## CPS 570: Artificial Intelligence

## Logic

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## Logic and AI

- Would like our Al to have knowledge about the world, and logically draw conclusions from it
- Search algorithms generate successors and evaluate them, but do not "understand" much about the setting
- Example question: is it possible for a chess player to have all her pawns and 2 queens?
- Search algorithm could search through tons of states to see if this ever happens, but...


## A story

- You roommate comes home; he/she is completely wet
- You know the following things:
- Your roommate is wet
- If your roommate is wet, it is because of rain, sprinklers, or both
- If your roommate is wet because of sprinklers, the sprinklers must be on
- If your roommate is wet because of rain, your roommate must not be carrying the umbrella
- The umbrella is not in the umbrella holder
- If the umbrella is not in the umbrella holder, either you must be carrying the umbrella, or your roommate must be carrying the umbrella
- You are not carrying the umbrella
- Can you conclude that the sprinklers are on?
- Can AI conclude that the sprinklers are on?


## Knowledge base for the story

- RoommateWet
- RoommateWet => (RoommateWetBecauseOfRain OR RoommateWetBecauseOfSprinklers)
- RoommateWetBecauseOfSprinklers => SprinklersOn
- RoommateWetBecauseOfRain => NOT(RoommateCarryingUmbrella)
- UmbrellaGone
- UmbrellaGone => (YouCarryingUmbrella OR RoommateCarryingUmbrella)
- NOT(YouCarryingUmbrella)


## Syntax

- What do well-formed sentences in the knowledge base look like?
- A BNF grammar:
- Symbol $\rightarrow$ P, Q, R, ..., RoommateWet, ...
- Sentence $\rightarrow$ True | False | Symbol| NOT(Sentence) | (Sentence AND Sentence)
| (Sentence OR Sentence) | (Sentence => Sentence)
- We will drop parentheses sometimes, but formally they really should always be there


# Semantics 

- A model specifies which of the proposition symbols are true and which are false
- Given a model, I should be able to tell you whether a sentence is true or false
- Truth table defines semantics of operators:

| a | b | NOT(a) | a AND b | a OR $b$ | $\mathrm{a}=>\mathrm{b}$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| false | false | true | false | false | true |
| false | true | true | false | true | true |
| true | false | false | false | true | false |
| true | true | false | true | true | true |

- Given a model, can compute truth of sentence recursively with these


## Caveats

- TwolsAnEvenNumber OR

ThreelsAnOddNumber
is true (not exclusive OR)

- TwolsAnOddNumber =>

ThreelsAnEvenNumber
is true (if the left side is false it's always true)
All of this is assuming those symbols are assigned their natural values...

## Tautologies

- A sentence is a tautology if it is true for any setting of its propositional symbols

| $P$ | $Q$ | PORQ | $\mathrm{NOT}(\mathrm{P})$ AND <br> $\mathrm{NOT}(\mathrm{Q})$ | $(\mathrm{P} \mathrm{OR} \mathrm{Q)}$ <br> OR (NOT(P) <br> AND <br> $\mathrm{NOT}(\mathrm{Q}))$ |
| :---: | :---: | :---: | :---: | :---: |
| false | false | false | true | true |
| false | true | true | false | true |
| true | false | true | false | true |
| true | true | true | false | true |

- ( P OR Q) OR (NOT(P) AND NOT(Q)) is a tautology


## Is this a tautology?

- $(P=>Q) O R(Q=>P)$


# Logical equivalences 

- Two sentences are logically equivalent if they have the same truth value for every setting of their propositional variables

| $P$ | Q | PORQ | $\operatorname{NOT}(\operatorname{NOT}(\mathrm{P})$ <br> $\operatorname{AND~NOT(Q))}$ |
| :---: | :---: | :---: | :---: |
| false | false | false | false |
| false | true | true | true |
| true | false | true | true |
| true | true | true | true |

- P OR Q and NOT(NOT(P) AND NOT(Q)) are logically equivalent
- Tautology = logically equivalent to True


## Famous logical equivalences

- $(\mathrm{a}$ OR b) $\equiv$ (b OR a) commutatitvity
- $(\mathrm{a}$ AND b) $\equiv$ (b AND a) commutatitvity
- ((a AND b) AND c) $\equiv(a$ AND (b AND c)) associativity
- ((a OR b) OR c) $\equiv(\mathrm{a}$ OR (b OR c)) associativity
- $\operatorname{NOT}(\operatorname{NOT}(\mathrm{a})) \equiv \mathrm{a}$ double-negation elimination
- $(\mathrm{a}=>\mathrm{b}) \equiv(\operatorname{NOT}(\mathrm{b})=>\operatorname{NOT}(\mathrm{a}))$ contraposition
- $(a=>b) \equiv(\operatorname{NOT}(a)$ OR b) implication elimination
- NOT(a AND b) $\equiv(\operatorname{NOT}(a)$ OR NOT(b)) De Morgan
- NOT(a OR b) $\equiv$ (NOT(a) AND NOT(b)) De Morgan
- $(\mathrm{a}$ AND $(\mathrm{b}$ OR c) $) \equiv((\mathrm{a}$ AND b) OR (a AND c)) distributitivity
- $(\mathrm{a}$ OR $(\mathrm{b}$ AND c) $) \equiv((\mathrm{a}$ OR b) AND $(\mathrm{a}$ OR c) ) distributitivity


## Inference

- We have a knowledge base of things that we know are true
- RoommateWetBecauseOfSprinklers
- RoommateWetBecauseOfSprinklers => SprinklersOn
- Can we conclude that SprinklersOn?
- We say SprinklersOn is entailed by the knowledge base if, for every setting of the propositional variables for which the knowledge base is true, SprinklersOn is also true

| RWBOS | SprinklersOn | Knowledge base |
| :---: | :---: | :---: |
| false | false | false |
| false | true | false |
| true | false | false |
| true | true | true |

- SprinklersOn is entailed!


## Simple algorithm for inference

- Want to find out if sentence a is entailed by knowledge base...
- For every possible setting of the propositional variables, - If knowledge base is true and a is false, return false
- Return true
- Not very efficient: $2^{\# p r o p o s i t i o n a l ~ v a r i a b l e s ~}$ settings
- Suppose we were careless in how we specified our knowledge base:
- PetOfRoommatelsABird => PetOfRoommateCanFly
- PetOfRoommatelsAPenguin => PetOfRoommatelsABird
- PetOfRoommatelsAPenguin => NOT(PetOfRoommateCanFly)
- PetOfRoommatelsAPenguin
- No setting of the propositional variables makes all of these true
- Therefore, technically, this knowledge base implies anything
- TheMoonlsMadeOfCheese


## Reasoning patterns

- Obtain new sentences directly from some other sentences in knowledge base according to reasoning patterns
- If we have sentences $a$ and $a=>b$, we can correctly conclude the new sentence b
- This is called modus ponens
- If we have a AND b, we can correctly conclude a
- All of the logical equivalences from before also give reasoning patterns


# Formal proof that the sprinklers are on 

1) RoommateWet
2) RoommateWet => (RoommateWetBecauseOfRain OR RoommateWetBecauseOfSprinklers)
3) RoommateWetBecauseOfSprinklers => SprinklersOn
4) RoommateWetBecauseOfRain $=>$ NOT(RoommateCarryingUmbrella)
5) UmbrellaGone
6) UmbrellaGone => (YouCarryingUmbrella OR RoommateCarryingUmbrella)
7) NOT (YouCarryingUmbrella)
8) YouCarryingUmbrella OR RoommateCarryingUmbrella (modus ponens on 5 and 6)
9) NOT(YouCarryingUmbrella) => RoommateCarryingUmbrella (equivalent to 8)
10) RoommateCarryingUmbrella (modus ponens on 7 and 9)
11) $\mathrm{NOT}(\mathrm{NOT}($ RoommateCarryingUmbrella) (equivalent to 10)
12) $\operatorname{NOT(NOT(RoommateCarryingUmbrella))~}=>$ NOT(RoommateWetBecauseOfRain) (equivalent to 4 by contraposition)
13) NOT(RoommateWetBecauseOfRain) (modus ponens on 11 and 12)
14) RoommateWetBecauseOfRain OR RoommateWetBecauseOfSprinklers (modus ponens on 1 and 2)
15) $\operatorname{NOT(RoommateWetBecauseOfRain)~=>~RoommateWetBecauseOfSprinklers~(equivalent~to~14)~}$
16) RoommateWetBecauseOfSprinklers (modus ponens on 13 and 15)
17) SprinklersOn (modus ponens on 16 and 3 )

## Reasoning about penguins

1) PetOfRoommatelsABird $=>$ PetOfRoommateCanFly
2) PetOfRoommatelsAPenguin => PetOfRoommatelsABird
3) PetOfRoommatelsAPenguin => NOT(PetOfRoommateCanFly)
4) PetOfRoommatelsAPenguin
5) PetOfRoommatelsABird (modus ponens on 4 and 2)
6) PetOfRoommateCanFly (modus ponens on 5 and 1)
7) NOT(PetOfRoommateCanFly) (modus ponens on 4 and 3)
8) $\mathrm{NOT}($ PetOfRoommateCanFly) $=>$ FALSE (equivalent to 6)
9) FALSE (modus ponens on 7 and 8)
10) FALSE => TheMoonIsMadeOfCheese (tautology)
11) TheMoonIsMadeOfCheese (modus ponens on 9 and 10)

## Systematic inference?

- General strategy: if we want to see if sentence a is entailed, add NOT(a) to the knowledge base and see if it becomes inconsistent (we can derive a contradiction)
- Any knowledge base can be written as a single formula in conjunctive normal form
RoommateWet => (RoommateWetBecauseOfRain OR RoommateWetBecauseOfSprinklers) becomes
(NOT(RoommateWet) OR RoommateWetBecauseOfRain OR RoommateWetBecauseOfSprinklers)
- Formula for modified knowledge base is satisfiable if and only if sentence a is not entailed


## Resolution

- Unit resolution: if we have
- $\mathrm{I}_{1}$ OR $\mathrm{I}_{2}$ OR $\ldots \mathrm{ORI}_{\mathrm{k}}$ and
- $\operatorname{NOT}\left(\mathrm{l}_{\mathrm{i}}\right)$
we can conclude
- $I_{1}$ OR $I_{2}$ OR $\ldots \mathrm{I}_{\mathrm{i}-1}$ OR $\mathrm{I}_{\mathrm{i}+1}$ OR $\ldots$ OR $\mathrm{I}_{\mathrm{k}}$
- Basically modus ponens


## Resolution...

- General resolution: if we have
- $I_{1}$ OR $I_{2}$ OR $\ldots$ OR $_{k}$ and
- $m_{1}$ OR $m_{2}$ OR $\ldots$ OR $m_{n}$ where for some $i, j, l_{i}=\operatorname{NOT}\left(m_{j}\right)$ we can conclude
- $I_{1}$ OR $I_{2}$ OR $\ldots \mathrm{I}_{\mathrm{i}-1}$ OR $\mathrm{I}_{\mathrm{i}+1}$ OR $\ldots$ OR $\mathrm{I}_{\mathrm{k}}$ OR $\mathrm{m}_{1}$ OR $\mathrm{m}_{2}$ OR $\ldots$ OR $_{\mathrm{j}-1}$ OR $\mathrm{m}_{\mathrm{j}+1}$ OR $\ldots$ OR $\mathrm{m}_{\mathrm{n}}$
- Same literal may appear multiple times; remove those


## Resolution algorithm

- Given formula in conjunctive normal form, repeat:
- Find two clauses with complementary literals,
- Apply resolution,
- Add resulting clause (if not already there)
- If the empty clause results, formula is not satisfiable - Must have been obtained from P and NOT(P)
- Otherwise, if we get stuck (and we will eventually), the formula is guaranteed to be satisfiable (proof in a couple of slides)


## Example

- Our knowledge base:
- 1) RoommateWetBecauseOfSprinklers
- 2) NOT(RoommateWetBecauseOfSprinklers) OR SprinklersOn
- Can we infer SprinklersOn? We add:
-3) NOT(SprinklersOn)
- From 2) and 3), get
-4) NOT(RoommateWetBecauseOfSprinklers)
- From 4) and 1), get empty clause


# If we get stuck, why is the formula satisfiable? 

- Consider the final set of clauses C
- Construct satisfying assignment as follows:
- Assign truth values to variables in order $\mathrm{x}_{1}, \mathrm{x}_{2}, \ldots, \mathrm{x}_{\mathrm{n}}$

If $x_{j}$ is the last chance to satisfy a clause (i.e., all the other variables in the clause came earlier and were set the wrong way), then set $\mathrm{x}_{\mathrm{j}}$ to satisfy it

- Otherwise, doesn't matter how it's set
- Suppose this fails (for the first time) at some point, i.e., $\mathrm{x}_{\mathrm{j}}$ must be set to true for one last-chance clause and false for another
- These two clauses would have resolved to something involving only up to $x_{j-1}$ (not to the empty clause, of course), which must be satisfied
- But then one of the two clauses must also be satisfied contradiction


## Special case: Horn clauses

- Horn clauses are implications with only positive literals
- $x_{1}$ AND $x_{2}$ AND $x_{4}=>x_{3}$ AND $x_{6}$
- TRUE => $x_{1}$
- Try to figure out whether some $x_{j}$ is entailed
- Simply follow the implications (modus ponens) as far as you can, see if you can reach $x_{j}$
- $x_{j}$ is entailed if and only if it can be reached (can set everything that is not reached to false)
- Can implement this more efficiently by maintaining, for each implication, a count of how many of the left-hand side variables have been reached

