

Relational Database Design Part I

CPS 116

Introduction to Database Systems

Announcements (Thu. Sep. 8)²

- ❖ Homework #1 is out
 - Due in 12 days, but start early—as soon as any portion has been covered
- ❖ Sign up for Gradiance
- ❖ Try our VM now
 - We can't help you on the due date if you run into last-minute installation problems
- ❖ Duke CS account created
 - Email me immediately if you still need one
- ❖ Readings: see Tentative Syllabus on course website

Relational model: review³

- ❖ A database is a collection of relations (or tables)
- ❖ Each relation has a list of attributes (or columns)
- ❖ Each attribute has a domain (or type)
- ❖ Each relation contains a set of tuples (or rows)

Keys⁴

- ❖ A set of attributes K is a key for a relation R if
 - In no instance of R will two different tuples agree on all attributes of K
 - That is, K is a “tuple identifier”
 - No proper subset of K satisfies the above condition
 - That is, K is minimal
- ❖ Example: *Student* (*SID*, *name*, *age*, *GPA*)
 - *SID* is a key of *Student*
 - *age* is not a key (not an identifier)
 - $\{\text{SID}, \text{name}\}$ is not a key (not minimal)

Schema vs. instance⁵

Student

<i>SID</i>	<i>name</i>	<i>age</i>	<i>GPA</i>
142	Bart	10	2.3
123	Milhouse	10	3.1
857	Lisa	8	4.3
456	Ralph	8	2.3

- ❖ Is *name* a key of *Student*?
 - Yes? Seems reasonable for this instance
 - No! Student names are not unique in general
- ❖ Key declarations are part of the schema

More examples of keys⁶

- ❖ *Enroll* (*SID*, *CID*)
 - $\{\text{SID}, \text{CID}\}$
 - ❖ A key can contain multiple attributes!
- ❖ *Address* (*street_address*, *city*, *state*, *zip*)
 - $\{\text{street_address}, \text{city}, \text{state}\}$
 - $\{\text{street_address}, \text{zip}\}$
 - ❖ A relation can have multiple keys!
 - We typically pick one as the “primary” key, and underline all its attributes, e.g., *Address* (*street_address*, *city*, *state*, *zip*)

7 Usage of keys

- ❖ More constraints on data, fewer mistakes
- ❖ Look up a row by its key value
 - Many selection conditions are “key = value”
- ❖ “Pointers”
 - Example: *Enroll (SID, CID)*
 - *SID* is a key of *Student*
 - *CID* is a key of *Course*
 - An *Enroll* tuple “links” a *Student* tuple with a *Course* tuple
 - Many join conditions are “key = key value stored in another table”

8 Database design

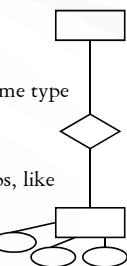
- ❖ Understand the real-world domain being modeled
- ❖ Specify it using a database design model
 - More intuitive and convenient for schema design
 - But not necessarily implemented by DBMS
- ❖ A few popular ones:
 - Entity/Relationship (E/R) model
 - Object Definition Language (ODL)
 - UML (Unified Modeling Language)
- ❖ Translate specification to the data model of DBMS
 - Relational, XML, object-oriented, etc.
- ❖ Create DBMS schema

9 Entity-relationship (E/R) model

- ❖ Historically and still very popular
- ❖ Can think of as a “watered-down” object-oriented design model
- ❖ Primarily a design model—not directly implemented by DBMS
- ❖ Designs represented by E/R diagrams
 - We use the style of E/R diagram covered by GMUW; there are other styles/extensions
 - Very similar to UML diagrams

10 E/R basics

- ❖ Entity: a “thing,” like an object
- ❖ Entity set: a collection of things of the same type, like a relation of tuples or a class of objects
 - Represented as a rectangle
- ❖ Relationship: an association among entities
- ❖ Relationship set: a set of relationships of the same type (among same entity sets)
 - Represented as a diamond
- ❖ Attributes: properties of entities or relationships, like attributes of tuples or objects
 - Represented as ovals



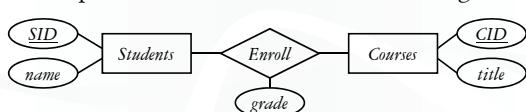
11 An example E/R diagram

- ❖ Students enroll in courses
- ❖ A key of an entity set is represented by underlining all attributes in the key
 - A key is a set of attributes whose values can belong to at most one entity in an entity set—like a key of a relation



12 Attributes of relationships

- ❖ Example: students take courses and receive grades
- ❖ Where do the grades go?
 - With *Students*?
 - But a student can have different grades for multiple courses
 - With *Courses*?
 - But a course can assign different grades for multiple students
 - With *Enroll*!



More on relationships

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- ❖ There could be multiple relationship sets between the same entity sets
 - Example: *Students Enroll Courses; Students TA Courses*
- ❖ In a relationship set, each relationship is uniquely identified by the entities it connects
 - Example: Between Bart and CPS116, there can be at most one *Enroll* relationship and at most one *TA* relationship
- ❖ What if Bart took CPS116 twice and got two different grades?

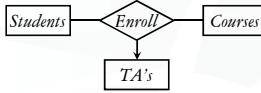
Multiplicity of relationships

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- ❖ E and F : entity sets
- ❖ Many-many: Each entity in E is related to 0 or more entities in F and vice versa
 - Example: 
- ❖ Many-one: Each entity in E is related to 0 or 1 entity in F , but each entity in F is related to 0 or more in E
 - Example: 
- ❖ One-one: Each entity in E is related to 0 or 1 entity in F and vice versa
 - Example: 
- ❖ “One” (0 or 1) is represented by an arrow →
- ❖ “Exactly one” is represented by a rounded arrow →

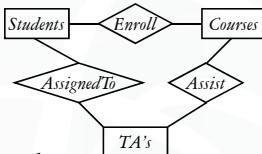
N-ary relationships

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- ❖ Example: Each course has multiple TA's; each student is assigned to one TA
 
- ❖ Meaning of an arrow into E : Pick one entity from each of the other entity sets; together they must be related to either 0 or 1 entity in E

N-ary versus binary relationships

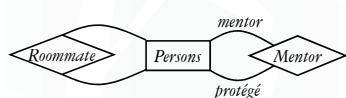
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- ❖ Can we model n -ary relationships using just binary relationships?
 
- ❖ No; for example:
 - Bart takes CPS116 and CPS114
 - Lisa TA's CPS116 and CPS114
 - Bart is assigned to Lisa in CPS116, but not in CPS114

Roles in relationships

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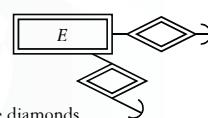
- ❖ An entity set may participate more than once in a relationship set
- ❖ May need to label edges to distinguish roles
- ❖ Examples
 - People mentor others; label needed
 - People are roommates of each other; label not needed



Weak entity sets

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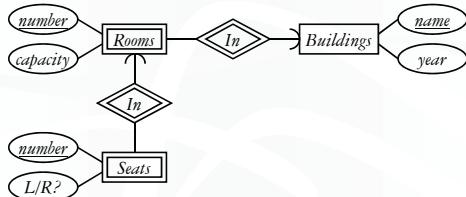
- ❖ Sometimes, the key of an entity set E comes not completely from its own attributes, but from the keys of other (one or more) entity sets
 - E must link to them via many-one (or one-one) relationship sets
 - Example: *Rooms inside Buildings* are partly identified by *Buildings'* name
 - E is called a weak entity set
 - Denoted by double rectangle
 - The relationship sets through which E obtains its key are drawn as double diamonds



Weak entity set examples

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- Seats in rooms in buildings



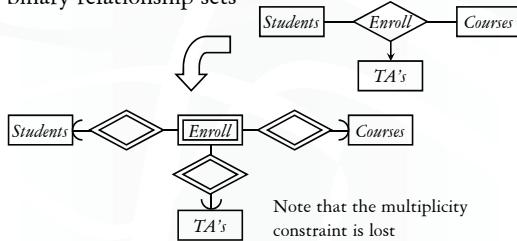
- Why must double diamonds be many-one/one-one?

- With many-many, we would not know which entity provides the key value!

Modeling n -ary relationships

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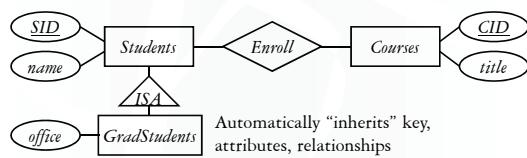
- An n -ary relationship set can be replaced by a weak entity set (called a connecting entity set) and n binary relationship sets



ISA relationships

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- Similar to the idea of subclasses in object-oriented programming: subclass = special case, fewer entities, and possibly more properties
 - Represented as a triangle (direction is important)
- Example: Graduate students are students, but they also have offices



Summary of E/R concepts

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- Entity sets
 - Keys
 - Weak entity sets
- Relationship sets
 - Attributes of relationships
 - Multiplicity
 - Roles
 - Binary versus n -ary relationships
 - Modeling n -ary relationships with weak entity sets and binary relationships
 - ISA relationships

Case study 1

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- Design a database representing cities, counties, and states
 - For states, record name and capital (city)
 - For counties, record name, area, and location (state)
 - For cities, record name, population, and location (county and state)
- Assume the following:
 - Names of states are unique
 - Names of counties are only unique within a state
 - Names of cities are only unique within a county
 - A city is always located in a single county
 - A county is always located in a single state

Case study 1: first design

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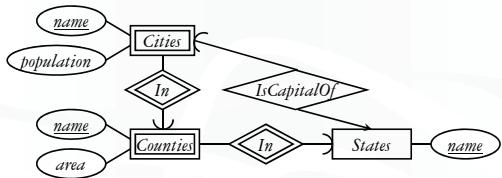
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- ```

 erDiagram
 class Cities {
 string name;
 number population;
 string county_name;
 number county_area;
 }
 class States {
 string name;
 string capital;
 }
 Cities }o--o{ States : In
 }

```
- County area information is repeated for every city in the county
    - Redundancy is bad (why?)
  - State capital should really be a city
    - Should "reference" entities through explicit relationships

## Case study 1: second design

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- ❖ Technically, nothing in this design could prevent a city in state X from being the capital of another state Y, but oh well...

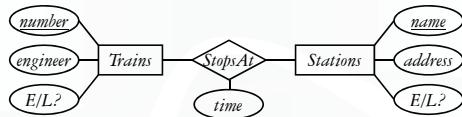
## Case study 2

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- ❖ Design a database consistent with the following:
  - A station has a unique name and an address, and is either an express station or a local station
  - A train has a unique number and an engineer, and is either an express train or a local train
  - A local train can stop at any station
  - An express train only stops at express stations
  - A train can stop at a station for any number of times during a day
  - Train schedules are the same everyday

## Case study 2: first design

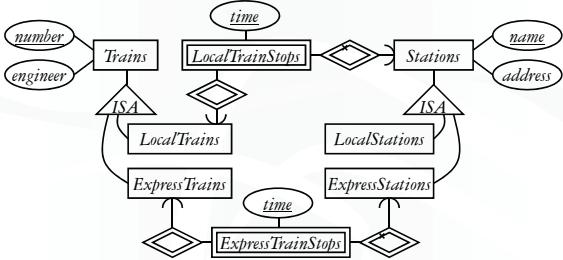
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- ❖ Nothing in this design prevents express trains from stopping at local stations
  - ❖ We should capture as many constraints as possible
- ❖ A train can stop at a station only once during a day
  - ❖ We should not introduce constraints

## Case study 2: second design

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Is the extra complexity worth it?