Compilers, Linkers, Loaders

Computer Science 104

Administrivia

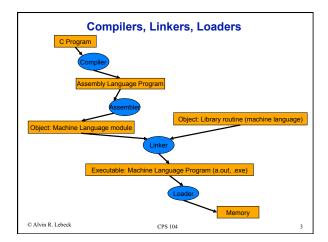
- Homework #2 Due Today
 - > Extra credit is on Blackboard
- Midterm: Monday Feb 16 in class
 - > Covers up through today's lecture
 - > Open book/notes
 - > Review session Friday
- · Friday: finish ASM programming on DE2 boards
- Reading
 - ➤ Compilers, linking & loading 2.12

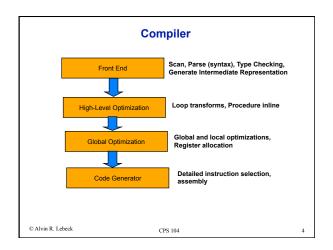
Next week

C Alvin R. Lebeck

- Logic Design
- Reading Appendix C.1-C.3, C.5

© Alvin R. Lebeck CPS 104





Compilers la \$t0, array li \$t1, 0 · Basic block = sequence of instructions with a single entry point and a single exit point lw \$t2, 0(\$t0) beq \$t2, \$r0, done First instruction is target of branch or jump or first instruction after branch or jump > Last instruction is branch or addi \$t2, \$t2, 10 sw \$t2, 0(\$t0) addi \$t1, \$t1, 4 Connect basic blocks to form control flow graph b loop Local optimizations are within a basic block Global optimizations are la \$t0, list lw \$t0, 0(list) Beq \$t0, \$r0, empty across basic block C Alvin R. Lebeck CPS 104

Common subexpression elimination > Array index address computation Strength reduction > Replace complex operations with simpler ops Constant Propagation > int x = 200; ... y = x + 40; z = y − 10; Copy Propagation Dead Store (code) elimination > Stores whose values are never used again Code that is never executed

CPS 104

Global Optimizations

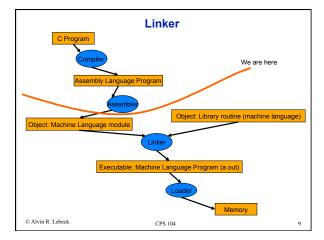
- · Occur across basic blocks
- · In addition to the previous set
- · Code Motion
 - > Find code that is loop invariant and move it before loop
 - > Computes the same value every iteration
- · Induction variable elimination
 - > Iterate over array using index k
 - > Could compute address using array_start plus offset (k * 4)
 - Use pointer based approach where you can just increment address by 4

© Alvin R. Lebeck CPS 104

Register Allocation

- Compiler first generates an intermediate representation (IR)
 - > Virtual registers (unlimited number of them...)
- · Must map from virtual registers to real set of registers
- · Goal is to reduce the number of loads & stores
- If not enough registers, must "spill" = save to stack and restore from stack
- · Sophisticated algorithms, reduces to graph coloring
 - > Given graph, color each node such that no two adjacent nodes have the same color
 - > Color = register number

© Alvin R. Lebeck CPS 104 8



Linker

- · Ability to resolve labels across multiple files
- · Compiler creates one object file per source file
- Includes symbol table that identifies labels within a file and any instructions that need to be "fixed"
- · Linker fills in values when they become known
- Static linking, all objects are linked to create executable file
- Dynamic linking, (DLL), occurs during execution
 Jump table

© Alvin R. Lebeck CPS 104 10

Loader: starting execution

- Part of Operating System that reads executable file off disk and starts execution
- Executable file has header information that identifies size of text and data in (ELF, COFF)
- · Works with OS to establish address space
 - > That ideal view of memory being 232 bytes large
- · Copies arguments into registers and stack
- Points PC to first instruction of startup code (which calls main)
 - > On return from main, this "startup" code executes an exit system call

© Alvin R. Lebeck CPS 104

Virtual Functions

· A little NiosII IDE and gnutools demo....

© Alvin R. Lebeck CPS 104

12

Summary

- Procedure calls
- Compilers, linkers, loaders

Next Time

Midterm Monday

Reading

Start Appendix C -- logic design

© Alvin R. Lebeck

CPS 104

13