

# Indexing: Part IV

CPS 216  
Advanced Database Systems

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## Announcements (February 12)

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- ❖ Reading assignments
  - Query processing survey (due next Monday)
  - Variant indexes (due next Wednesday)
- ❖ Homework #2 assigned today
  - Due February 26 (in two weeks)
- ❖ Homework #1
  - Sample solution available next Tuesday
  - Grades will be posted on Blackboard
- ❖ Recitation session tomorrow (will announce by email too)
  - D240 1-2pm
- ❖ Midterm and course project proposal in 3 weeks
- ❖ Message board

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## Keyword search

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The screenshot shows a search engine interface with several search results. The search query entered is "database AND search". The results include:

- The Internet Movie Database (IMDb)...
- ... Search the Internet Movie Database. For more search options, please visit Search central...
- Images | Google...
- CPS 216: Advanced Database Systems (Fall 2001)
  - Course Information
  - Course Description / Time and Place / Books
  - Resources: Staff...
- Association for Computing Machinery
  - Founded in 1947, ACM is the world's educational and scientific computing society. Today, our members—...

At the bottom of the search results, there is a search bar containing the text "database AND search" and a "Search" button.

What are the documents containing both "database" and "search"?

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## Keywords × documents

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All keywords

All documents

Document 1    Document 2    Document 3    Document #

"a"	1	1	1	...	1
"cat"	1	1	0	...	0
"database"	0	0	1	...	0
"dog"	0	1	0	...	1
"search"	0	0	1	...	0
...	...	...	...	...	...

1 means keyword appears in the document  
0 means otherwise

- ❖ Inverted lists: store the matrix by rows
  - ❖ Signature files: store the matrix by columns
- With compression, of course!

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## Inverted lists

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- ❖ Store the matrix by rows
- ❖ For each keyword, store an inverted list
  - $\langle \text{keyword}, \text{doc-id-list} \rangle$
  - $\langle \text{"database"}, \{3, 7, 142, 857, \dots\} \rangle$
  - $\langle \text{"search"}, \{3, 9, 192, 512, \dots\} \rangle$
  - It helps to sort *doc-id-list* (why?)
- ❖ Vocabulary index on keywords
  - B<sup>+</sup>-tree or hash-based
- ❖ How large is an inverted list index?

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## Using inverted lists

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- ❖ Documents containing "database"
  - Use the vocabulary index to find the inverted list for "database"
  - Return documents in the inverted list
- ❖ Documents containing "database" AND "search"
  - Return documents in the intersection of the two inverted lists
- ❖ OR? NOT?

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## What are “all” the keywords?

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- ❖ All sequences of letters (up to a given length)?
  - ... that actually appear in documents!
- ❖ All words in English?
- ❖ Plus all phrases?
  - Alternative: approximate phrase search by proximity
- ❖ Minus all stop words
  - They appear in nearly every document; not useful in search
  - Example: a, of, the, it
- ❖ Combine words with common stems
  - They can be treated as the same for the purpose of search
  - Example: database, databases

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## Frequency and proximity

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- ❖ Frequency
  - $\langle \text{keyword}, \{ \langle \text{doc-id}, \text{number-of-occurrences} \rangle, \langle \text{doc-id}, \text{number-of-occurrences} \rangle, \dots \} \rangle$
- ❖ Proximity (and frequency)
  - $\langle \text{keyword}, \{ \langle \text{doc-id}, \langle \text{position-of-occurrence}_1, \text{position-of-occurrence}_2, \dots \rangle \rangle, \langle \text{doc-id}, \langle \text{position-of-occurrence}_1, \dots \rangle \rangle, \dots \} \rangle$
  - When doing AND, check for positions that are near

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## Signature files

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- ❖ Store the matrix by columns and compress them
- ❖ For each document, store a  $w$ -bit signature
- ❖ Each word is hashed into a  $w$ -bit value, with only  $s < w$  bits turned on
- ❖ Signature is computed by taking the bit-wise OR of the hash values of all words on the document

Does  $doc_3$  contain "database"?  
 $hash(\text{"database"}) = 0110$       $doc_1$  contains "database": 0110  
 $hash(\text{"dog"}) = 1100$       $doc_2$  contains "dog": 1100  
 $hash(\text{"cat"}) = 0010$       $doc_3$  contains "cat" and "dog": 1110

☞ Some false positives; no false negatives

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## Bit-sliced signature files

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### ❖ Motivation

- To check if a document contains a word, we only need to check the bits that are set in the word's hash value
- So why bother retrieving all  $w$  bits of the signature?

doc	signature
1	0000111100
2	0000111100
3	0001111100
4	0111011100
...	...
$n$	0000111100

↑ Slice 7 ... Slice 0

Bit-sliced signature files

Starting to look like an inverted list again!

### ❖ Instead of storing $n$ signature files, store $w$ bit slices

### ❖ Only check the slices that correspond to the set bits in the word's hash value

### ❖ Start from the sparse slices

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## Inverted lists versus signatures

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### ❖ Inverted lists are better for most purposes (*TODS*, 1998)

### ❖ Problems of signature files

### ❖ Saving grace of signature files

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## Ranking result pages

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### ❖ A single search may return many pages

- A user will not look at all result pages
- Complete result may be unnecessary
- ☞ Result pages need to be ranked

### ❖ Possible ranking criteria

- Based on content
  - Number of occurrences of the search terms
  - Similarity to the query text
- Based on link structure
  - Backlink count
  - PageRank
- And more...

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## Textual similarity

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- ❖ Vocabulary:  $\{w_1, \dots, w_n\}$
- ❖ IDF (Inverse Document Frequency):  $\{f_1, \dots, f_n\}$ 
  - $f_i = 1 /$  the number of times  $w_i$  appears on the Web
- ❖ Significance of words on page  $p$ :  $\{p_1 f_1, \dots, p_n f_n\}$ 
  - $p_i$  is the number of times  $w_i$  appears on  $p$
- ❖ Textual similarity between two pages  $p$  and  $q$  is defined to be  $\{p_1 f_1, \dots, p_n f_n\} \cdot \{q_1 f_1, \dots, q_n f_n\} = p_1 q_1 f_1^2 + \dots + p_n q_n f_n^2$ 
  - $q$  could be the query text

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## Why weight significance by IDF?

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## Problems with content-based ranking

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## Backlink

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- ❖ A page with more backlinks is ranked higher
- ❖ Intuition: Each backlink is a “vote” for the page’s importance

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## Google’s PageRank

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- ❖ Main idea: Pages pointed by high-ranking pages are ranked higher
  - Definition is recursive by design
  - Based on global link structure; hard to spam
- ❖ Naïve PageRank
  - $N(p)$ : number of outgoing links from page  $p$
  - $B(p)$ : set of pages that point to  $p$
  - $\text{PageRank}(p) = \sum_{q \in B(p)} (\text{PageRank}(q) / N(q))$
  - ☞ Each page  $p$  gets a boost of its importance from each page that points to  $p$
  - ☞ Each page  $q$  evenly distributes its importance to all pages that  $q$  points to

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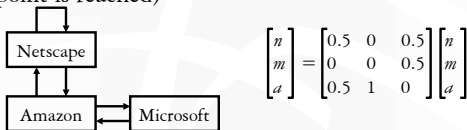
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## Calculating naïve PageRank

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- ❖ Initially, set all PageRank’s to 1; then evaluate  $\text{PageRank}(p) \leftarrow \sum_{q \in B(p)} (\text{PageRank}(q) / N(q))$  repeatedly until the values converge (i.e. a fixed point is reached)



$$\begin{bmatrix} n \\ m \\ a \end{bmatrix} = \begin{bmatrix} 0.5 & 0 & 0.5 \\ 0 & 0 & 0.5 \\ 0.5 & 1 & 0 \end{bmatrix} \begin{bmatrix} n \\ m \\ a \end{bmatrix}$$

$$\begin{bmatrix} n \\ m \\ a \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}, \begin{bmatrix} 1 \\ 0.5 \\ 1.5 \end{bmatrix}, \begin{bmatrix} 1.25 \\ 0.75 \\ 1 \end{bmatrix}, \begin{bmatrix} 1.125 \\ 0.5 \\ 1.375 \end{bmatrix}, \begin{bmatrix} 1.25 \\ 0.6875 \\ 1.0625 \end{bmatrix}, \dots, \begin{bmatrix} 1.2 \\ 0.6 \\ 1.2 \end{bmatrix}$$

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# Random surfer model

- ❖ A random surfer
  - Starts with a random page
  - Randomly selects a link on the page to visit next
  - Never uses the “back” button
- ❖ PageRank( $p$ ) measures the probability that a random surfer visits page  $p$

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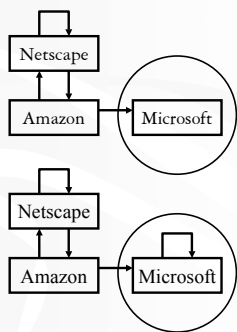
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# Problems with the naïve PageRank

- ❖ Dead end: a page with no outgoing links
  - A dead end causes all importance to “leak” eventually out of the Web
- ❖ Spider trap: a group of pages with no links out of the group
  - A spider trap will eventually accumulate all importance of the Web




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# Practical PageRank

- ❖  $d$ : decay factor
- ❖ PageRank( $p$ ) =  $d \cdot \sum_{q \in B(p)} (\text{PageRank}(q) / N(q)) + (1 - d)$
- ❖ Intuition in the random surfer model
  - A surfer occasionally gets bored and jump to a random page on the Web instead of following a random link on the current page

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## Google (1998)

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- ❖ Inverted lists in practice contain a lot of context information

	Hit: 2 bytes	Relative Capitalization	font size		
In URL/title/meta tag	plain: cap:1	imp:3	position: 12		within the page
	fancy: cap:1	imp = 7	type: 4	position: 8	within the page
In anchor text	anchor: cap:1	imp = 7	type: 4	hash:4   pos: 4	within the anchor

URL associated with the anchor

- ❖ PageRank is not the final ranking
  - Type-weight: depends on the type of the occurrence
    - For example, large font weights more than small font
  - Count-weight: depends on the number of occurrences
    - Increases linearly first but then tapers off
  - For multiple search terms, nearby occurrences are matched together and a proximity measure is computed
    - Closer proximity weights more

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## Suffix arrays (SODA, 1990)

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- ❖ Another index for searching text
- ❖ Conceptually, to construct a suffix array for string  $S$ 
  - Enumerate all  $|S|$  suffixes of  $S$
  - Sort these suffixes in lexicographical order
- ❖ To search for occurrences of a substring
  - Do a binary search on the suffix array

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## Suffix array example

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$S = \text{mississippi}$       $q = \text{sip}$

Suffixes:	Sorted suffixes:	Suffix array:	
mississippi	i	10	
ississippi	ippi	7	
ssissippi	issippi	4	No need to store
sissippi	ississippi	1	the suffix strings;
issippi	mississippi	0	just store where
ssippi	⇒ pi	9	they start
sippi	ppi	8	
ippi	⇒ sippi	6	$O( q  \cdot \log  S )$
ppi	⇒ sissippi	3	
pi	ssippi	5	
i	ssissippi	2	

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## One improvement

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- ❖ Remember how much of the query string has been matched

$q = \text{sisterhood}$

*low*:     $\Rightarrow$  sissippi...    Matched 3 characters

*middle*:  $\Rightarrow$  sisterhood...    Start checking from the 4<sup>th</sup> character

*high*:     $\Rightarrow$  sistering...    Matched 5 characters

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## Another improvement

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- ❖ Pre-compute the longest common prefix information between suffixes
  - For all (*low*, *middle*) and (*middle*, *high*) pairs that can come up in a binary search

$q = \text{sisterhood}$                        $O(|q| + \log |S|)$

*low*:     $\Rightarrow$  sissippi...    Matched 3 characters

*middle*:  $\Rightarrow$  sisterhood...    Start checking from the 7<sup>th</sup> character

...  $\rightarrow$  Matched 6 characters (pre-computed)

*high*:     $\Rightarrow$  sistering...    Matched 6 characters

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## Suffix arrays versus inverted lists

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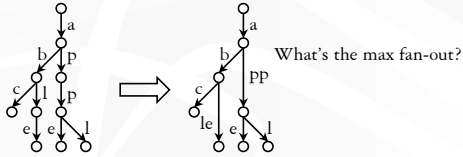
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### Trie: a string index

- ❖ A tree with edges labeled by characters
- ❖ A node represents the string obtained by concatenating all characters along the path from the root



- ❖ Compact trie: replace a path without branches by a single edge labeled by a string

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### Suffix tree

Index all suffixes of a large string in a compact trie

- ☞ Can support the same queries as a suffix array
- ❖ Internal nodes have fan-out  $\geq 2$  (except the root)
- ❖ No two edges out of the same node can share the same first character

To get linear space

- ❖ Instead of inlining the string labels, store pointers to them in the original string

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### Patricia trie, Pat tree, String B-tree

A Patricia trie is just like a compact trie, but

- ❖ Instead of labeling each edge by a string, only label by the first character and the string length
- ❖ Leaves point to strings
- ☞ Faster search (especially for external memory) because of inlining of the first character
- ☞ But

- ❖ A Pat tree indexes all suffixes of a large string in a Patricia trie
- ❖ A String B-tree uses a Patricia trie to store and compare strings in B-tree nodes

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## Summary

- ❖ General tree-based string indexing tricks
  - Trie, Patricia trie, String B-tree
  - Good exercise: put them in a GiST! ☺
- ❖ Two general ways to index for substring queries
  - Index words: inverted lists, signature files
  - Index all suffixes: suffix array, suffix tree, Pat tree
- ❖ Web search and information retrieval go beyond substring queries
  - IDF, PageRank, ...

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