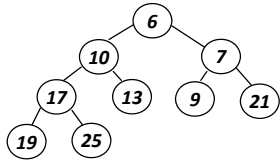


CompSci 100e

Program Design and Analysis II

| | | | | | | | | | | |
|---|---|----|---|----|----|---|----|----|----|----|
| | 6 | 10 | 7 | 17 | 13 | 9 | 21 | 19 | 25 | |
| 0 | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |



March 31, 2011

Prof. Rodger

CPS100E Spring11

1

Announcements

- Written linked lists/trees due today
- APT BSTCount due Tuesday, April 2
- Boggle assignment due in one week
 - Will discuss more in lab

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Abstract Data Type

- Stack (LIFO)
 - Push (add), pop (remove)
- Queue (FIFO)
 - Enqueue (add), dequeue (remove)
- Priority queue – queue, but best item dequeued (example: delete and return the minimum each time)
 - Enqueue (add), deleteMin (remove)

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Priority queue implementation

- Operations: add and delete min
- Want to do these operations efficiently

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Priority Queue sorting

code below sorts, complexity?

```
String[] array = {...}; // array filled with data
PriorityQueue<String> pq = new
    PriorityQueue<String>();
for(String s : array) pq.add(s);
for(int k=0; k < array.length; k++){
    array[k] = pq.remove();
}
```

Priority Queue top-M sorting

- What if we have *lots and lots and lots* of data
– code below sorts top-M elements, complexity?

```
Scanner s = ... // initialize;
PriorityQueue<String> pq =
    new PriorityQueue<String>();
while (s.hasNext()) {
    pq.add(s.next());
    if (pq.size() > M) pq.remove();
}
```

- What's advantageous about this code?
 - Store everything and sort everything?
 - Store everything, sort first M?
 - What is complexity of sort: $O(n \log n)$

PriorityQueue.java (Java 5+)

- What about objects inserted into pq?
 - Comparable, e.g., essentially sortable
 - How can we change what *minimal* means?
 - Implementation uses *heap*, tree stored in an array
- Use a Comparator for comparing entries we can make a min-heap act like a max-heap, see PQDemo
 - Where is class Comparator declaration? How used?
 - What if we didn't know about Collections.reverseOrder?
 - How do we make this ourselves?

Priority Queue implementation

- Heap data structure is fast and reasonably simple
 - Why not use inheritance hierarchy as was used with Map?
 - Trade-offs when using HashMap and TreeMap:
 - Time, space, ordering properties, TreeMap support?
- Changing comparison when calculating priority?
 - Create object to replace, or in lieu of compareTo
 - Comparable interface compares *this* to passed object
 - Comparator interface compares two passed objects
 - Both comparison methods: compareTo() and compare()
 - Compare two objects (parameters or self and parameter)
 - Returns -1, 0, +1 depending on <, ==, >

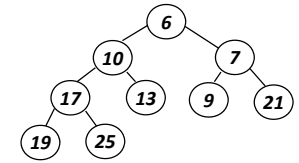
Creating Heaps

- Heap: array-based implementation of binary tree used for implementing priority queues:
 - add/insert, peek/getmin, remove/deletemin, $O(\log n)$
- Array minimizes storage (no explicit pointers), faster too, contiguous (cache) and indexing
- Heap has *shape* property and *heap/value* property
 - shape: tree filled at all levels (except perhaps last) and filled left-to-right (complete binary tree)
 - each node has value smaller than both children



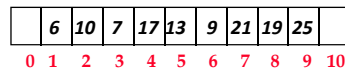
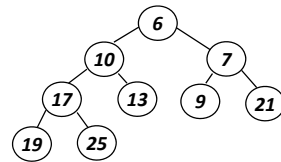
Array-based heap – one implementation for priority queue

- store “node values” in array beginning at index 1
- for node with index k
 - left child: index $2*k$
 - right child: index $2*k+1$
- why is this conducive for maintaining heap shape?
- what about heap property?
- is the heap a search tree?
- where is minimal node?
- where are nodes added? deleted?



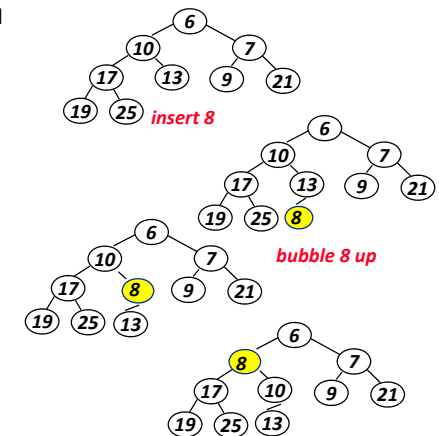
Thinking about heaps

- Where is minimal element?
- Where is maximal element?
- How many leaves are there in an N -node heap (big- O)?
- What is complexity of find max in a minheap? Why?
- Where is second smallest element? Why?

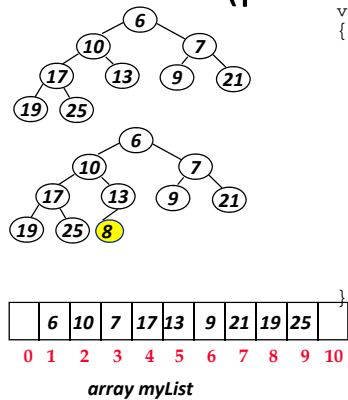


Adding values to heap

- to maintain heap shape, must add new value in left-to-right order of last level
 - could violate *heap property*
 - move value “up” if too small
- change places with parent if heap property violated
 - stop when parent is smaller
 - stop when root is reached
- pull parent down, swapping isn’t necessary (optimization)



Adding values, details (pseudocode)

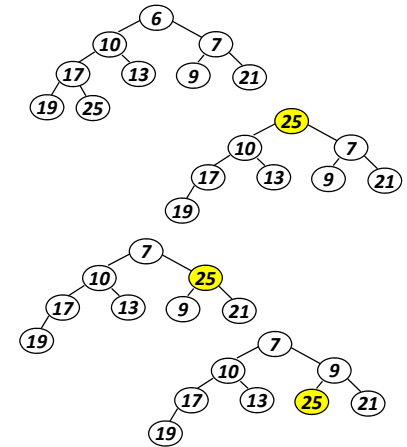


```
void add(Object elt)
{
    // add elt to heap in myList
    myList.add(elt);
    int loc = myList.size()-1;

    while (1 < loc &&
           elt < myList.get(loc/2)){
        myList.set(loc,myList.get(loc/2));
        loc = loc/2; // go to parent
    }
    // what's true here?
    myList.set(loc,elt);
}
```

Removing minimal element

- Where is minimal element?
 - If we remove it, what changes, shape/property?
- How can we maintain shape?
 - “last” element moves to root
 - What property is violated?
- After moving last element, subtrees of root are heaps, why?
 - Move root down (pull child up) does it matter where?
- When can we stop “re-heaping”?
 - Less than both children
 - Reach a leaf



Priority Queue implementations

- Priority queues: average and worst case

| | Insert average | Getmin (delete) | Insert worst | Getmin (delete) |
|---------------|-------------------|--------------------|-----------------|--------------------|
| Unsorted list | | | | |
| Sorted list | | | | |
| Search tree | | | | |
| Balanced tree | | | | |
| Heap | | | | |

- *Heap has $O(n)$ build heap from n elements*

Anita Borg 1949-2003

- “Dr. Anita Borg tenaciously envisioned and set about to change the world for women and for technology. ... she fought tirelessly for the development technology with positive social and human impact.”
- “Anita Borg sought to revolutionize the world and the way we think about technology and its impact on our lives.”
- http://www.youtube.com/watch?v=1yPxd5jqz_Q

