Relational Database Design using E/R

Introduction to Databases CompSci 316 Spring 2019



Announcements (Thu. Jan. 17)

- Sign up for Piazza and Gradiance
- Set up your VM
- Homework 1 Problems 1 and 2 posted (after class!)
 - Due in 2.5 weeks
 - Problem 1 is already posted on gradiance (attempt as many times as you need!)
- Office hours have been posted on website
 - First TA office hour tomorrow

Relational model: review

- A database is a collection of relations (or tables)
- Each relation has a set of attributes (or columns)
- Each attribute has a name and a domain (or type)
- Each relation contains a set of tuples (or rows)

Keys

- A set of attributes K is a key for a relation R if
 - In no instance of R will two different tuples agree on all attributes of K
 - That is, K can serve as a "tuple identifier"
 - No proper subset of K satisfies the above condition
 - That is, K is minimal
- Example: User (uid, name, age, pop)
 - uid is a key of User
 - age is not a key (not an identifier)
 - {uid, name} is not a key (not minimal)

Schema vs. instance

uid	name	age	рор
142	Bart	10	0.9
123	Milhouse	10	0.2
857	Lisa	8	0.7
456	Ralph	8	0.3

- Is name a key of User?
 - Yes? Seems reasonable for this instance
 - No! User names are not unique in general
- Key declarations are part of the schema

More examples of keys

- Member (uid, gid)
 - {uid, gid}
 - A key can contain multiple attributes
- Address (street_address, city, state, zip)
 - {street_address, city, state}
 - {street_address, zip}
 - A relation can have multiple keys!
 - We typically pick one as the "primary" key, and <u>underline</u> all its attributes, e.g., Address (<u>street_address</u>, city, state, <u>zip</u>)

Use of keys

- More constraints on data, fewer mistakes
- Look up a row by its key value
 - Many selection conditions are "key = value"
- "Pointers" to other rows (often across tables)
 - Example: Member (uid, gid)
 - uid is a key of User
 - gid is a key of Group
 - A Member row "links" a User row with a Group row
 - Many join conditions are "key = key value stored in another table"

Database design

- Understand the real-world domain being modeled
- Specify it using a database design model
 - More intuitive and convenient for schema design
 - But not necessarily implemented by DBMS
 - A few popular ones:
 - Entity/Relationship (E/R) model
 - Object Definition Language (ODL)
 - UML (Unified Modeling Language)
- Translate specification to the data model of DBMS
 - Relational, XML, object-oriented, etc.
- Create DBMS schema

But what about ORM?

- Automatic object-relational mappers are made popular by rapid Web development frameworks
 - For example, with Python SQLAlchemy:
 - You declare Python classes and their relationships
 - It automatically converts them into database tables
 - If you want, you can just work with Python objects, and never need to be aware of the database schema or write SQL
- But you still need designer discretion in all but simple cases
- Each language/library has its own syntax for creating schema and for querying/modifying data
 - Quirks and limitations cause portability problems
 - They are not necessarily easier to learn than SQL

Entity-relationship (E/R) model

- Historically and still very popular
- Concepts applicable to other design models as well
- Can think of as a "watered-down" object-oriented design model
- Primarily a design model—not directly implemented by DBMS
- Designs represented by E/R diagrams
 - We use the style of E/R diagram covered by the GMUW book; there are other styles/extensions
 - Very similar to UML diagrams

E/R basics

- Entity: a "thing," like an object
- Entity set: a collection of things of the same type, like a relation of tuples or a class of objects
 - Represented as a rectangle
- Relationship: an association among entities
- Relationship set: a set of relationships of the same type (among same entity sets)
 - Represented as a diamond
- Attributes: properties of entities or relationships, like attributes of tuples or objects
 - Represented as ovals

An example E/R diagram

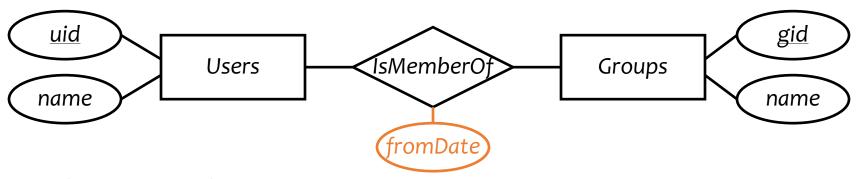
Users are members of groups



- A key of an entity set is represented by underlining all attributes in the key
 - A key is a set of attributes whose values can belong to at most one entity in an entity set—like a key of a relation

Attributes of relationships

Example: a user belongs to a group since a particular date



- Where do the dates go?
 - With Users?
 - But a user can join multiple groups on different dates
 - With Groups?
 - But different users can join the same group on different dates
 - With IsMemberOf!

More on relationships

- There could be multiple relationship sets between the same entity sets
 - Example: Users IsMemberOf Groups; Users Likes Groups
- In a relationship set, each relationship is uniquely identified by the entities it connects
 - Example: Between Bart and "Dead Putting Society", there can be at most one *IsMemberOf* relationship and at most one *Likes* relationship
 - What if Bart joins DPS, leaves, and rejoins? How can we modify the design to capture historical membership information?
 - Make an entity set of MembershipRecords

Multiplicity of relationships

- *E* and *F*: entity sets
- Many-many: Each entity in E is related to 0 or more entities in F and vice versa
 - Example:



- Many-one: Each entity in *E* is related to 0 or 1 entity in *F*, but each entity in *F* is related to 0 or more in *E*
 - Example:



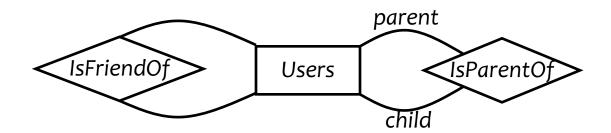
- One-one: Each entity in E is related to 0 or 1 entity in F and vice versa
 - Example:



- "One" (o or 1) is represented by an arrow
- "Exactly one" is represented by a rounded arrow

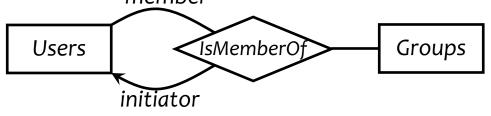
Roles in relationships

- An entity set may participate more than once in a relationship set
- May need to label edges to distinguish roles
- Examples
 - Users may be parents of others; label needed
 - Users may be friends of each other; label not needed



n-ary relationships

• Example: a user must have an initiator in order to join a group



Rule for interpreting an arrow into entity set E in an n-ary relationship:

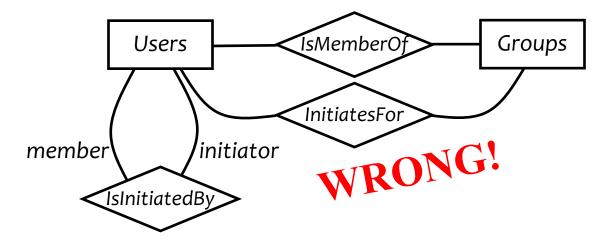
Pick one entity from each of the other entity sets;
 together they can be related to at most one entity in E

• Exercise: hypothetically, member what do these arrows imply?

Users IsMemberOf Groups initiator

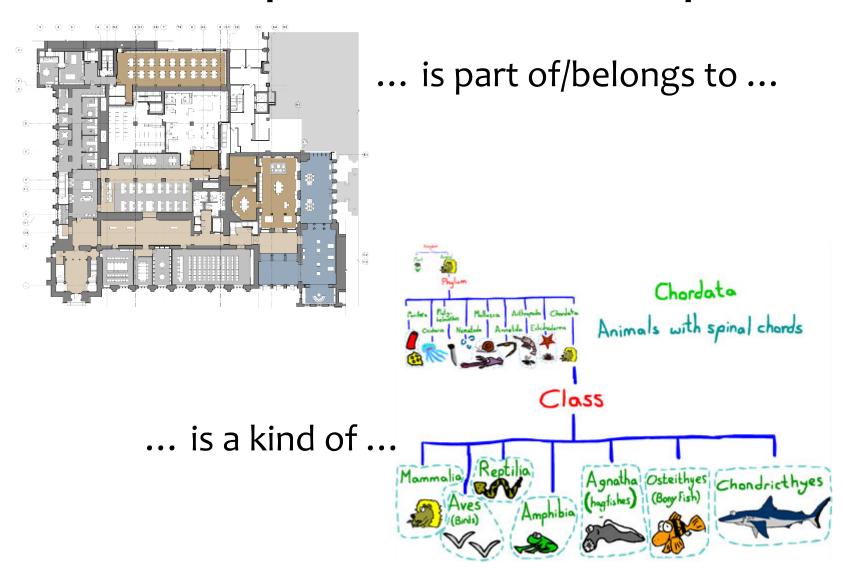
n-ary versus binary relationships

• Can we model *n*-ary relationships using just binary relationships?



- No; for example:
 - Ralph is in both abc and gov
 - Lisa has served as initiator in both abc and gov
 - Ralph was initiated by Lisa in abc, but not by her in gov

Next: two special relationships



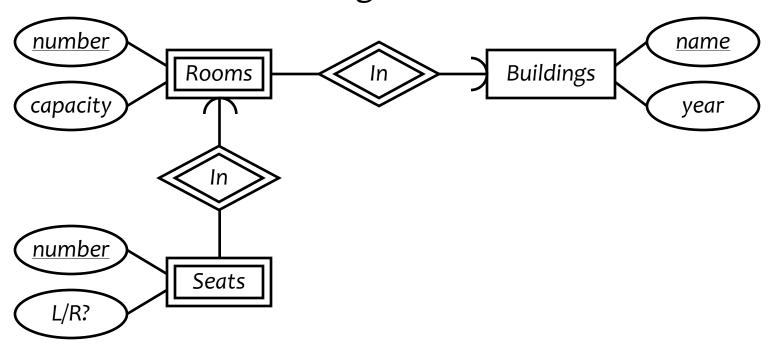
Weak entity sets

Sometimes, an entity's identity depends on some others'

- The key of a weak entity set *E* comes not completely from its own attributes, but from the keys of one or more other entity sets
 - E must link to them via many-one or one-one relationship sets
- Example: Rooms inside Buildings are partly identified by Buildings' name
- A weak entity set is drawn as a double rectangle
- The relationship sets through which it obtains its key are called supporting relationship sets, drawn as double diamonds

Weak entity set examples

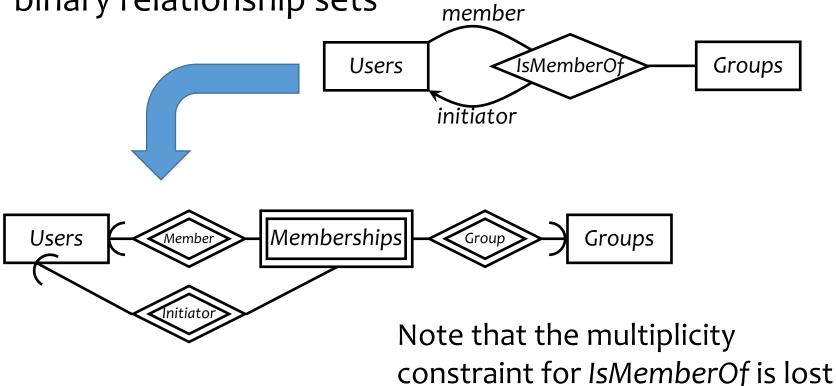
Seats in rooms in building



- Why must double diamonds be many-one/one-one?
 - With many-many, we would not know which entity provides the key value!

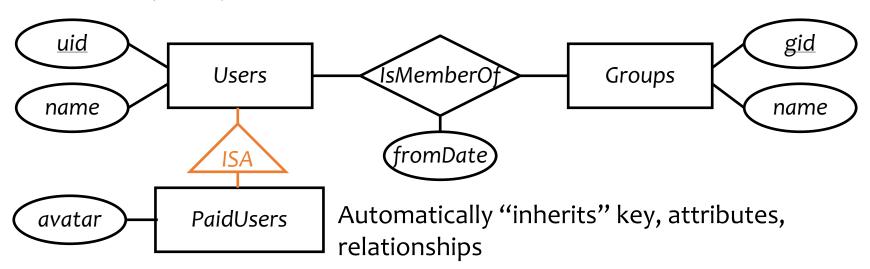
Remodeling n-ary relationships

• An n-ary relationship set can be replaced by a weak entity set (called a connecting entity set) and n binary relationship sets



ISA relationships

- Similar to the idea of subclasses in object-oriented programming: subclass = special case, fewer entities, and possibly more properties
 - Represented as a triangle (direction is important)
- Example: paid users are users, but they also get avatars (yay!)



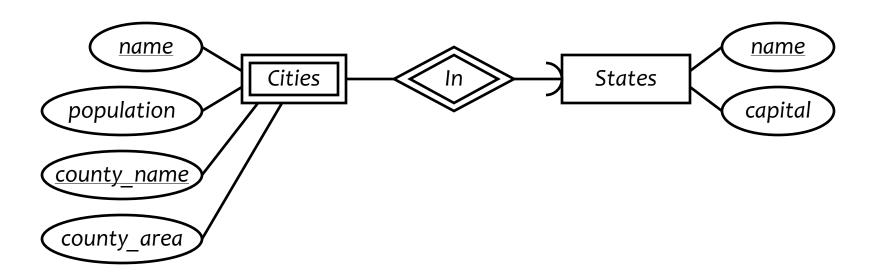
Summary of E/R concepts

- Entity sets
 - Keys
 - Weak entity sets
- Relationship sets
 - Attributes of relationships
 - Multiplicity
 - Roles
 - Binary versus *n*-ary relationships
 - Modeling n-ary relationships with weak entity sets and binary relationships
 - ISA relationships

Case study 1

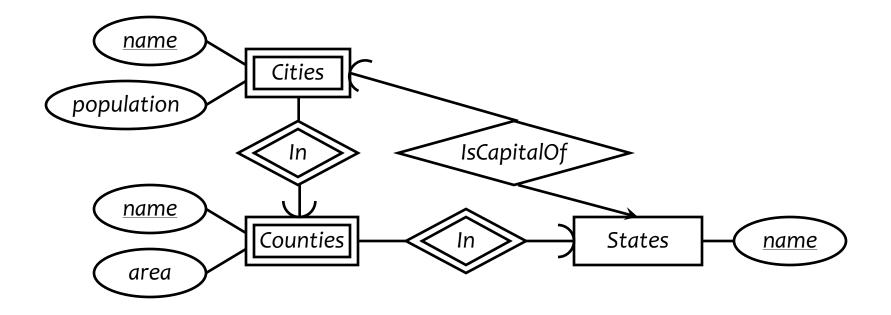
- Design a database representing cities, counties, and states
 - For states, record name and capital (city)
 - For counties, record name, area, and location (state)
 - For cities, record name, population, and location (county and state)
- Assume the following:
 - Names of states are unique
 - Names of counties are only unique within a state
 - Names of cities are only unique within a county
 - A city is always located in a single county
 - A county is always located in a single state

Case study 1: first design



- County area information is repeated for every city in the county
 - Redundancy is bad (why?)
- State capital should really be a city
 - Should "reference" entities through explicit relationships

Case study 1: second design

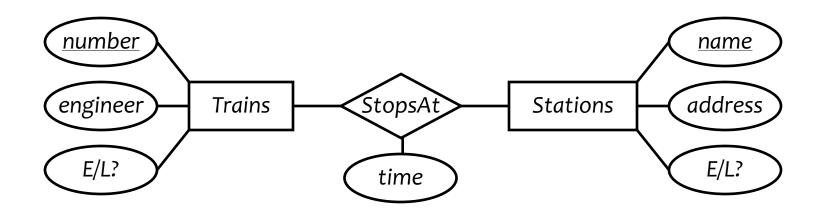


• Technically, nothing in this design prevents a city in state *X* from being the capital of another state *Y*, but oh well...

Case study 2

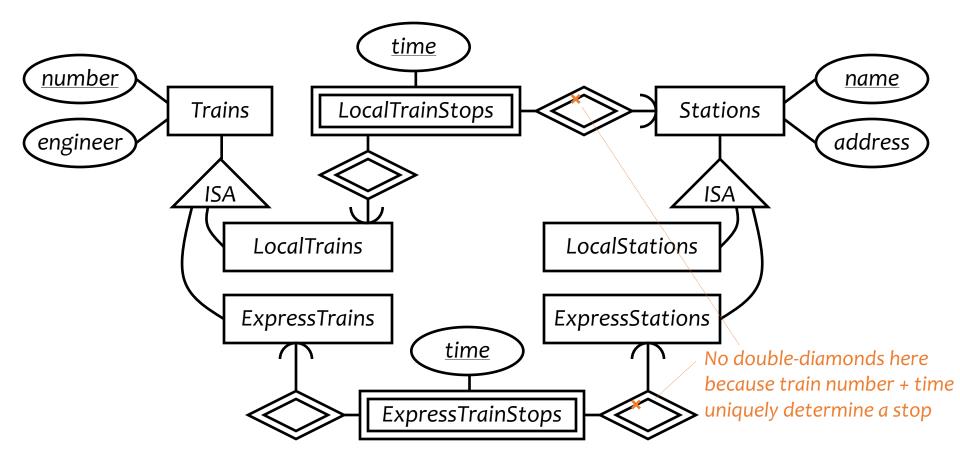
- Design a database consistent with the following:
 - A station has a unique name and an address, and is either an express station or a local station
 - A train has a unique number and an engineer, and is either an express train or a local train
 - A local train can stop at any station
 - An express train only stops at express stations
 - A train can stop at a station for any number of times during a day
 - Train schedules are the same everyday

Case study 2: first design



- Nothing in this design prevents express trains from stopping at local stations
 - *We should capture as many constraints as possible
- A train can stop at a station only once during a day
 We should not introduce unintended constraints

Case study 2: second design



Is the extra complexity worth it?