

## Can We Do Better? Mining frequent patterns without candidate generation

s Database projection and compression
s Project the database based on its frequent patterns
s Compress a database into a compact, Frequent-Pattern tree (FP-tree)
s condensed, but complete for frequent pattern mining
s no candidate generation: test projected database only!
s $\mathbb{B D}^{\text {Divide-and-conquer } \square}$
s decompose both the task and DB according to the frequent patterns obtained so far

## Benefits of the FP-tree Structure

## s Completeness

s Preserves complete information for frequent pattern mining

- Compactness
s Reducing irrelevant info $\square$ infrequent items are gone
s Items in frequency descent order: the more frequently occurring, the more likely to be shared
s Never be larger than the original database (not count node-links and the count field)
s For Connect-4 DB, compression ratio could be over 100


## Is Apriori Efficient Enough? <br> $\square$ Performance Bottlenecks!!

s Basic Idea: Candidate generation-and-test
$s$ Use frequent ( $k \square 1$ )-itemsets to generate candidate frequent $k$-itemsets
s Use database scan and pattern matching to test (i.e., collect counts for the candidate itemsets)
s Bottleneck:
s Generation may lead to huge candidate sets
s $n$ frequent 1-itemset will generate $n(n-1) / 2$ candidate 2itemsets
$s$ To discover a frequent pattern of length 100, e.g., $\left\{a_{1}\right.$, $\left.a_{2}, \square, a_{100}\right\}$, we need to generate $2^{100} \approx 10^{30}$ candidates.
s Test involves multiple scans of the entire database s Needs ( $n+1$ ) scans, $n$ is the length of the longest pattern

## Construction of FP-tree from a Transaction Database

$\begin{array}{ll}- \text { TID } & \text { Items bought } \quad \text { (ordered) frequent items }\end{array}$
$100 \quad\{f, a, c, d, g, i, m, p\} \quad\{f, c, a, m, p\}$
$200 \quad\{a, b, c, f, l, m, o\} \quad\{f, c, a, b, m\}$
$\{f, c, a, b, m\}$
$300 \quad\{b, f, h, j, o, w\}$
$\{f, b\}$
$\{b, c, k, s, p\}$
$\{c, b, p\}$
$500 \quad\{a, f, c, e, l, p, m, n\}$
$\{f, c, a, m, p\}$
1 Scan DB once, find frequent 1-itemset (single item pattern)
2 Order frequent items in frequency descending order
3 Scan DB again, construct FP-tree


## Mining Frequent Patterns with FP-trees

Idea: Frequent pattern growth
s Recursively grow frequent patterns by pattern and database partition
Method
s For each frequent item, construct its conditional pattern-base, and then its conditional FP-tree
s Repeat the process on each newly created conditional FP-tree
s Until the resulting FP-tree is empty, or it contains only one path $\square$ single path will generate all the combinations of its sub-paths, each of which is a frequent pattern

## From FP-tree to Conditional Pattern-Base

Starting at the frequent item header table in the FP-tree
Traverse the FP-tree by following the link of each frequent item $p$
s Accumulate all of transformed prefix paths of item $p$ to form $p \square \mathrm{~s}$ conditional pattern base


## From Conditional Pattern-Bases to Conditional FP-trees

For each pattern-base
s Accumulate the count for each item in the base
s Construct the FP-tree for the frequent items of the pattern base


## A Special Case: Single FP-tree Path

Suppose a (conditional) FP-tree $T$ has a single path $P$
$s$ The complete set of frequent patterns of $T$ can be generated by enumeration of all the combinations of the sub-paths of $P$


## Transformed Prefix Paths

Derive the transformed prefix paths of item $p$
s For each item $p$ in the tree, collect $p \square \mathrm{~s}$ prefix path with count $=p \square \mathrm{~s}$ frequency



FP-Growth vs. Apriori: Scalability With the Number of Transactions


FP-Growth vs. Tree-Projection:
Scalability with the Support Threshold


## Why Is FP-Growth the Winner?

decompose both the mining task and DB according to the frequent patterns obtained so far
no redundant counting
leads to focused search of smaller databases no candidate generation, no candidate test compressed database: FP-tree structure no repeated scan of entire database

## I/O-Bound FP-Growth: Scaling FPGrowth by DB Projection

FP-tree cannot fit in memory? $\square$ DB projection
s first partition a database into a set of projected DBs
$s$ then construct and mine FP-tree for each projected DB
s Alternative methods
s Construction of a disk-resident FP-tree
s Materialization and incremental update of an FP-tree

Partition-Based Projection



