❖ *Enroll(SID, CID), ClubMember(club, SID)*

- (SELECT SID FROM ClubMember)  
EXCEPT  
(SELECT SID FROM Enroll);
  - SID's of students who are in clubs but not taking any classes
- (SELECT SID FROM ClubMember)  
EXCEPT ALL  
(SELECT SID FROM Enroll);
  - SID's of students who are in more clubs than classes

## Summary of SQL features covered so far <sup>20</sup>

- ❖ SELECT-FROM-WHERE statements (select-project-join queries)
- ❖ Set and bag operations
- ☞ Next: how to nest SQL queries

## Table expression <sup>21</sup>

- ❖ Use query result as a table
  - In set and bag operations, FROM clauses, etc.
  - A way to “nest” queries
- ❖ Example: names of students who are in more clubs than classes

```
SELECT DISTINCT name
FROM Student,
  ((SELECT SID FROM ClubMember)
  EXCEPT ALL
  (SELECT SID FROM Enroll)) AS S
WHERE Student.SID = S.SID;
```

## Scalar subqueries <sup>22</sup>

- ❖ A query that returns a single row can be used as a value in WHERE, SELECT, etc.

❖ Example: students at the same age as Bart

```
SELECT *           What's Bart's age?
FROM Student
WHERE age = (SELECT age
             FROM Student
             WHERE name = 'Bart');
```

- ❖ Runtime error if subquery returns more than one row
  - Under what condition will this runtime error never occur?
    - *name* is a key of *Student*
- ❖ What if subquery returns no rows?
  - The return value is treated as a special value NULL, and the comparison fails
- ❖ Can be used in SELECT to compute a value for an output column

## IN subqueries <sup>23</sup>

- ❖  $x$  IN (*subquery*) checks if  $x$  is in the result of *subquery*
- ❖ Example: students at the same age as (some) Bart

```
SELECT *           What's Bart's age?
FROM Student
WHERE age IN (SELECT age
             FROM Student
             WHERE name = 'Bart');
```

## EXISTS subqueries <sup>24</sup>

- ❖ EXISTS (*subquery*) checks if the result of *subquery* is non-empty

❖ Example: students at the same age as (some) Bart

```
▪ SELECT *
  FROM Student AS s ←
  WHERE EXISTS (SELECT * FROM Student
               WHERE name = 'Bart'
               AND age = s.age);
```

- This happens to be a correlated subquery—a subquery that references tuple variables in surrounding queries

## Operational semantics of subqueries <sup>25</sup>

- ❖ 

```
SELECT *
FROM Student AS s
WHERE EXISTS (SELECT * FROM Student
              WHERE name = 'Bart'
              AND age = s.age);
```
- ❖ For each row *s* in Student
  - Evaluate the subquery with the appropriate value of *s.age*
  - If the result of the subquery is not empty, output *s.\**
- ❖ The DBMS query optimizer may choose to process the query in an equivalent, but more efficient way (example?)

## Scoping rule of subqueries <sup>26</sup>

- ❖ To find out which table a column belongs to
  - Start with the immediately surrounding query
  - If not found, look in the one surrounding that; repeat if necessary
- ❖ Use *table\_name.column\_name* notation and AS (renaming) to avoid confusion

## Another example <sup>27</sup>

```
SELECT * FROM Student s
WHERE EXISTS
  (SELECT * FROM Enroll e
   WHERE [SID] = s.SID
   AND EXISTS
     (SELECT * FROM Enroll
      WHERE [SID] = [s.SID]
      AND CID <> e.CID));
```

Students who are taking at least two courses

## Quantified subqueries <sup>28</sup>

- ❖ A quantified subquery can be used as a value in a WHERE condition
- ❖ Universal quantification (for all):  
... WHERE *x op ALL (subquery)* ...
  - True iff for all *t* in the result of *subquery*, *x op t*
- ❖ Existential quantification (exists):  
... WHERE *x op ANY (subquery)* ...
  - True iff there exists some *t* in the result of *subquery* such that *x op t*
- ☞ Beware
  - In common parlance, "any" and "all" seem to be synonyms
  - In SQL, ANY really means "some"

## Examples of quantified subqueries <sup>29</sup>

- ❖ Which students have the highest GPA?
  - ```
SELECT *
FROM Student
WHERE GPA >= ALL
  (SELECT GPA FROM Student);
```
  - ```
SELECT *
FROM Student
WHERE NOT
  (GPA < ANY (SELECT GPA FROM Student));
```
- ☞ Use NOT to negate a condition

## More ways of getting the highest GPA <sup>30</sup>

- ❖ Which students have the highest GPA?
  - ```
SELECT *
FROM Student AS s
WHERE NOT EXISTS
  (SELECT * FROM Student
   WHERE GPA > s.GPA);
```
  - ```
SELECT * FROM Student
WHERE SID NOT IN
  (SELECT s1.SID
   FROM Student AS s1, Student AS s2
   WHERE s1.GPA < s2.GPA);
```

## Summary of SQL features covered so far <sup>31</sup>

- ❖ SELECT-FROM-WHERE statements
- ❖ Set and bag operations
- ❖ Table expressions, subqueries
  - Subqueries allow queries to be written in more declarative ways (recall the highest GPA query)
  - But they do not add much expressive power
    - Try translating other forms of subqueries into [NOT] EXISTS, which in turn can be translated into join (and difference)

☞ Next: aggregation and grouping

## Aggregates <sup>32</sup>

- ❖ Standard SQL aggregate functions: COUNT, SUM, AVG, MIN, MAX
- ❖ Example: number of students under 18, and their average GPA
  - SELECT COUNT(\*), AVG(GPA)  
FROM Student  
WHERE age < 18;
  - COUNT(\*) counts the number of rows

## Aggregates with DISTINCT <sup>33</sup>

- ❖ Example: How many students are taking classes?

- SELECT COUNT(DISTINCT SID)  
FROM Enroll;

is equivalent to:

- SELECT COUNT(\*)  
FROM (SELECT DISTINCT SID FROM Enroll);

## GROUP BY <sup>34</sup>

- ❖ SELECT ... FROM ... WHERE ...  
GROUP BY *list\_of\_columns*;
- ❖ Example: find the average GPA for each age group
  - SELECT age, AVG(GPA)  
FROM Student  
GROUP BY age;

## Operational semantics of GROUP BY <sup>35</sup>

SELECT ... FROM ... WHERE ... GROUP BY ...;

- ❖ Compute FROM ( $\times$ )
  - ❖ Compute WHERE ( $\sigma$ )
  - ❖ Compute GROUP BY: group rows according to the values of GROUP BY columns
  - ❖ Compute SELECT for each group ( $\pi$ )
    - For aggregation functions with DISTINCT inputs, first eliminate duplicates within the group
- ☞ Number of groups = number of rows in the final output

## Example of computing GROUP BY <sup>36</sup>

SELECT age, AVG(GPA) FROM Student GROUP BY age;

SID	name	age	GPA
142	Bart	10	2.3
857	Lisa	8	4.3
123	Wilhouse	10	3.1
456	Ralph	8	2.3
--	--	--	--

Compute GROUP BY: group rows according to the values of GROUP BY columns



SID	name	age	GPA
142	Bart	10	2.3
123	Wilhouse	10	3.1
857	Lisa	8	4.3
456	Ralph	8	2.3
--	--	--	--

Compute SELECT for each group



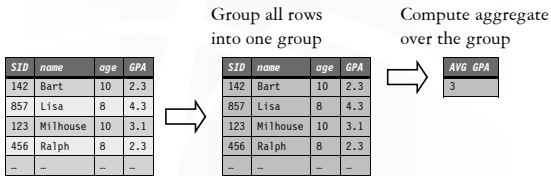
age	AVG GPA
10	2.7
8	3.3
--	--

## Aggregates with no GROUP BY

37

- ❖ An aggregate query with no GROUP BY clause represent a special case where all rows go into one group

```
SELECT AVG(GPA) FROM Student;
```



## Restriction on SELECT

38

- ❖ If a query uses aggregation/group by, then every column referenced in SELECT must be either
  - Aggregated, or
  - A GROUP BY column
- ☞ This restriction ensures that any SELECT expression produces only one value for each group

## Examples of invalid queries

39

- ❖ ~~SELECT~~ ~~SID~~, age FROM Student GROUP BY age;
  - Recall there is one output row per group
  - There can be multiple SID values per group
- ❖ ~~SELECT~~ ~~SID~~, MAX(GPA) FROM Student;
  - Recall there is only one group for an aggregate query with no GROUP BY clause
  - There can be multiple SID values
  - Wishful thinking (that the output SID value is the one associated with the highest GPA) does NOT work
  - ☞ Another way of writing the max GPA query?

## HAVING

40

- ❖ Used to filter groups based on the group properties (e.g., aggregate values, GROUP BY column values)
- ❖ SELECT ... FROM ... WHERE ... GROUP BY ... HAVING *condition*;
  - Compute FROM ( $\times$ )
  - Compute WHERE ( $\sigma$ )
  - Compute GROUP BY: group rows according to the values of GROUP BY columns
  - Compute HAVING (another  $\sigma$  over the groups)
  - Compute SELECT ( $\pi$ ) for each group that passes HAVING

## HAVING examples

41

- ❖ Find the average GPA for each age group over 10
  - SELECT age, AVG(GPA)  
FROM Student  
GROUP BY age  
HAVING age > 10;
  - Can be written using WHERE without table expressions
- ❖ List the average GPA for each age group with more than a hundred students
  - SELECT age, AVG(GPA)  
FROM Student  
GROUP BY age  
HAVING COUNT(\*) > 100;
  - Can be written using WHERE and table expressions

## Summary of SQL features covered so far

42

- ❖ SELECT-FROM-WHERE statements
- ❖ Set and bag operations
- ❖ Table expressions, subqueries
- ❖ Aggregation and grouping
  - More expressive power than relational algebra
- ☞ Next: ordering output rows

## ORDER BY <sup>43</sup>

- ❖ SELECT {DISTINCT} ...  
FROM ... WHERE ... GROUP BY ... HAVING ...  
ORDER BY *output\_column* [ASC | DESC], ...;
- ❖ ASC = ascending, DESC = descending
- ❖ Operational semantics
  - After SELECT list has been computed and optional duplicate elimination has been carried out, sort the output according to ORDER BY specification

## ORDER BY example <sup>44</sup>

- ❖ List all students, sort them by GPA (descending) and name (ascending)
  - SELECT SID, name, age, GPA  
FROM Student  
ORDER BY GPA DESC, name;
  - ASC is the default option
  - Strictly speaking, only output columns can appear in ORDER BY clause (although some DBMS support more)
  - Can use sequence numbers instead of names to refer to output columns: ORDER BY 4 DESC, 2;

## Summary of SQL features covered so far <sup>45</sup>

- ❖ SELECT-FROM-WHERE statements
  - ❖ Set and bag operations
  - ❖ Table expressions, subqueries
  - ❖ Aggregation and grouping
  - ❖ Ordering
- ☞ Next: NULL's, outerjoins, data modification, constraints, ...