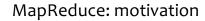
# Map-Reduce and Spark

Introduction to Databases CompSci 316 Spring 2019



#### Announcements (Tue., Apr. 16)

- Project demos—sign-up instructions to be emailed soon
- Homework #4 final due dates
  - Problem 3: today 04/16
  - Problems 4, 5, 6: next Monday 04/22
  - Problem X1: next Wednesday 04/24

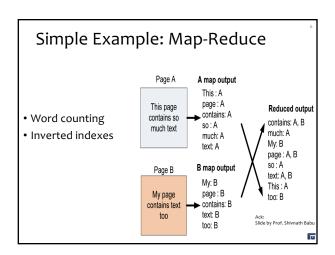


- Many problems can be processed in this pattern:
  - Given a lot of unsorted data
  - Map: extract something of interest from each record
  - Shuffle: group the intermediate results in some way
  - Reduce: further process (e.g., aggregate, summarize, analyze, transform) each group and write final results (Customize map and reduce for problem at hand)
- Make this pattern easy to program and efficient to run
  - Original Google paper in OSDI 2004
  - Hadoop has been the most popular open-source implementation
  - · Spark still supports it

# The Good

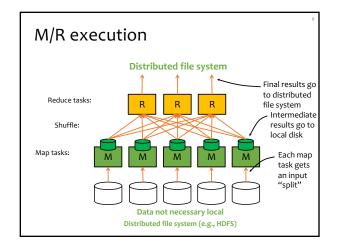
# M/R programming model

- Input/output: each a collection of key/value pairs
- Programmer specifies two functions
  - $map(k_1, v_1) \rightarrow list(k_2, v_2)$ 
    - Processes each input key/value pair, and produces a list of intermediate key/value pairs
  - $\operatorname{reduce}(k_2, \operatorname{list}(v_2)) \to \operatorname{list}(v_3)$ 
    - Processes all intermediate values associated with the same key, and produces a list of result values (usually just one for the key)



#### M/R example: word count

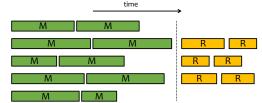
- Expected input: a huge file (or collection of many files) with millions of lines of English text
- Expected output: list of (word, count) pairs
- Implementation
  - map(\_, line) → list(word, count)
    - Given a line, split it into words, and output (w, 1) for each word w in the line
  - reduce(word, list(count)) → (word, count)
  - Given a word w and list L of counts associated with it, compute  $s = \sum_{\text{count} \in L} \text{count}$  and output (w, s)
  - Optimization: before shuffling, map can pre-aggregate word counts locally so there is less data to be shuffled
    - This optimization can be implemented in Hadoop as a "combiner"



#### Some implementation details

- There is one "master" node
- Input file gets divided into m "splits," each a contiguous piece of the file
- Master assigns m map tasks (one per split) to "workers" and tracks their progress
- ullet Map output is partitioned into r "regions"
- Master assigns r reduce tasks (one per region) to workers and tracks their progress
- Reduce workers read regions from the map workers' local disks

# M/R execution timeline



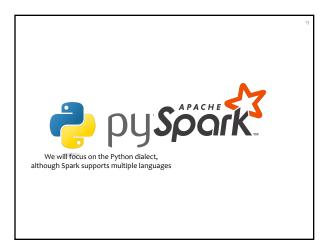
- When there are more tasks than workers, tasks execute in "waves"
  - Boundaries between waves are usually blurred
- Reduce tasks can't start until all map tasks are done

#### More implementation details

- Numbers of map and reduce tasks
  - Larger is better for load balancing
  - But more tasks add overhead and communication
- Worker failure
  - Master pings workers periodically
  - If one is down, reassign its split/region to another worker
- "Straggler": a machine that is exceptionally slow
  - Pre-emptively run the last few remaining tasks redundantly as backup

# M/R example: Hadoop TeraSort

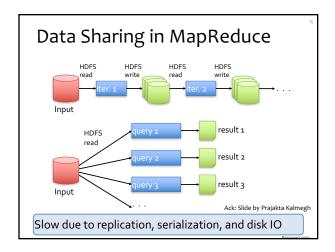
- Expected input: a collection of (key, payload) pairs
- Expected output: sorted (key, payload) pairs
- Implementation
  - Using a small sample of input, find r-1 key values that divides the key range into r subranges where # pairs is roughly equal across them
  - $map(k, payload) \rightarrow (j, \langle k, payload \rangle)$ 
    - If k falls within the j-th subrange
  - reduce $(j, list(\langle k, payload \rangle)) \rightarrow list(k, payload)$ 
    - Sort the list of (k, payload) pairs by k and output

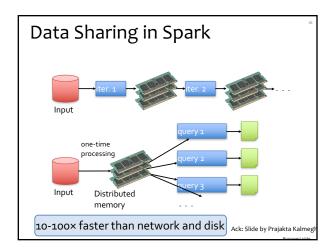


# Why a New Programming Model?

- MapReduce greatly simplified big data analysis
- But as soon as it got popular, users wanted more:
- More complex, multi-stage iterative applications (graph algorithms, machine learning)
- More interactive ad-hoc queries
- More real-time online processing
- All three of these apps require fast data sharing across parallel jobs

Ack: Slide by Prajakta Kalmegh





# Addressing inefficiencies in Hadoop

- Hadoop: no automatic optimization
   Spark
  - Allow program to be a DAG of DB-like operators, with less reliance on black-box code
  - Delay evaluation as much as possible
  - Fuse operators into stages and compile each stage
- Hadoop: too many I/Os
  - E.g., output of each M/R job is always written to disk
  - But such checkpointing simplifies failure recovery

#### ® Spark

- Keep intermediate results in memory
- Instead of checkpointing, use "lineage" for recovery

#### **RDDs**

- Spark stores all intermediate results as Resilient Distributed Datasets (RDDs)
  - Immutable, memory-resident, and distributed across multiple nodes
- Spark also tracks the "lineage" of RDDs, i.e., what expressions computed them
  - Can be done at the partition level

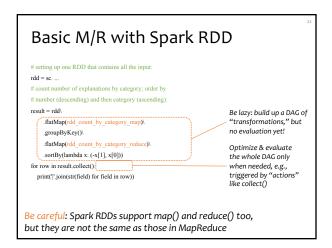
# What happens to a RDD if a node crashes?

- The partition of this RDD on this node will be lost
- But with lineage, the master simply recomputes the a lost partition when needed
  - Requires recursive recomputation if input RDD partitions are also missing

#### Example: votes & explanations

- Raw data reside in lots of JSON files obtained from ProPublica API
- Each vote: URI (id), question, description, date, time, result
- Each explanation: member id, name, state, party, vote URI, date, text, category
  - E.g., "Pooo523", "David E. Price", "NC", "D", "https://api.propublica.org/congress/v1/115/house/sessio ns/2/votes/269.json", "2018-06-20", "Mr. Speaker, due to adverse weather and numerous flight delays and cancellations in North Carolina, I was unable to vote yesterday during Roll Call 269, the motion...", "Travel difficulties"

# Basic M/R with Spark RDD explain\_fields = ('member\_id', 'name', 'state', 'party', 'vote\_api\_uri', 'date', 'text', 'category') # Map function: $map(k_1, v_1) \rightarrow list(k_2, v_2)$ def rdd\_count\_by\_category\_map(record): if len(record) == len(explain\_fields): return [(record[explain\_fields index'(category')], 1)] else: return [] # Reduce function: $reduce(k_2, list(v_2)) \rightarrow list(v_3)$ def rdd\_count\_by\_category\_reduce(record): key, vals = record return [(key, len(vals))]



# Moving "BD" to "DB"

Each element in a RDD is an opaque object—hard to program

- Why don't we make each element a "row" with named columns—easier to refer to in processing
  - RDD becomes a *DataFrame* (name from the R language)
  - Still immutable, memory-resident, and distributed
- Then why don't we have database-like operators instead of just MapReduce?
  - Knowing their semantics allows more optimization
- Spark in fact pushed the idea further
  - Spark Dataset = DataFrame with type-checking
  - And just run SQL over Datasets using SparkSQL!

# Spark DataFrame

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