

Brief Comments on Game Theory

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CPS 170

What is Game Theory

- Very general mathematical framework to study situations where multiple agents interact, including:
 - Popular notions of games
 - Everything up to and including multistep, multiagent, simultaneous move, partial information games
 - Can even include negotiating, posturing and uncertainty about the players and game itself
- von Neumann and Morgenstern (1944) was a major launching point for modern game theory
- Nash: Existence of equilibria in [general sum](#) games

Covered Today

- 2 player, zero sum **simultaneous move** games
- Example: Rock, Paper, Scissors
- Linear programming solution







Linear Programs (max formulation)

$$\begin{aligned} &\text{maximize : } c^T x \\ &\text{subject to : } Ax \leq b \\ &\quad \quad \quad : x \geq 0 \end{aligned}$$

- Note: min formulation also possible
 - Min: $c^T x$
 - Subject to: $Ax \geq b$
- Some use equality as the canonical representation (introducing slack variables)
- LP tricks
 - Multiply by -1 to reverse inequalities
 - Can easily introduce equality constraints, or arbitrary domain constraints

Rock, Paper, Scissors Zero Sum Formulation

- In zero sum games, one player's loss is other's gain
- Payoff matrix:

	 <i>R</i>	 <i>P</i>	 <i>S</i>
 <i>R</i>	0	-1	1
 <i>P</i>	1	0	-1
 <i>S</i>	-1	1	0

- Minimax solution maximizes worst case outcome

Rock, Paper, Scissors Equations

- R, P, S = probability that we play rock, paper, or scissors respectively ($R + P + S = 1$)
- U is our expected utility
- Bounding our utility:
 - Opponent rock case: $U \leq P - S$
 - Opponent paper case: $U \leq S - R$
 - Opponent scissors case: $U \leq R - P$
- Want to maximize U subject to constraints
- Solution: $(1/3, 1/3, 1/3)$

Rock, Paper, Scissors LP Formulation

- Our variables are: $x = [U, R, P, S]^T$
- We want:
 - Maximize U
 - $U \leq P - S$
 - $U \leq S - R$
 - $U \leq R - P$
 - $R + P + S = 1$
- How do we make this fit:
 maximize : $c^T x$
 subject to : $Ax \leq b$
 $: x \geq 0$
 ?

Rock Paper Scissors LP Formulation

$$x = [U, R, P, S]^T$$

$$A = \begin{pmatrix} 1 & 0 & -1 & 1 \\ 1 & -1 & 0 & 1 \\ 1 & -1 & 1 & 0 \\ 0 & 1 & 1 & 1 \\ 0 & -1 & -1 & -1 \end{pmatrix}$$

$$b = [0, 0, 0, 1, -1]^T$$

$$c = [1, 0, 0, 0]^T$$

maximize : $c^T x$
 subject to : $Ax \leq b$
 $: x \geq 0$

Rock, Paper, Scissors Solution

- If we feed this LP to an LP solver we get:
 - $R=P=S=1/3$
 - $U=0$
- Solution for the other player is:
 - The same...
 - By symmetry
- This is the minimax solution
- This is also an equilibrium
 - No player has an incentive to deviate
 - (Defined more precisely later in the course)

Tangent: Why is RPS Fun?

- OK, it's not...
- Why *might* RPS be fun?
 - Try to exploit non-randomness in your friends
 - Try to be random yourself

Minimax Solutions in General

- Minimax solutions for 2-player zero-sum games can always be found by solving a linear program
- The minimax solutions will also be equilibria
- For general sum games:
 - Minimax does not apply
 - Equilibria may not be unique
 - Need to search for equilibria using more computationally intensive methods