Methods for Transforming Grammars

We will consider CFL without $\lambda$. It would be easy to add $\lambda$ to any grammar by adding a new start symbol $S_0$,

$$ S_0 \rightarrow S \mid \lambda $$
Theorem (Substitution) Let $G$ be a CFG. Suppose $G$ contains

$$A \rightarrow x_1Bx_2$$

where $A$ and $B$ are different variables, and $B$ has the productions

$$B \rightarrow y_1 \mid y_2 \mid \ldots \mid y_n$$

Then can construct $G'$ from $G$ by deleting

$$A \rightarrow x_1Bx_2$$

from $P$ and adding to it

$$A \rightarrow x_1y_1x_2 \mid x_1y_2x_2 \mid \ldots \mid x_1y_nx_2$$

Then, $L(G) = L(G')$. 
Example:

\[ S \rightarrow aBa \] becomes
\[ B \rightarrow aS \mid a \]

Definition: A production of the form
\[ A \rightarrow Ax, \ A \in V, \ x \in (V \cup T)^* \] is left recursive.
Example Previous expression grammar was left recursive.

\[
\begin{align*}
E & \rightarrow E + T \mid T \\
T & \rightarrow T \times F \mid F \\
F & \rightarrow I \mid (E) \\
I & \rightarrow a \mid b
\end{align*}
\]

Derivation of \( a + b + a + a \) is:

\[
\begin{align*}
E & \Rightarrow E + T \Rightarrow E + T + T \Rightarrow E + T + T + T \\
\Rightarrow^* & \ a + T + T + T 
\end{align*}
\]
Theorem (Removing Left recursion)
Let G=(V,T,S,P) be a CFG. Divide productions for variable A into left-recursive and non left-recursive productions:

\[
\begin{align*}
A & \to A x_1 \mid A x_2 \mid \ldots \mid A x_n \\
A & \to y_1 \mid y_2 \mid \ldots \mid y_m
\end{align*}
\]

where \(x_i, y_i\) are in \((V \cup T)^*\).

Then G’=(V\cup\{Z\}, T, S, P’) and P’ replaces rules of form above by

\[
\begin{align*}
A & \to y_i y_i Z, \text{ } i=1,2,\ldots,m \\
Z & \to x_i x_i Z, \text{ } i=1,2,\ldots,n
\end{align*}
\]
Example:

\[ E \rightarrow E + T | T \] becomes

\[ T \rightarrow T \ast F | F \] becomes

Now, Derivation of \( a + b + a + a \) is:
Useless productions

S \rightarrow aB \mid bA
A \rightarrow aA
B \rightarrow Sa
C \rightarrow cBc \mid a

What can you say about this grammar?

Theorem (useless productions) Let G be a CFG. Then \exists \ G’ that does not contain any useless variables or productions s.t. L(G)=L(G’).
To Remove Useless Productions:
Let $G = (V, T, S, P)$.

I. Compute $V_1 = \{\text{Variables that can derive strings of terminals}\}$

1. $V_1 = \emptyset$

2. Repeat until no more variables added
   • For every $A \in V$ with $A \rightarrow x_1 x_2 \ldots x_n$, $x_i \in (T^* \cup V_1)$, add $A$ to $V_1$

3. $P_1 = \text{all productions in } P \text{ with symbols in } (V_1 \cup T)^*$

Then $G_1 = (V_1, T, S, P_1)$ has no variables that can’t derive strings.
II. Draw Variable Dependency Graph

For $A \rightarrow xBy$, draw $A \rightarrow B$.

Remove productions for $V$ if there is no path from $S$ to $V$ in the dependency graph. Resulting Grammar $G'$ is s.t. $L(G) = L(G')$ and $G'$ has no useless productions.
Example:

\[
\begin{align*}
S & \rightarrow aB \mid bA \\
A & \rightarrow aA \\
B & \rightarrow Sa \mid b \\
C & \rightarrow cBc \mid a \\
D & \rightarrow bCb \\
E & \rightarrow Aa \mid b
\end{align*}
\]
Theorem (remove λ productions) Let G be a CFG with λ not in L(G). Then ∃ a CFG G’ having no λ-productions s.t. L(G)=L(G’).

To Remove λ-productions

1. Let $V_n = \{A \mid \exists \text{ production } A \rightarrow \lambda \}$
2. Repeat until no more additions
   • if $B \rightarrow A_1A_2\ldots A_m$ and $A_i \in V_n$ for all $i$, then put $B$ in $V_n$
3. Construct G’ with productions $P’$ s.t.
   • If $A \rightarrow x_1x_2\ldots x_m \in P$, $m \geq 1$, then put all productions formed when $x_j$ is replaced by λ (for all $x_j \in V_n$) s.t. $|\text{rhs}| \geq 1$ into $P’$. 
Example:

\[ S \rightarrow Ab \]
\[ A \rightarrow BCB \mid Aa \]
\[ B \rightarrow b \mid \lambda \]
\[ C \rightarrow cC \mid \lambda \]
Definition Unit Production

\[ A \rightarrow B \]

where \( A, B \in V \).

Consider removing unit productions:

Suppose we have

\[ A \rightarrow B \]
\[ B \rightarrow a \mid ab \]

But what if we have

\[ A \rightarrow B \] becomes
\[ B \rightarrow C \]
\[ C \rightarrow A \]
Theorem (Remove unit productions)
Let $G=(V,T,S,P)$ be a CFG without $\lambda$-productions. Then $\exists$ CFG $G’=(V’,T’,S,P’)$ that does not have any unit-productions and $L(G)=L(G’)$. 

To Remove Unit Productions:

1. Find for each $A$, all $B$ s.t. $A \Rightarrow^* B$
   (Draw a dependency graph)

2. Construct $G’=(V’,T’,S,P’)$ by
   (a) Put all non-unit productions in $P’$
   (b) For all $A \Rightarrow^* B$ s.t. $B \rightarrow y_1 | y_2 | \ldots | y_n \in P’$, put $A \rightarrow y_1 | y_2 | \ldots | y_n \in P’$
Example:

S → AB
A → B
B → C | Bb
C → A | c | Da
D → A
Theorem Let L be a CFL that does not contain $\lambda$. Then $\exists$ a CFG for L that does not have any useless productions, $\lambda$-productions, or unit-productions.

Proof

1. Remove $\lambda$-productions
2. Remove unit-productions
3. Remove useless productions

Note order is very important. Removing $\lambda$-productions can create unit-productions! QED.
Definition: A CFG is in Chomsky Normal Form (CNF) if all productions are of the form

\[ A \rightarrow BC \text{ or } A \rightarrow a \]

where \( A, B, C \in V \) and \( a \in T \).

Theorem: Any CFG \( G \) with \( \lambda \) not in \( L(G) \) has an equivalent grammar in CNF.

Proof:

1. Remove \( \lambda \)-productions, unit productions, and useless productions.

2. For every rhs of length \( \geq 1 \), replace each terminal \( x_i \) by a new variable \( C_j \) and add the production \( C_j \rightarrow x_i \).

3. Replace every rhs of length \( \geq 2 \) by a series of productions, each with rhs of length 2. QED.
Example:

\[ S \rightarrow CBcd \]
\[ B \rightarrow b \]
\[ C \rightarrow Cc \mid e \]
Definition: A CFG is in Greibach normal form (GNF) if all productions have the form

\[ A \rightarrow ax \]

where \( a \in T \) and \( x \in V^* \)

Theorem For every CFG \( G \) with \( \lambda \) not in \( L(G) \), \( \exists \) a grammar in GNF.

Proof:

1. Rewrite grammar in CNF.
2. Relabel Variables \( A_1, A_2, \ldots, A_n \)
3. Eliminate left recursion and use substitution to get all productions into the form:

\[ A_i \rightarrow A_j x_j, \ j > i \]
\[ Z_i \rightarrow A_j x_j, \ j \leq n \]
\[ A_i \rightarrow a x_i \]

where \( a \in T, \ x_i \in V^*, \) and \( Z_i \) are new variables introduced for left recursion.

4. All productions with \( A_n \) are in the correct form, \( A_n \rightarrow a x_n \). Use these productions as substitutions to get \( A_{n-1} \) productions in the correct form. Repeat with \( A_{n-2}, A_{n-3}, \) etc until all productions are in the correct form.