SQL: Part I

Introduction to Databases CompSci 316 Fall 2016



Announcements (Tue. Sep. 20)

- Homework #1 due today 11:59pm
- Homework #2 assigned
- Project mixer next Tuesday in class; details to follow in email

SQL

- SQL: Structured Query Language
 - Pronounced "S-Q-L" or "sequel"
 - The standard query language supported by most DBMS
- A brief history
 - IBM System R
 - ANSI SQL89
 - ANSI SQL92 (SQL2)
 - ANSI SQL99 (SQL3)
 - ANSI SQL 2003 (added OLAP, XML, etc.)
 - ANSI SQL 2006 (added more XML)
 - ANSI SQL 2008, ...

Creating and dropping tables

- CREATE TABLE table_name (..., column_name column_type, ...);
- DROP TABLE table_name;
- Examples

Basic queries: SFW statement

- SELECT A_1 , A_2 , ..., A_n FROM R_1 , R_2 , ..., R_m WHERE condition;
- Also called an SPJ (select-project-join) query
- Corresponds to (but not really equivalent to) relational algebra query:

$$\pi_{A_1,A_2,...,A_n} (\sigma_{condition}(R_1 \times R_2 \times \cdots \times R_m))$$

Example: reading a table

- SELECT * FROM User;
 - Single-table query, so no cross product here
 - WHERE clause is optional
 - * is a short hand for "all columns"

Example: selection and projection

- Name of users under 18
 - SELECT name FROM User WHERE age<18;
- When was Lisa born?
 - SELECT 2016-age FROM User WHERE name = 'Lisa';
 - SELECT list can contain expressions
 - Can also use built-in functions such as SUBSTR, ABS, etc.
 - String literals (case sensitive) are enclosed in single quotes

Example: join

- ID's and names of groups with a user whose name contains "Simpson"
 - SELECT Group.gid, Group.name
 FROM User, Member, Group
 WHERE User.uid = Member.uid
 AND Member.gid = Group.gid
 AND User.name LIKE '%Simpson%';
 - LIKE matches a string against a pattern
 - % matches any sequence of zero or more characters
 - Okay to omit *table_name* in *table_name*. *column_name* if *column_name* is unique

Example: rename

- ID's of all pairs of users that belong to one group
 - Relational algebra query:

```
\pi_{m_1.uid,m_2.uid} (\rho_{m_1}Member\bowtie_{m_1.gid=m_2.gid\land m_1.uid>m_2.uid}\rho_{m_2}Member) • SQL: SELECT ml.uid AS uidl, m2.uid AS uid2
```

FROM Member AS ml, Member AS m2
WHERE ml.gid = m2.gid
AND ml.uid > m2.uid;

AS keyword is completely optional

A more complicated example

Names of all groups that Lisa and Ralph are both in

```
SELECT g.name
FROM User ul, User u2, Member ml, Member m2, Group g
WHERE ul.name = 'Lisa' AND u2.name = 'Ralph'
AND ul.uid = ml.uid AND u2.uid = m2.uid
AND ml.gid = g.gid AND m2.gid = g.gid;
```

Tip: Write the FROM clause first, then WHERE, and then SELECT

Why SFW statements?

- Out of many possible ways of structuring SQL statements, why did the designers choose SELECT-FROM-WHERE?
 - A large number of queries can be written using only selection, projection, and cross product (or join)
 - Any query that uses only these operators can be written in a canonical form: $\pi_L\left(\sigma_p(R_1\times\cdots\times R_m)\right)$
 - Example: $\pi_{R.A,S.B}(R \bowtie_{p_1} S) \bowtie_{p_2} (\pi_{T.C}\sigma_{p_3}T)$ = $\pi_{R.A,S.B,T.C}\sigma_{p_1 \wedge p_2 \wedge p_3}(R \times S \times T)$
 - SELECT-FROM-WHERE captures this canonical form

Set versus bag semantics

- Set
 - No duplicates
 - Relational model and algebra use set semantics
- Bag
 - Duplicates allowed
 - Number of duplicates is significant
 - SQL uses bag semantics by default

Set versus bag example

Member

uid	gid
142	dps
123	gov
857	abc
857	gov
456	abc
456	gov
•••	•••

 $\pi_{gid}Member$ gid dps gov abc

SELECT gid FROM Member;

gid
dps
gov
abc
gov
abc
gov
...

A case for bag semantics

- Efficiency
 - Saves time of eliminating duplicates
- Which one is more useful?
 - $\pi_{a,ge}User$
 - SELECT age FROM User;
 - The first query just returns all possible user ages
 - The second query returns the user age distribution
- Besides, SQL provides the option of set semantics with DISTINCT keyword

Forcing set semantics

- ID's of all pairs of users that belong to one group
 - SELECT ml.uid AS uidl, m2.uid AS uid2 FROM Member AS ml, Member AS m2 WHERE ml.gid = m2.gid AND ml.uid > m2.uid;
 - Say Lisa and Ralph are in both the book club and the student government
 - SELECT DISTINCT ml.uid AS uidl, m2.uid AS uid2 ...
 - With DISTINCT, all duplicate (uid1, uid2) pairs are removed from the output

Semantics of SFW

```
FROM R_1, R_2, ..., R_m
 WHERE condition;
• For each t_1 in R_1:
   For each t_2 in R_2: ...
      For each t_m in R_m:
        If condition is true over t_1, t_2, ..., t_m:
          Compute and output E_1, E_2, ..., E_n as a row
 If DISTINCT is present
   Eliminate duplicate rows in output
```

• SELECT [DISTINCT] E_1 , E_2 , ..., E_n

• $t_1, t_2, ..., t_m$ are often called tuple variables

SQL set and bag operations

- UNION, EXCEPT, INTERSECT
 - Set semantics
 - Duplicates in input tables, if any, are first eliminated
 - Duplicates in result are also eliminated (for UNION)
 - Exactly like set ∪, –, and ∩ in relational algebra
- UNION ALL, EXCEPT ALL, INTERSECT ALL
 - Bag semantics
 - Think of each row as having an implicit count (the number of times it appears in the table)
 - Bag union: sum up the counts from two tables
 - Bag difference: proper-subtract the two counts
 - Bag intersection: take the minimum of the two counts

Examples of bag operations

Bagl	Bag2	
fruit	fruit	
apple	apple	
apple	orange	
orange	orange	

```
(SELECT * FROM Bag1)
UNION ALL
(SELECT * FROM Bag2);

fruit
    apple
    apple
    orange
```

apple

orange

orange

```
(SELECT * FROM Bag1)

EXCEPT ALL

(SELECT * FROM Bag2);

fruit

apple
```

```
(SELECT * FROM Bag1)
INTERSECT ALL
(SELECT * FROM Bag2);

fruit
apple
```

orange

Examples of set versus bag operations

Poke (uid1, uid2, timestamp)

```
• (SELECT uid1 FROM Poke)

EXCEPT

(SELECT uid2 FROM Poke);
```

- Users who poked others but never got poked by others
- (SELECT uidl FROM Poke) EXCEPT ALL (SELECT uid2 FROM Poke);
 - Users who poked others more than others poke them

SQL features covered so far

- SELECT-FROM-WHERE statements (select-projectjoin queries)
- Set and bag operations

Next: how to nest SQL queries

Table subqueries

- Use query result as a table
 - In set and bag operations, FROM clauses, etc.
 - A way to "nest" queries
- Example: names of users who poked others more than others poked them

```
• SELECT DISTINCT name
FROM User,

((SELECT uidl AS uid FROM Poke))
EXCEPT ALL
(SELECT uid2 AS uid FROM Poke))
AS T
WHERE User.uid = T.uid;
```

Scalar subqueries

- A query that returns a single row can be used as a value in WHERE, SELECT, etc.
- Example: users at the same age as Bart

```
• SELECT *
FROM User
What's Bart's age?
WHERE age = (SELECT age
FROM User
WHERE name = 'Bart');
```

- Runtime error if subquery returns more than one row
 - Under what condition will this error never occur?
- What if the subquery returns no rows?
 - The answer is treated as a special value NULL, and the comparison with NULL will fail

IN subqueries

- *x* IN (*subquery*) checks if *x* is in the result of *subquery*
- Example: users at the same age as (some) Bart

```
• SELECT *
FROM User
WHERE age IN (SELECT age
FROM User
WHERE name = 'Bart');
```

EXISTS subqueries

- EXISTS (*subquery*) checks if the result of *subquery* is non-empty
- Example: users at the same age as (some) Bart
 - SELECT *
 FROM User AS u
 WHERE EXISTS (SELECT * FROM User
 WHERE name = 'Bart'
 AND age = u.age);
 - This happens to be a correlated subquery—a subquery that references tuple variables in surrounding queries

Semantics of subqueries

```
• SELECT *
FROM User AS u
WHERE EXISTS (SELECT * FROM User
WHERE name = 'Bart'
AND age = u.age);
```

- For each row u in User
 - Evaluate the subquery with the value of u.age
 - If the result of the subquery is not empty, output u.*
- The DBMS query optimizer may choose to process the query in an equivalent, but more efficient way (example?)

Scoping rule of subqueries

- To find out which table a column belongs to
 - Start with the immediately surrounding query
 - If not found, look in the one surrounding that; repeat if necessary
- Use table_name.column_name notation and AS (renaming) to avoid confusion

Another example

```
• SELECT * FROM User u
WHERE EXISTS
(SELECT * FROM Member m
WHERE uid = u.uid
AND EXISTS
(SELECT * FROM Member
WHERE uid = u.uid
AND gid <> m.gid));
```

Users who join at least two groups

Quantified subqueries

- A quantified subquery can be used syntactically as a value in a WHERE condition
- Universal quantification (for all):
 - ... WHERE x op ALL (subquery) ...
 - True iff for all t in the result of subquery, x op t
- Existential quantification (exists):
 - ... WHERE x op ANY (subquery) ...
 - True iff there exists some t in subquery result such that x op t
 - Beware
 - In common parlance, "any" and "all" seem to be synonyms
 - In SQL, ANY really means "some"

Examples of quantified subqueries

Which users are the most popular?

More ways to get the most popular

Which users are the most popular?

```
• SELECT *
FROM User AS u
WHERE NOT EXISTS
(SELECT * FROM User
WHERE pop > u.pop);
```

• SELECT * FROM User
WHERE uid NOT IN
 (SELECT ul.uid
 FROM User AS ul, User AS u2
WHERE ul.pop < u2.pop);</pre>

SQL features covered so far

- SELECT-FROM-WHERE statements
- Set and bag operations
- Subqueries
 - Subqueries allow queries to be written in more declarative ways (recall the "most popular" query)
 - But in many cases they don't add expressive power
 - Try translating other forms of subqueries into [NOT] EXISTS,
 which in turn can be translated into join (and difference)
 - Watch out for number of duplicates though

Next: aggregation and grouping

Aggregates

- Standard SQL aggregate functions: COUNT, SUM, AVG, MIN, MAX
- Example: number of users under 18, and their average popularity
 - SELECT COUNT(*), AVG(pop)
 FROM User
 WHERE age < 18;
 - COUNT(*) counts the number of rows

Aggregates with DISTINCT

Example: How many users are in some group?

```
• SELECT COUNT(DISTINCT uid) FROM Member; is equivalent to:
```

• SELECT COUNT(*)
FROM (SELECT DISTINCT uid FROM Member);

Grouping

• SELECT ... FROM ... WHERE ... GROUP BY list_of_columns;

- Example: compute average popularity for each age group
 - SELECT age, AVG(pop) FROM User GROUP BY age;

Semantics of GROUP BY

SELECT ... FROM ... WHERE ... GROUP BY ...;

- Compute FROM (×)
- Compute WHERE (σ)
- Compute GROUP BY: group rows according to the values of GROUP BY columns
- Compute SELECT for each group (π)
 - For aggregation functions with DISTINCT inputs, first eliminate duplicates within the group
- Number of groups = number of rows in the final output

Example of computing GROUP BY

SELECT age, AVG(pop) FROM User GROUP BY age;

uid	name	age	рор
142	Bart	10	0.9
857	Lisa	8	0.7
123	Milhouse	10	0.2
456	Ralph	8	0.3

Compute GROUP BY: group rows according to the values of GROUP BY columns

Compute SELECT for each group

age	avg_pop
10	0.55
8	0.50

uid	name	age	рор
142	Bart	10	0.9
123	Milhouse	10	0.2
857	Lisa	8	0.7
456	Ralph	8	0.3

Aggregates with no GROUP BY

 An aggregate query with no GROUP BY clause = all rows go into one group

SELECT AVG(pop) FROM User;

Group all rows into one group

Aggregate over the whole group

avg pop

uid	name	age	рор
142	Bart	10	0.9
857	Lisa	8	0.7
123	Milhouse	10	0.2
456	Ralph	8	0.3

uid	name	age	рор	
142	Bart	10	0.9	
857	Lisa	8	0.7	
123	Milhouse	10	0.2	
456	Ralph	8	0.3	

Restriction on SELECT

- If a query uses aggregation/group by, then every column referenced in SELECT must be either
 - Aggregated, or
 - A GROUP BY column
- This restriction ensures that any SELECT expression produces only one value for each group

Examples of invalid queries

- SELECT uid, age FROM User GROUP BY age;
 - Recall there is one output row per group
 - There can be multiple uid values per group
- SELECT und, MAX(pop) FROM User;
 - Recall there is only one group for an aggregate query with no GROUP BY clause
 - There can be multiple *uid* values
 - Wishful thinking (that the output *uid* value is the one associated with the highest popularity) does NOT work
 - Another way of writing the "most popular" query?

HAVING

- Used to filter groups based on the group properties (e.g., aggregate values, GROUP BY column values)
- SELECT ... FROM ... WHERE ... GROUP BY ... HAVING condition;
 - Compute FROM (×)
 - Compute WHERE (σ)
 - Compute GROUP BY: group rows according to the values of GROUP BY columns
 - Compute HAVING (another σ over the groups)
 - Compute SELECT (π) for each group that passes HAVING

HAVING examples

- List the average popularity for each age group with more than a hundred users
 - SELECT age, AVG(pop)
 FROM User
 GROUP BY age
 HAVING COUNT(*) > 100;
 - Can be written using WHERE and table subqueries
- Find average popularity for each age group over 10
 - SELECT age, AVG(pop)
 FROM User
 GROUP BY age
 HAVING age > 10;
 - Can be written using WHERE without table subqueries

SQL features covered so far

- SELECT-FROM-WHERE statements
- Set and bag operations
- Subqueries
- Aggregation and grouping
 - More expressive power than relational algebra

Next: ordering output rows

ORDER BY

- SELECT [DISTINCT] ...
 FROM ... WHERE ... GROUP BY ... HAVING ...
 ORDER BY output_column [ASC | DESC], ...;
- ASC = ascending, DESC = descending
- Semantics: After SELECT list has been computed and optional duplicate elimination has been carried out, sort the output according to ORDER BY specification

ORDER BY example

- List all users, sort them by popularity (descending) and name (ascending)
 - SELECT uid, name, age, pop FROM User ORDER BY pop DESC, name;
 - ASC is the default option
 - Strictly speaking, only output columns can appear in ORDER BY clause (although some DBMS support more)
 - Can use sequence numbers instead of names to refer to output columns: ORDER BY 4 DESC, 2;

SQL features covered so far

- SELECT-FROM-WHERE statements
- Set and bag operations
- Subqueries
- Aggregation and grouping
- Ordering

Next: NULL's, outerjoins, data modification, constraints, ...