

Relational Database Design using E/R

Introduction to Databases

CompSci 316 Spring 2019



DUKE
COMPUTER SCIENCE

Announcements (Thu. Jan. 17)

- Sign up for Piazza and Gradiance
- Set up your VM
- Homework 1 – Problems 1 and 2 posted (after class!)
 - Due in 2.5 weeks
 - Problem 1 is already posted on gradiance (attempt as many times as you need!)
- Office hours have been posted on website
 - First TA office hour tomorrow

Relational model: review

- A database is a collection of **relations** (or **tables**)
- Each relation has a set of **attributes** (or **columns**)
- Each attribute has a name and a **domain** (or **type**)
- Each relation contains a set of **tuples** (or **rows**)

Keys

→ minimal & unique

- A set of attributes K is a **key** for a relation R if
 - In no instance of R will two different tuples agree on all attributes of K
 - That is, K can serve as a “**tuple identifier**”
 - No proper subset of K satisfies the above condition
 - That is, K is **minimal**
- Example: $User(uid, name, age, pop)$
 - uid is a key of $User$
 - age is not a key (not an identifier)
 - $\{uid, name\}$ is not a key (not minimal)

Superkey

$User$
 $uid, name, pop$
 uid
 $(uid, name)$

Schema vs. instance

<i>uid</i>	<i>name</i>	<i>age</i>	<i>pop</i>
142	Bart	10	0.9
123	Milhouse	10	0.2
857	Lisa	8	0.7
456	Ralph	8	0.3

- Is *name* a key of *User*?
 - Yes? Seems reasonable for this instance
 - No! User names are not unique **in general**
- Key declarations are part of the schema

More examples of keys

- `Member (uid, gid)`

- `{uid, gid}`

- ☞ A key can contain multiple attributes

- `Address (street_address, city, state, zip)`

- `{street_address, city, state}`

- `{street_address, zip}`

- ☞ A relation can have multiple keys!

- We typically pick one as the “primary” key, and underline all its attributes, e.g., `Address (street_address, city, state, zip)`

① key
② superkey
③ primary key

④ key = candidate key

Use of keys

- More constraints on data, fewer mistakes
- Look up a row by its key value
 - Many selection conditions are “key = value”
- “Pointers” to other rows (often across tables)
 - Example: *Member (uid, gid)*
 - *uid* is a key of *User*
 - *gid* is a key of *Group*
 - A *Member* row “links” a *User* row with a *Group* row
 - Many join conditions are “key = key value stored in another table”

Database design

- Understand the real-world domain being modeled
- Specify it using a database **design model**
 - More intuitive and convenient for schema design
 - But not necessarily implemented by DBMS
 - A few popular ones:
 - Entity/Relationship (E/R) model
 - Object Definition Language (ODL)
 - UML (Unified Modeling Language)
- Translate specification to the data model of DBMS
 - Relational, XML, object-oriented, etc.
- Create DBMS schema

But what about ORM?

- Automatic **object-relational mappers** are made popular by rapid Web development frameworks
 - For example, with Python SQLAlchemy:
 - You declare Python classes and their relationships
 - It automatically converts them into database tables
 - If you want, you can just work with Python objects, and never need to be aware of the database schema or write SQL
- But you still need designer discretion in all but simple cases
- Each language/library has its own syntax for creating schema and for querying/modifying data
 - Quirks and limitations cause portability problems
 - They are not necessarily easier to learn than SQL

Entity-relationship (E/R) model

- Historically and still very popular
- Concepts applicable to other design models as well
- Can think of as a “watered-down” object-oriented design model
- Primarily a design model—not directly implemented by DBMS
- Designs represented by E/R diagrams
 - We use the style of E/R diagram covered by the GMUW book; there are other styles/extensions
 - Very similar to UML diagrams

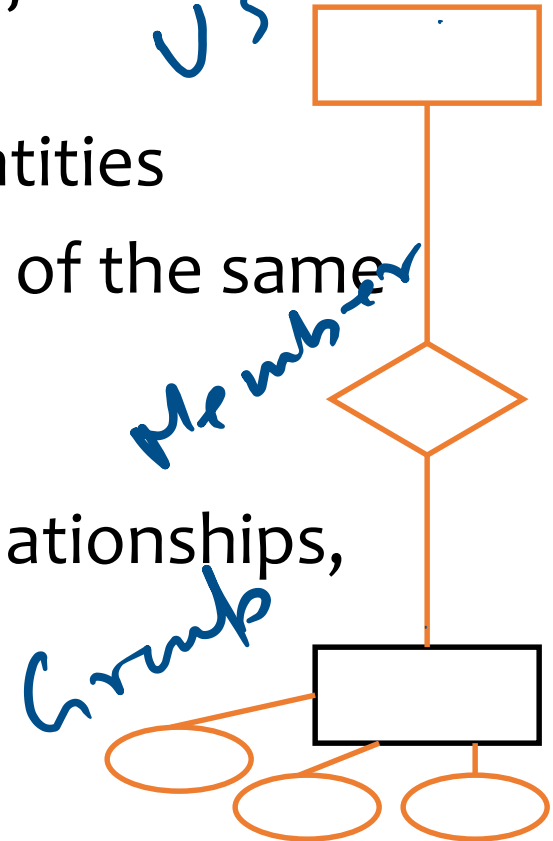
E/R basics

Rel data model.

relationship

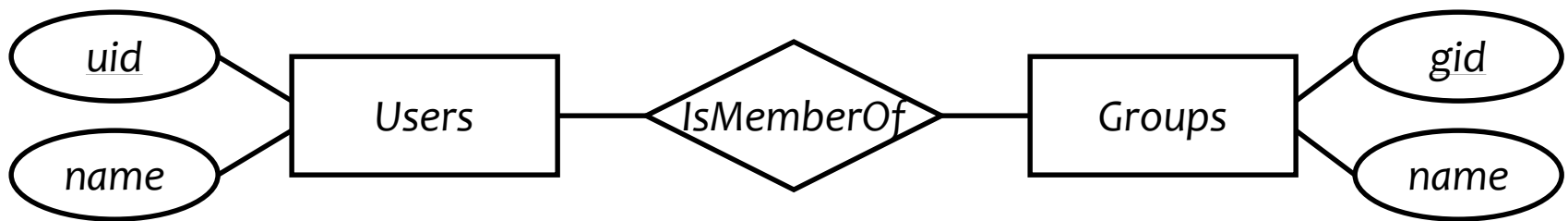
relationship set

- **Entity**: a “thing,” like an object
- **Entity set**: a collection of things of the same type, like a relation of tuples or a class of objects
 - Represented as a rectangle
- **Relationship**: an association among entities
- **Relationship set**: a set of relationships of the same type (among same entity sets)
 - Represented as a diamond
- **Attributes**: properties of entities or relationships, like attributes of tuples or objects
 - Represented as ovals



An example E/R diagram

- Users are members of groups

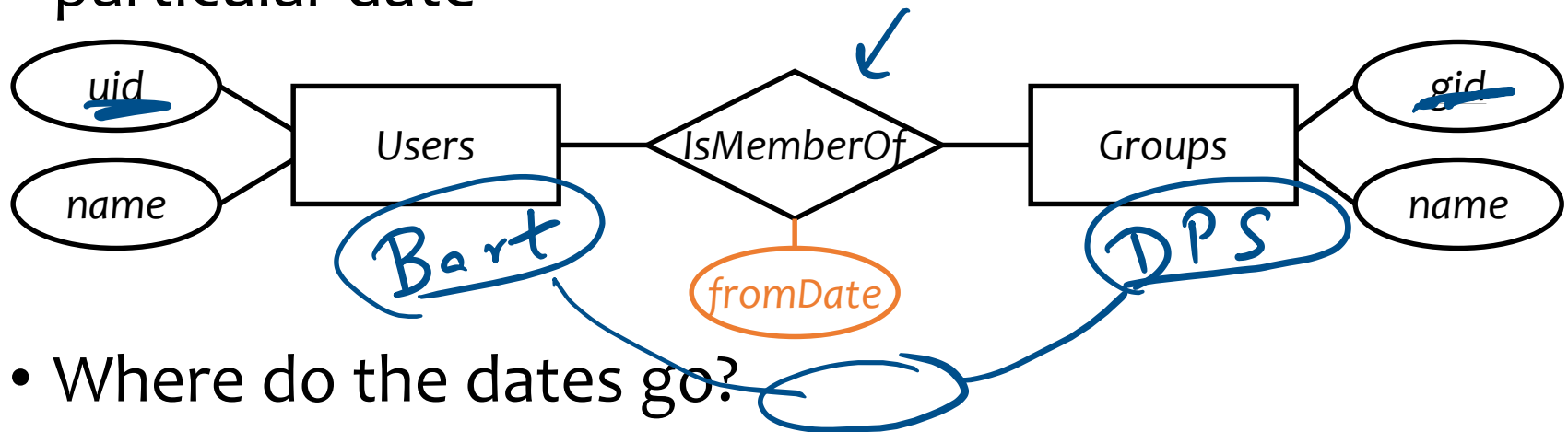


- A key of an entity set is represented by underlining all attributes in the key
 - A key is a set of attributes whose values can belong to at most one entity in an entity set—like a key of a relation

Attributes of relationships

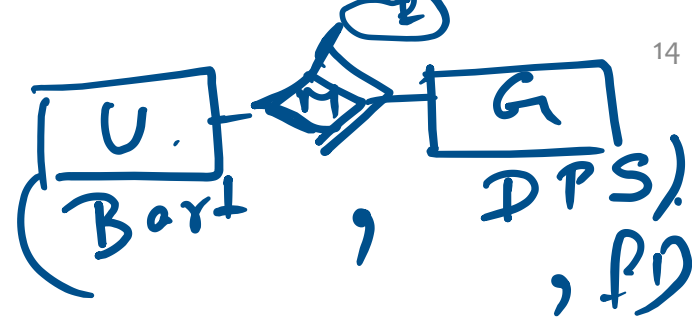
(uid, gid, date)

- Example: a user belongs to a group since a particular date



- Where do the dates go?
 - With *Users*?
 - But a user can join multiple groups on different dates
 - With *Groups*?
 - But different users can join the same group on different dates
 - With *IsMemberOf*!

More on relationships

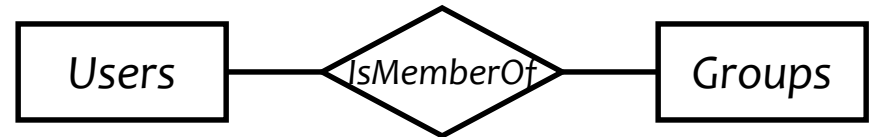


- There could be multiple relationship sets between the same entity sets
 - Example: *Users IsMemberOf Groups*; *Users Likes Groups*
- In a relationship set, each relationship is uniquely identified by the entities it connects
 - Example: Between Bart and “Dead Putting Society”, there can be at most one *IsMemberOf* relationship and at most one *Likes* relationship
- ☞ What if Bart joins DPS, leaves, and rejoins? How can we modify the design to capture historical membership information?
 - ☞ Make an entity set of *MembershipRecords*

Multiplicity of relationships

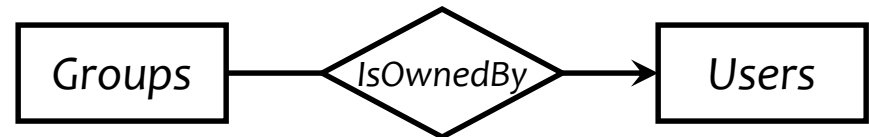
- E and F : entity sets
- **Many-many**: Each entity in E is related to 0 or more entities in F and vice versa

- Example:



- **Many-one**: Each entity in E is related to 0 or 1 entity in F , but each entity in F is related to 0 or more in E

- Example:



- **One-one**: Each entity in E is related to 0 or 1 entity in F and vice versa

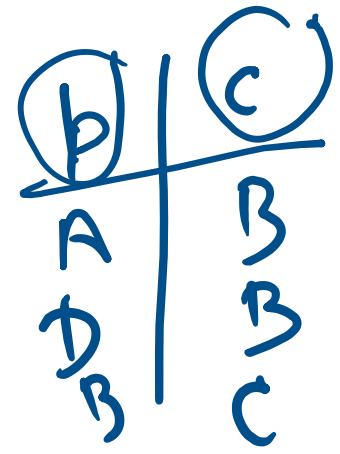
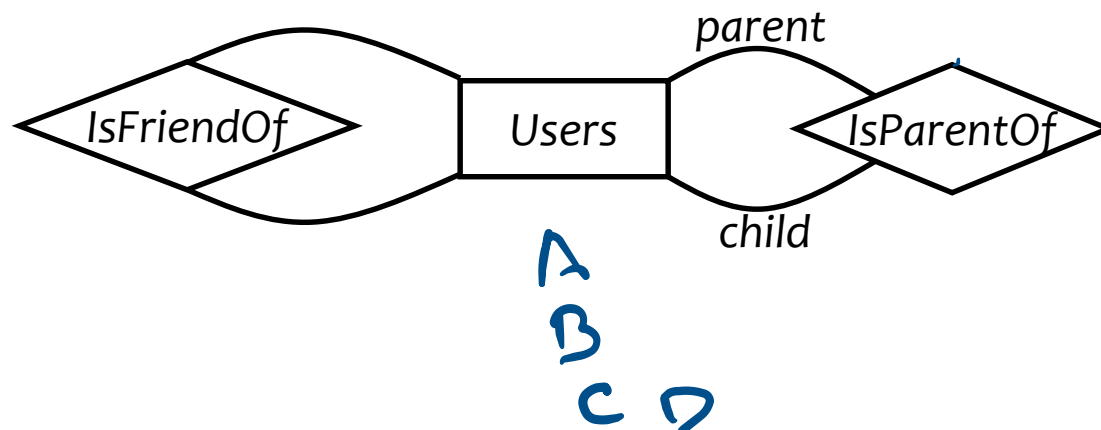
- Example:



- “One” (0 or 1) is represented by an arrow \longrightarrow
- “Exactly one” is represented by a rounded arrow \longrightarrow

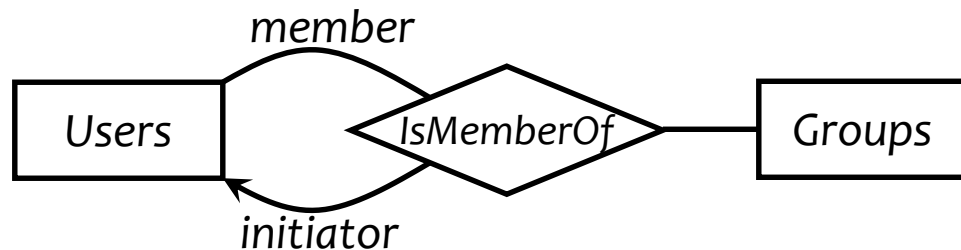
Roles in relationships

- An entity set may participate more than once in a relationship set
- ☞ May need to label edges to distinguish **roles**
- Examples
 - Users may be parents of others; label needed
 - Users may be friends of each other; label not needed



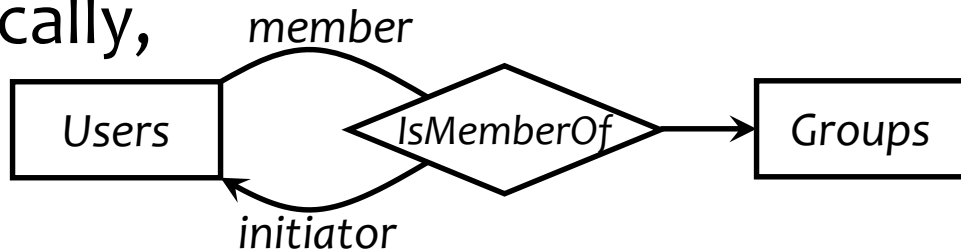
n -ary relationships

- Example: a user must have an initiator in order to join a group



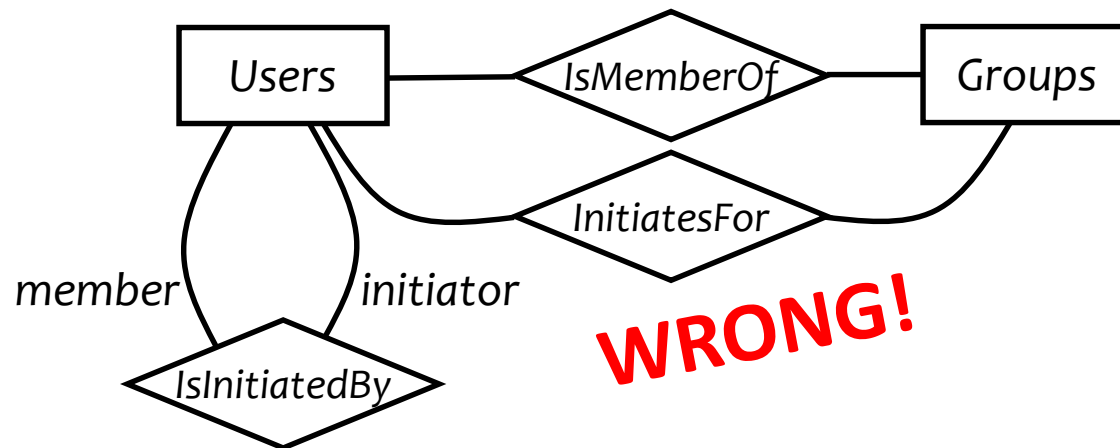
Rule for interpreting an arrow into entity set E in an n -ary relationship:

- Pick one entity from each of the other entity sets; together they can be related to at most one entity in E
- Exercise: hypothetically, what do these arrows imply?



n -ary versus binary relationships

- Can we model n -ary relationships using just binary relationships?

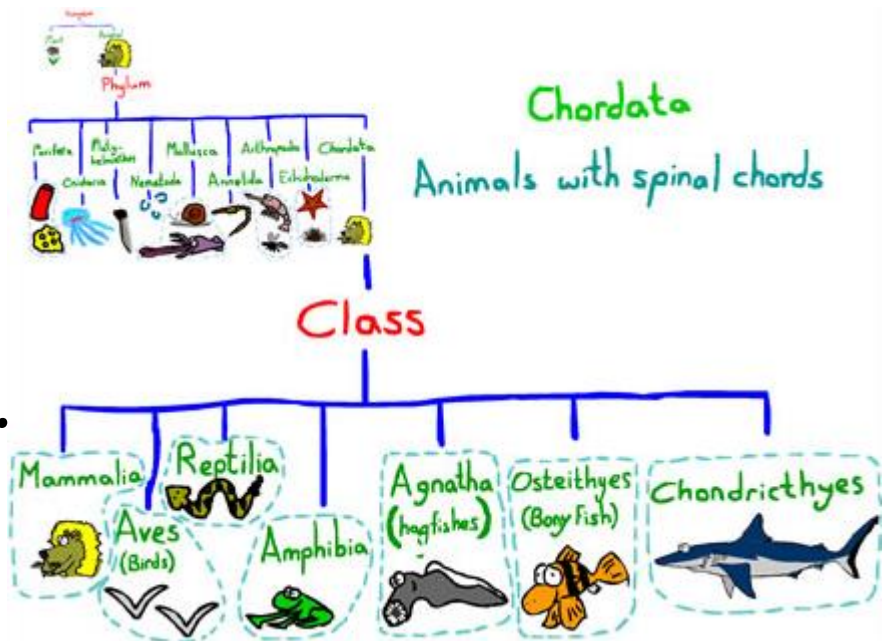


- No; for example:
 - Ralph is in both abc and gov
 - Lisa has served as initiator in both abc and gov
 - Ralph was initiated by Lisa in abc, but not by her in gov

Next: two special relationships



... is part of/belongs to ...

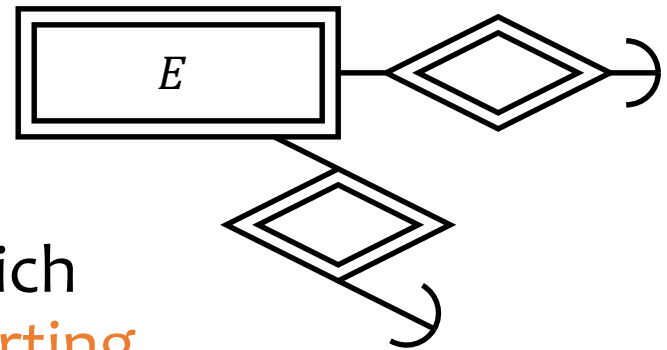


... is a kind of ...

Weak entity sets

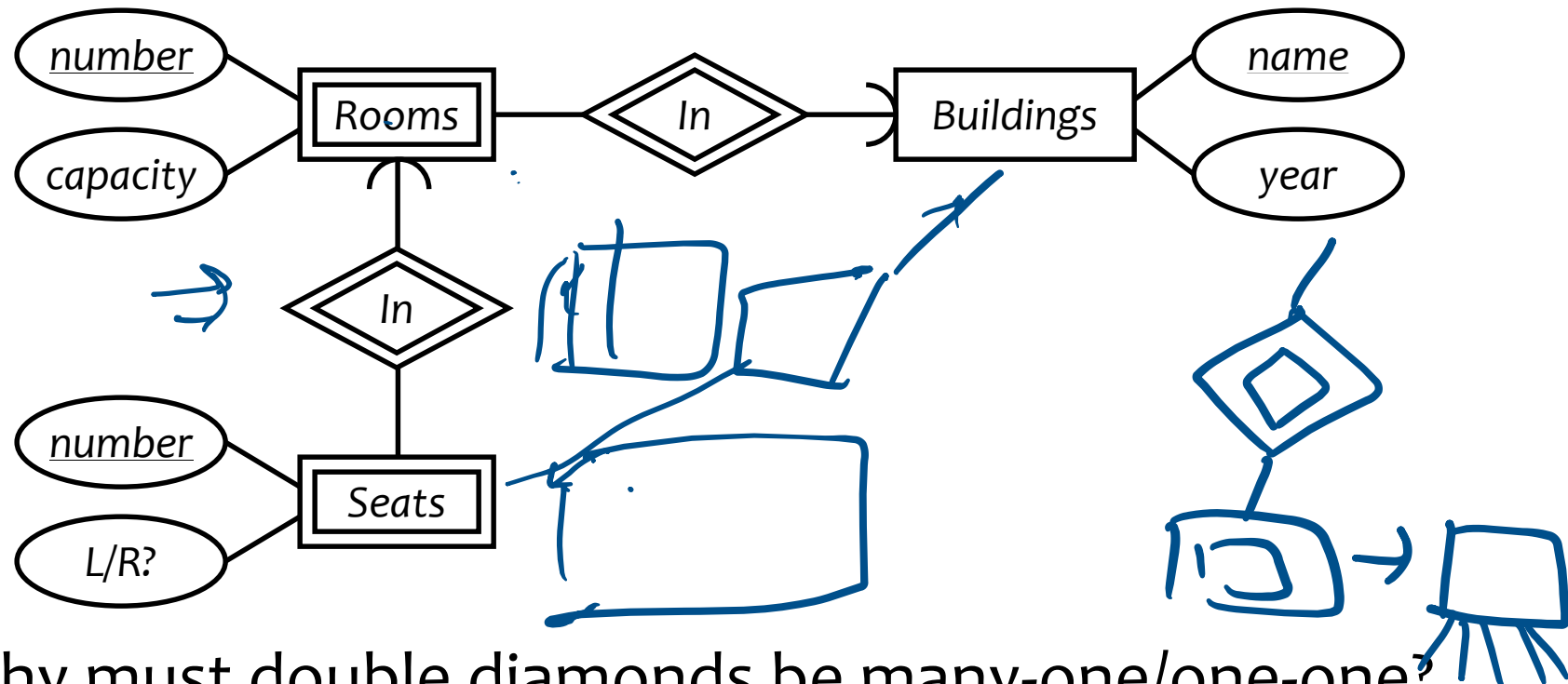
Sometimes, an entity's identity depends on some others'

- The key of a **weak entity set** E comes not completely from its own attributes, but from the keys of one or more other entity sets
 - E must link to them via many-one or one-one relationship sets
- Example: *Rooms* inside *Buildings* are partly identified by *Buildings*' name
- A weak entity set is drawn as a double rectangle
- The relationship sets through which it obtains its key are called **supporting relationship sets**, drawn as double diamonds



Weak entity set examples

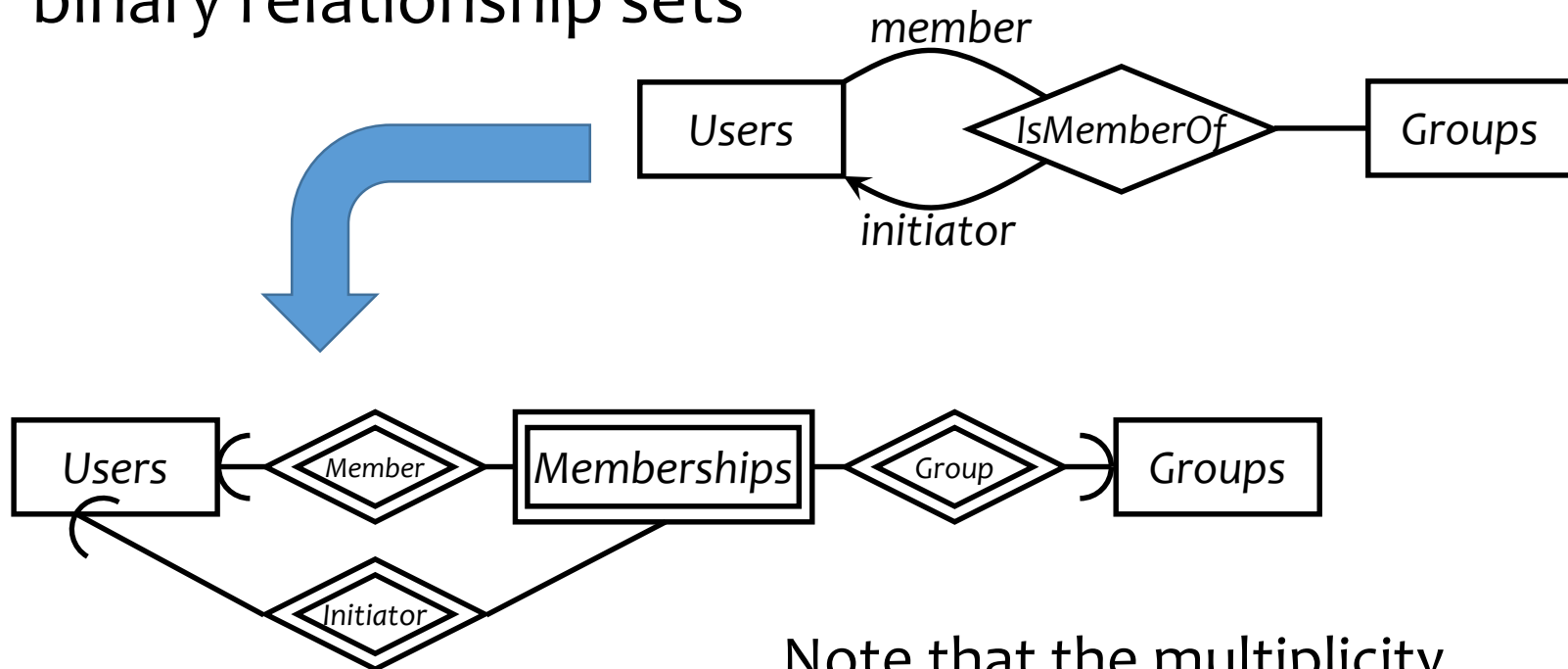
- Seats in rooms in building



- Why must double diamonds be many-one/one-one?
 - With many-many, we would not know which entity provides the key value!

Remodeling n -ary relationships

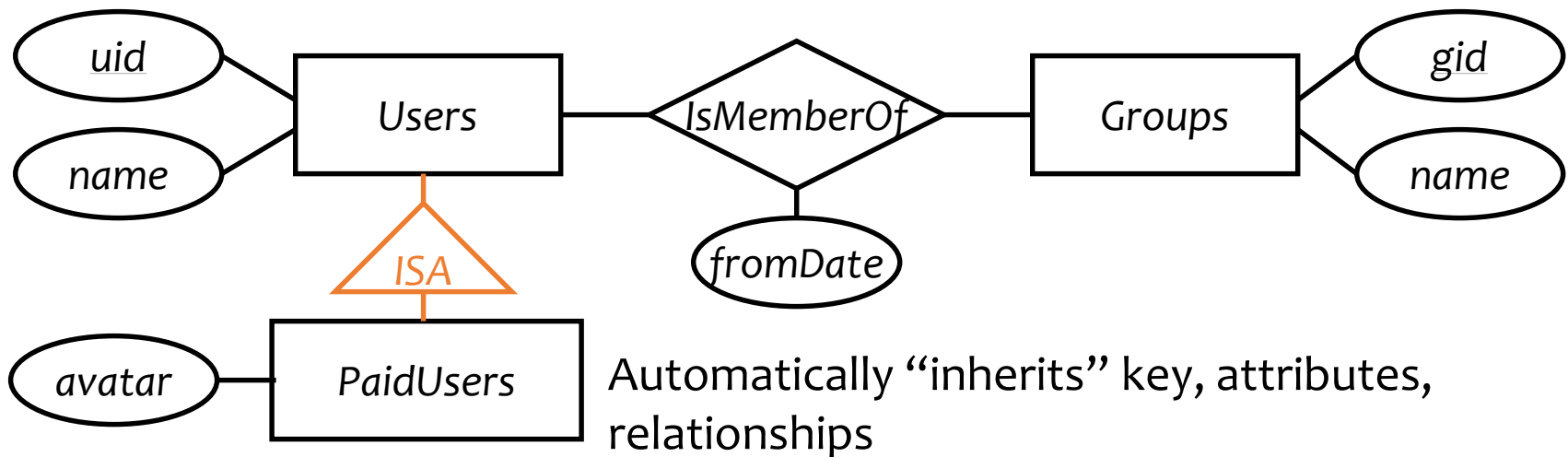
- An n -ary relationship set can be replaced by a weak entity set (called a **connecting entity set**) and n binary relationship sets



Note that the multiplicity constraint for *IsMemberOf* is lost

ISA relationships

- Similar to the idea of subclasses in object-oriented programming: subclass = special case, fewer entities, and possibly more properties
 - Represented as a triangle (direction is important)
- Example: paid users are users, but they also get avatars (yay!)



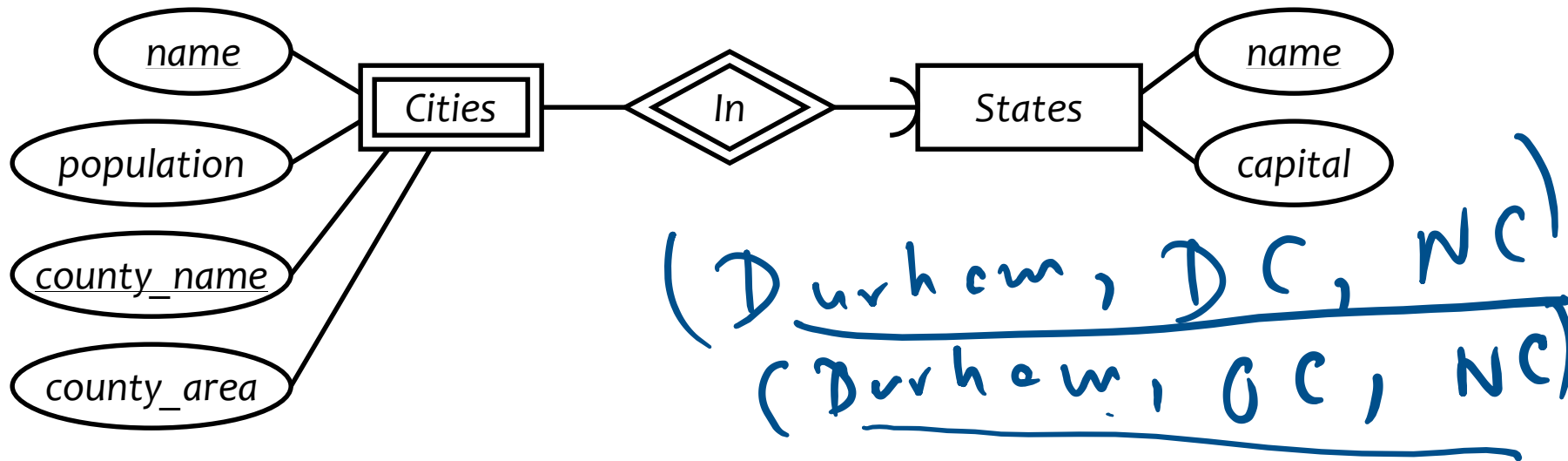
Summary of E/R concepts

- Entity sets
 - Keys
 - Weak entity sets
- Relationship sets
 - Attributes of relationships
 - Multiplicity
 - Roles
 - Binary versus n -ary relationships
 - Modeling n -ary relationships with weak entity sets and binary relationships
 - ISA relationships

Case study 1

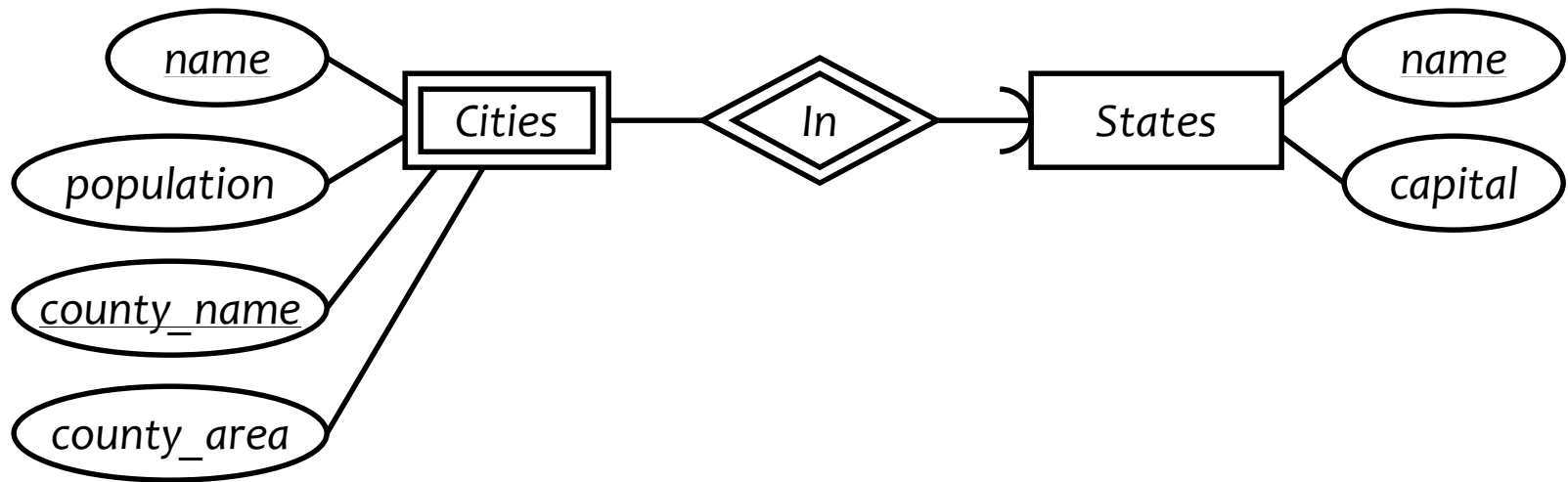
- Design a database representing cities, counties, and states
 - For states, record name and capital (city)
 - For counties, record name, area, and location (state)
 - For cities, record name, population, and location (county and state)
- Assume the following:
 - Names of states are unique
 - Names of counties are only unique within a state
 - Names of cities are only unique within a county
 - A city is always located in a single county
 - A county is always located in a single state

Case study 1: first design



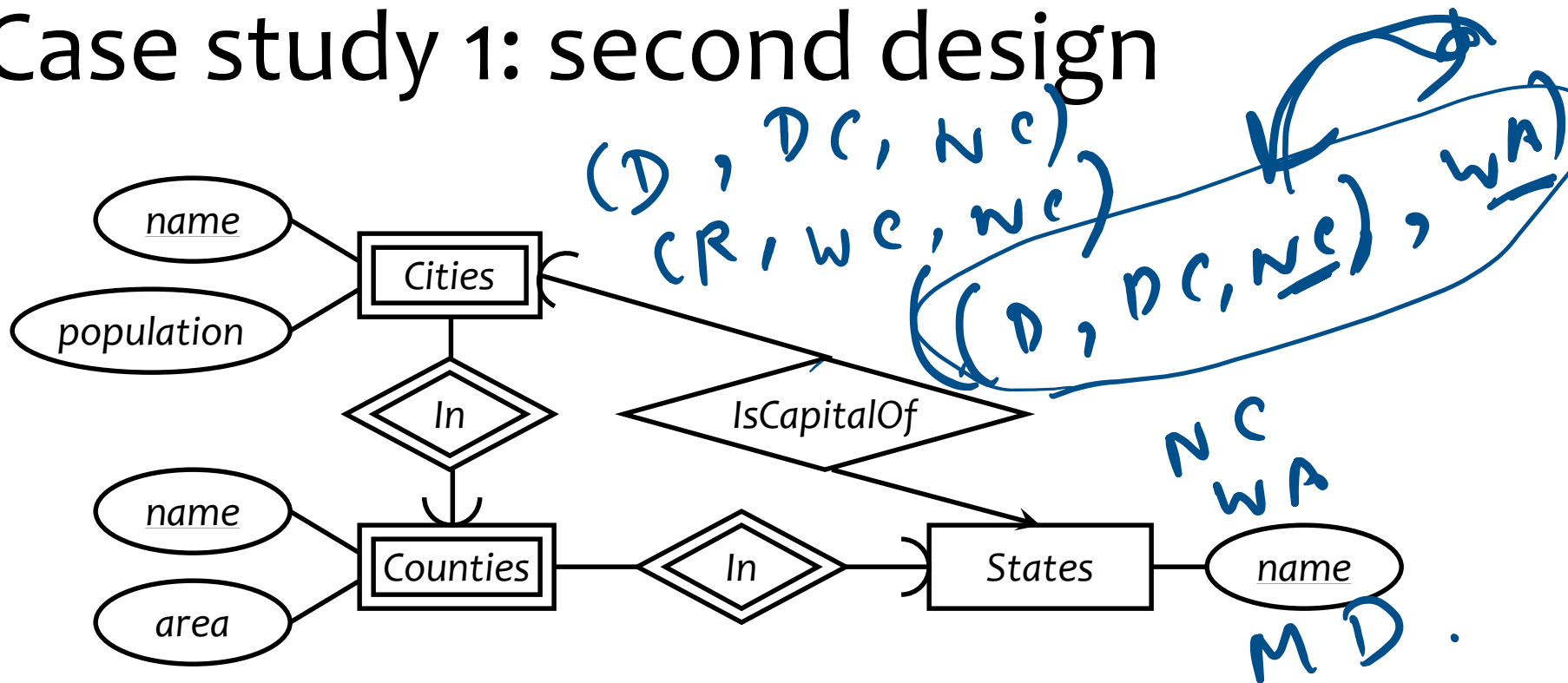
- Design a database representing cities, counties, and states
 - For states, record name and capital (city)
 - For counties, record name, area, and location (state)
 - For cities, record name, population, and location (county and state)
- Assume the following:
 - Names of states are unique
 - Names of counties are only unique within a state
 - Names of cities are only unique within a county
 - A city is always located in a single county
 - A county is always located in a single state

Case study 1: first design



- County area information is repeated for every city in the county
 - ☞ Redundancy is bad (why?)
- State capital should really be a city
 - ☞ Should “reference” entities through explicit relationships

Case study 1: second design

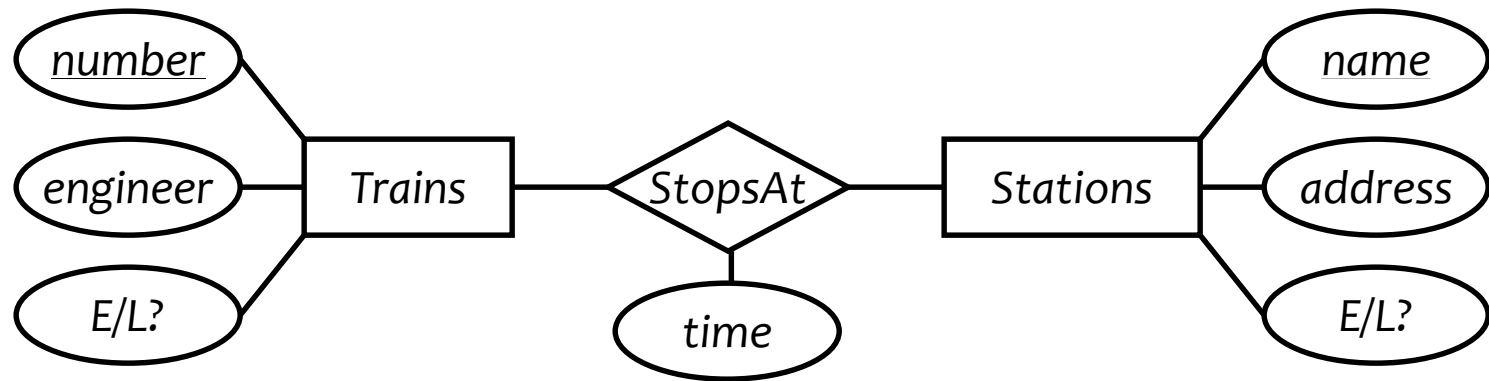


- Technically, nothing in this design prevents a city in state X from being the capital of another state Y , but...

Case study 2

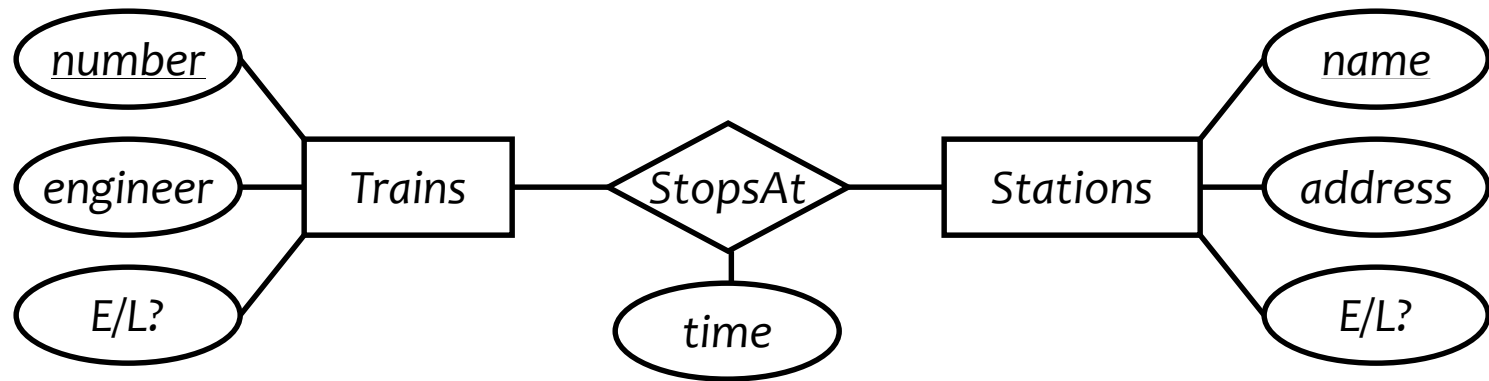
- Design a database consistent with the following:
 - A station has a unique name and an address, and is either an express station or a local station
 - A train has a unique number and an engineer, and is either an express train or a local train
 - A local train can stop at any station
 - An express train only stops at express stations
 - A train can stop at a station for any number of times during a day
 - Train schedules are the same everyday

Case study 2: first design



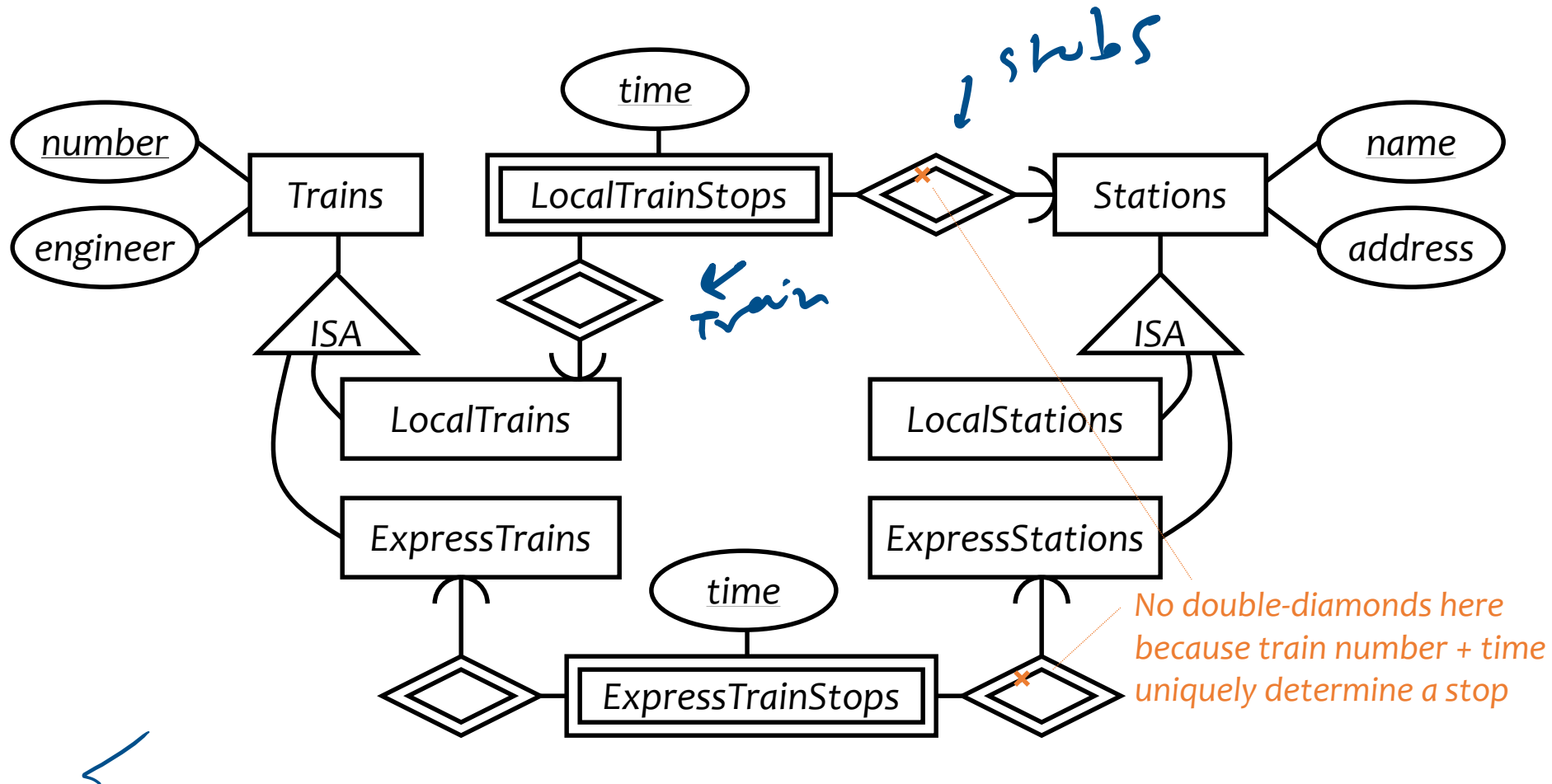
- Design a database consistent with the following:
 - A station has a unique name and an address, and is either an express station or a local station
 - A train has a unique number and an engineer, and is either an express train or a local train
 - A local train can stop at any station
 - An express train only stops at express stations
 - A train can stop at a station for any number of times during a day
 - Train schedules are the same everyday

Case study 2: first design



- Nothing in this design prevents express trains from stopping at local stations
 - ☞ We should capture as many constraints as possible
- A train can stop at a station only once during a day
 - ☞ We should not introduce unintended constraints

Case study 2: second design



Is the extra complexity worth it?