SQL: Part I

CPS 116

Introduction to Database Systems

Announcements (Tue. Sep. 20)

- Homework #1 due today
  - Deadline for the Gradiance portion has been extended
- Homework #2 will be assigned Thursday

SQL

- SQL: Structured Query Language
  - Pronounced "S-Q-L" or "sequel"
  - The standard query language supported by most commercial DBMS
- A brief history
  - IBM System R
  - ANSI SQL89
  - ANSI SQL92 (SQL2)
  - ANSI SQL99 (SQL3)
  - ANSI SQL 2003 (added OLAP, XML, etc.)
  - ANSI SQL 2006 (added more XML)
  - ANSI SQL 2008, …

Creating and dropping tables

4

- CREATE TABLE *table_name* (*…*, *column_name*$_i$ *column_type*$_i$, *…*);
- DROP TABLE *table_name*;
- Examples

```
create table Student (SID integer,
                      name varchar(30), email varchar(30),
                      age integer, GPA float);
create table Course (CID char(10), title varchar(100));
create table Enroll (SID integer, CID char(10));
drop table Student;
drop table Course;
drop table Enroll;
-- everything from -- to the end of the line is ignored.
-- SQL is insensitive to white space.
-- SQL is insensitive to case (e.g., ...Course... is equivalent to
-- ...COURSE...)
```

Basic queries: SFW statement

5

- SELECT A_1, A_2, …, A_n
  FROM R_1, R_2, …, R_m
  WHERE *condition*;
- Also called an SPJ (select-project-join) query
- Equivalent (not really!) to relational algebra query
  $\pi_{A_1, A_2, \ldots, A_n} (\sigma_{condition} (R_1 \times R_2 \times \ldots \times R_m))$

Example: reading a table

6

- SELECT * FROM Student;
  - Single-table query, so no cross product here
  - WHERE clause is optional
  - * is a short hand for "all columns"

Example: selection and projection

7

- Name of students under 18
  - `SELECT name FROM Student WHERE age < 18;`
- When was Lisa born?
  - ```
    SELECT 2009 – age
    FROM Student
    WHERE name = 'Lisa';
    ```
  - `SELECT` list can contain expressions
    - Can also use built-in functions such as `SUBSTR`, `ABS`, etc.
  - String literals (case sensitive) are enclosed in single quotes

Example: join

8

- SID's and names of students taking courses with the word "Database" in their titles
  - ```
    SELECT Student.SID, Student.name
    FROM Student, Enroll, Course
    WHERE Student.SID = Enroll.SID
    AND Enroll.CID = Course.CID
    AND title LIKE '%Database%';
    ```
  - `LIKE` matches a string against a pattern
    - `%` matches any sequence of 0 or more characters
  - Okay to omit *table_name* in *table_name*`.`*column_name* if *column_name* is unique

Example: rename

9

- SID's of all pairs of classmates
  - Relational algebra query:
  - SQL:
    ```
    SELECT e1.SID AS SID1, e2.SID AS SID2
    FROM Enroll AS e1, Enroll AS e2
    WHERE e1.CID = e2.CID
    AND e1.SID > e2.SID;
    ```
  - `AS` keyword is completely optional

A more complicated example

10

- Titles of all courses that Bart and Lisa are taking together

```
SELECT
FROM
WHERE
```

Tip: Write the FROM clause first, then WHERE, and then SELECT

Why SFW statements?

11

- Out of many possible ways of structuring SQL statements, why did the designers choose SELECT-FROM-WHERE?
  - A large number of queries can be written using only selection, projection, and cross product (or join)
  - Any query that uses only these operators can be written in a canonical form: $\pi_L (\sigma_p (R_1 \times \ldots \times R_m))$
    - Example: $\pi_{R.A, S.B} (R \not\bowtie_{p1} S) \not\bowtie_{p2} (\pi_{T.C} \sigma_{p3} T) = \pi_{R.A, S.B, T.C} \sigma_{p1 \wedge p2 \wedge p3} (R \times S \times T)$
  - SELECT-FROM-WHERE captures this canonical form

Set versus bag semantics

12

- Set
  - No duplicates
  - Relational model and algebra use set semantics
- Bag
  - Duplicates allowed
  - Number of duplicates is significant
  - SQL uses bag semantics by default

Set versus bag example

13

Enroll

SID	CID
142	CPS116
142	CPS114
123	CPS116
857	CPS116
857	CPS130
456	CPS114
…	…

π_{SID} *Enroll*

SID
142
123
857
456
…

```
SELECT SID
FROM Enroll;
```

SID
142
142
123
857
857
456
…

A case for bag semantics

14

- Besides, SQL provides the option of set semantics with `DISTINCT` keyword

Forcing set semantics

15

- SID's of all pairs of classmates
  - ```
    SELECT e1.SID AS SID1, e2.SID AS SID2
    FROM Enroll AS e1, Enroll AS e2
    WHERE e1.CID = e2.CID
    AND e1.SID > e2.SID;
    ```
  - `SELECT DISTINCT e1.SID AS SID1, e2.SID AS SID2`
    `...`
    - With `DISTINCT`, all duplicate (`SID1`, `SID2`) pairs are removed from the output
    ```

Operational semantics of SFW

16

- SELECT [DISTINCT] E_1, E_2, …, E_n
  FROM R_1, R_2, …, R_m
  WHERE *condition*;
- For each t_1 in R_1:
  For each t_2 in R_2: … …
  For each t_m in R_m:
  If *condition* is true over $t_1, t_2, \ldots, t_m$:
  Compute and output $E_1, E_2, \ldots, E_n$ as a row
  If DISTINCT is present
  Eliminate duplicate rows in output
- $t_1, t_2, \ldots, t_m$ are often called tuple variables

SQL set and bag operations

17

- UNION, EXCEPT, INTERSECT
  - Set semantics
    - Duplicates in input tables, if any, are first eliminated
  - Exactly like set ∪, −, and ∩ in relational algebra
- UNION ALL, EXCEPT ALL, INTERSECT ALL
  - Bag semantics
  - Think of each row as having an implicit count (the number of times it appears in the table)
  - Bag union: sum up the counts from two tables
  - Bag difference: proper-subtract the two counts
  - Bag intersection: take the minimum of the two counts

Examples of bag operations

18

Bag1

fruit
apple
apple
orange

Bag2

fruit
apple
orange
orange

Bag1 UNION ALL Bag2

fruit
apple
apple
orange
apple
orange
orange

Bag1 EXCEPT ALL Bag2

fruit
apple

Bag1 INTERSECT ALL Bag2

fruit
apple
orange

Examples of set versus bag operations

19

- *Enroll*(*SID*, *CID*), *ClubMember*(*club*, *SID*)
  - `(SELECT SID FROM ClubMember) EXCEPT (SELECT SID FROM Enroll);`
  - `(SELECT SID FROM ClubMember) EXCEPT ALL (SELECT SID FROM Enroll);`

Summary of SQL features covered so far

20

- `SELECT-FROM-WHERE` statements (select-project-join queries)
- Set and bag operations

☞ Next: how to nest SQL queries

Table expression

21

- Use query result as a table
  - In set and bag operations, `FROM` clauses, etc.
  - A way to "nest" queries
- Example: names of students who are in more clubs than classes

```
SELECT DISTINCT name
FROM Student,
     ((SELECT SID FROM ClubMember)
      EXCEPT ALL
      (SELECT SID FROM Enroll)) AS S
WHERE Student.SID = S.SID;
```

Scalar subqueries

- A query that returns a single row can be used as a value in `WHERE`, `SELECT`, etc.
- Example: students at the same age as Bart

```
SELECT *                    What's Bart's age?
FROM Student
WHERE age = (SELECT age
             FROM Student
             WHERE name = 'Bart');
```

- Runtime error if subquery returns more than one row
  - Under what condition will this runtime error never occur?
- What if subquery returns no rows?
  - The return value is treated as a special value `NULL`, and the comparison fails
- Can be used in `SELECT` to compute a value for an output column

`IN` subqueries

- *x* `IN (`*subquery*`)` checks if *x* is in the result of *subquery*
- Example: students at the same age as (some) Bart

```
SELECT *                    What's Bart's age?
FROM Student
WHERE age IN (SELECT age
              FROM Student
              WHERE name = 'Bart');
```

`EXISTS` subqueries

- `EXISTS (`*subquery*`)` checks if the result of *subquery* is non-empty
- Example: students at the same age as (some) Bart
  -
```
SELECT *
FROM Student AS s
WHERE EXISTS (SELECT * FROM Student
              WHERE name = 'Bart'
              AND age = s.age);
```
  - This happens to be a correlated subquery—a subquery that references tuple variables in surrounding queries

Operational semantics of subqueries

25

- ```
  SELECT *
  FROM Student AS s
  WHERE EXISTS (SELECT * FROM Student
                WHERE name = 'Bart'
                AND age = s.age);
  ```
- For each row `s` in `Student`
  - Evaluate the subquery with the appropriate value of `s.age`
  - If the result of the subquery is not empty, output `s.*`
- The DBMS query optimizer may choose to process the query in an equivalent, but more efficient way (example?)

Scoping rule of subqueries

26

- To find out which table a column belongs to
  - Start with the immediately surrounding query
  - If not found, look in the one surrounding that; repeat if necessary
- Use *table_name*`.`*column_name* notation and `AS` (renaming) to avoid confusion

Another example

27

```
SELECT * FROM Student s
WHERE EXISTS
    (SELECT * FROM Enroll e
     WHERE SID = s.SID
     AND EXISTS
         (SELECT * FROM Enroll
          WHERE SID = s.SID
          AND CID <> e.CID));
```
```

Quantified subqueries

28

- A quantified subquery can be used as a value in a `WHERE` condition
- Universal quantification (for all):
  … `WHERE` *x op* `ALL` (*subquery*) …
  - True iff for all *t* in the result of *subquery*, *x op t*
- Existential quantification (exists):
  … `WHERE` *x op* `ANY` (*subquery*) …
  - True iff there exists some *t* in the result of *subquery* such that *x op t*
  - ☞Beware
    - In common parlance, "any" and "all" seem to be synonyms
    - In SQL, `ANY` really means "some"

Examples of quantified subqueries

29

- Which students have the highest GPA?
  - ```
    SELECT *
    FROM Student
    WHERE GPA >= ALL
        (SELECT GPA FROM Student);
    ```
  - ```
    SELECT *
    FROM Student
    WHERE NOT
        (GPA < ANY (SELECT GPA FROM Student);
    ```
  - ☞Use `NOT` to negate a condition

More ways of getting the highest GPA

30

- Which students have the highest GPA?

Summary of SQL features covered so far

31

- SELECT-FROM-WHERE statements
- Set and bag operations
- Table expressions, subqueries
  - Subqueries allow queries to be written in more declarative ways (recall the highest GPA query)
  - But they do not add much expressive power
    - Try translating other forms of subqueries into [NOT] EXISTS, which in turn can be translated into join (and difference)

☞ Next: aggregation and grouping

Aggregates

32

- Standard SQL aggregate functions: COUNT, SUM, AVG, MIN, MAX
- Example: number of students under 18, and their average GPA
  - ```
    SELECT COUNT(*), AVG(GPA)
    FROM Student
    WHERE age < 18;
    ```
  - COUNT(*) counts the number of rows

Aggregates with DISTINCT

33

- Example: How many students are taking classes?
  - ```
    SELECT COUNT(DISTINCT SID)
    FROM Enroll;
    ```

  is equivalent to:

GROUP BY

34

- SELECT … FROM … WHERE …
  GROUP BY *list_of_columns*;
- Example: find the average GPA for each age group
  - SELECT age, AVG(GPA)
    FROM Student
    GROUP BY age;

Operational semantics of GROUP BY

35

SELECT … FROM … WHERE … GROUP BY …;

- Compute FROM ($\times$)
- Compute WHERE (σ)
- Compute GROUP BY: group rows according to the values of GROUP BY columns
- Compute SELECT for each group (π)
  - For aggregation functions with DISTINCT inputs, first eliminate duplicates within the group

☞ Number of groups = number of rows in the final output

Example of computing GROUP BY

36

SELECT age, AVG(GPA) FROM Student GROUP BY age;

SID	name	age	GPA
142	Bart	10	2.3
857	Lisa	8	4.3
123	Milhouse	10	3.1
456	Ralph	8	2.3
…	…	…	…

Compute GROUP BY: group rows according to the values of GROUP BY columns

⇨

SID	name	age	GPA
142	Bart	10	2.3
123	Milhouse	10	3.1
857	Lisa	8	4.3
456	Ralph	8	2.3
…	…	…	…

Compute SELECT for each group

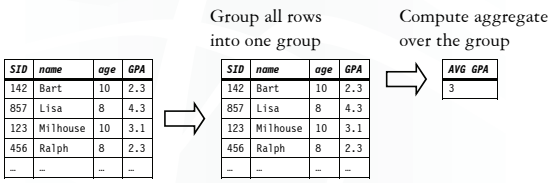
⇦

age	AVG GPA
10	2.7
8	3.3
…	…

Aggregates with no GROUP BY

- An aggregate query with no GROUP BY clause represent a special case where all rows go into one group

```
SELECT AVG(GPA) FROM Student;
```



Restriction on SELECT

- If a query uses aggregation/group by, then every column referenced in SELECT must be either
  - Aggregated, or
  - A GROUP BY column

Examples of invalid queries

- SELECT ~~SID,~~ age FROM Student GROUP BY age;
  - Recall there is one output row per group
  - There can be multiple SID values per group
- SELECT ~~SID,~~ MAX(GPA) FROM Student;
  - Recall there is only one group for an aggregate query with no GROUP BY clause
  - There can be multiple SID values
  - Wishful thinking (that the output SID value is the one associated with the highest GPA) does NOT work

☞Another way of writing the max GPA query?

HAVING

- Used to filter groups based on the group properties (e.g., aggregate values, `GROUP BY` column values)
- `SELECT ... FROM ... WHERE ... GROUP BY ... HAVING` *condition*`;`
  - Compute `FROM` ($\times$)
  - Compute `WHERE` (σ)
  - Compute `GROUP BY`: group rows according to the values of `GROUP BY` columns
  - Compute `HAVING` (another σ over the groups)
  - Compute `SELECT` (π) for each group that passes `HAVING`

HAVING examples

- Find the average GPA for each age group over 10
  - ```
    SELECT age, AVG(GPA)
    FROM Student
    GROUP BY age
    HAVING age > 10;
    ```
  - Can be written using `WHERE` without table expressions
- List the average GPA for each age group with more than a hundred students
  - ```
    SELECT age, AVG(GPA)
    FROM Student
    GROUP BY age
    HAVING COUNT(*) > 100;
    ```
  - Can be written using `WHERE` and table expressions

Summary of SQL features covered so far

- `SELECT-FROM-WHERE` statements
- Set and bag operations
- Table expressions, subqueries
- Aggregation and grouping
  - More expressive power than relational algebra

☞ Next: ordering output rows

ORDER BY

- SELECT [DISTINCT] ...
  FROM ... WHERE ... GROUP BY ... HAVING ...
  ORDER BY *output_column* [ASC | DESC], ...;
- ASC = ascending, DESC = descending
- Operational semantics
  - After SELECT list has been computed and optional duplicate elimination has been carried out, sort the output according to ORDER BY specification

ORDER BY example

- List all students, sort them by GPA (descending) and name (ascending)
  - SELECT SID, name, age, GPA
    FROM Student
    ORDER BY GPA DESC, name;
  - ASC is the default option
  - Strictly speaking, only output columns can appear in ORDER BY clause (although some DBMS support more)
  - Can use sequence numbers instead of names to refer to output columns: ORDER BY 4 DESC, 2;

Summary of SQL features covered so far

- SELECT-FROM-WHERE statements
- Set and bag operations
- Table expressions, subqueries
- Aggregation and grouping
- Ordering

☞ Next: NULL's, outerjoins, data modification, constraints, …