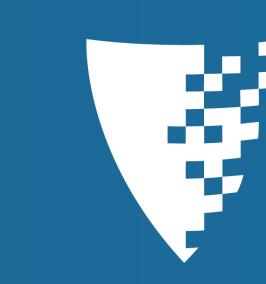


Multi-Level Data Translocation for Faster Processing of Scattered Data on Shared-Memory Computers

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Problem description

- Points $\mathcal{X} = \{\mathbf{x}_i\}_{i=1}^n$ initially stored in order $a(\mathcal{X})$ Access order $b(\mathcal{X})$ far different from $a(\mathcal{X})$
- inter-procedure or inter-operation
- Indirect indexing $a(\mathcal{X}) \mapsto b(\mathcal{X})$
- invokes irregular memory access patterns
- computation becomes acutely memory bounded
- difficult to parallelize
- Common solutions: reordering operations, data, or both
- loop transformations: splitting, fusion, skewing, distribution
- strip-mining, tiling and permutation
- Common solutions are challenged by scattered data
- Fast data translocation: Physical data relocation $\Pi: a(\mathcal{X}) \mapsto b(\mathcal{X})$

Objective

Improve performance of operations on scattered data

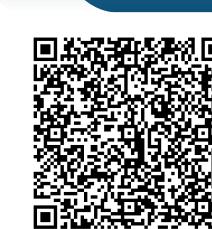
- optimal data locality for minimal memory access latency
- maximal utilization of parallel resources & scheduling schemes

Applications

- Scattered data samples acquired/generated in various applications
- 3D scans, magnetic resonance imaging
- Molecular/celestial dynamics simulations
- Integral imaging, augmented or virtual reality
- Graph embedding
- Processing typically involves calculation of all-point interactions
- direct evaluation too expensive, $O(n^2)$ operations
- approximation/compression techniques: $O(n \log n)$ or O(n)
- * arithmetic operations becomes very small memory operations

computation becomes acutely memory bounded

Contact



Dimitris Floros

Methodology

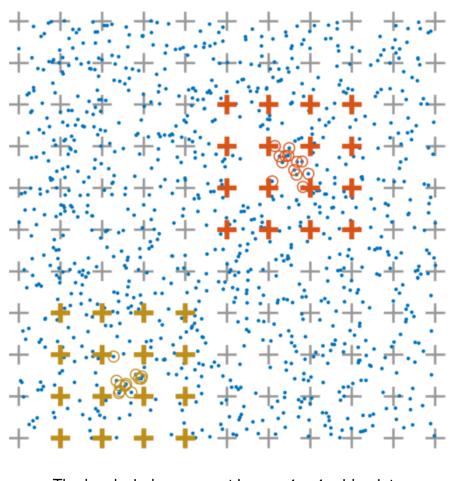
Demo case: scatter-grid translations

Data translation between scattered data points and regular points on an auxiliary grid (externally specified or internally determined)

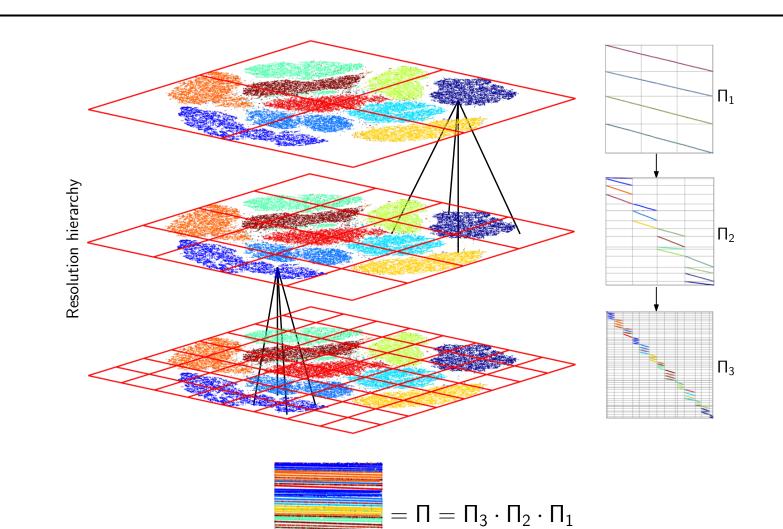
Scattered interactions are decomposed into

- local translations between scattered and grid points (S2G & G2S)
- global interactions among the equispaced grid points (G2G)

This poster focuses on local translations S2G and G2S



The local window support is $w = 4 \times 4$ grid point

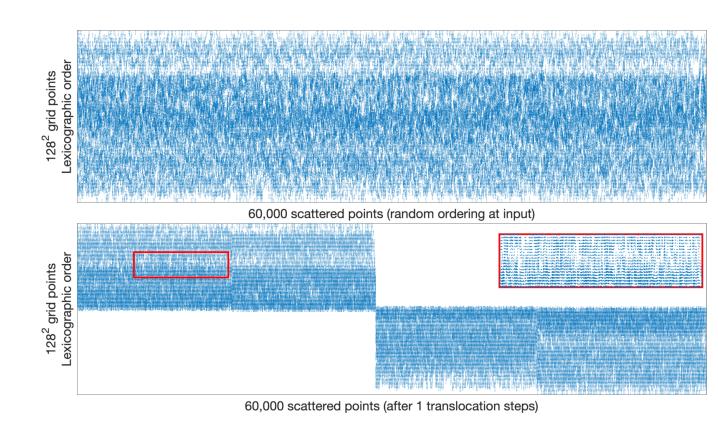


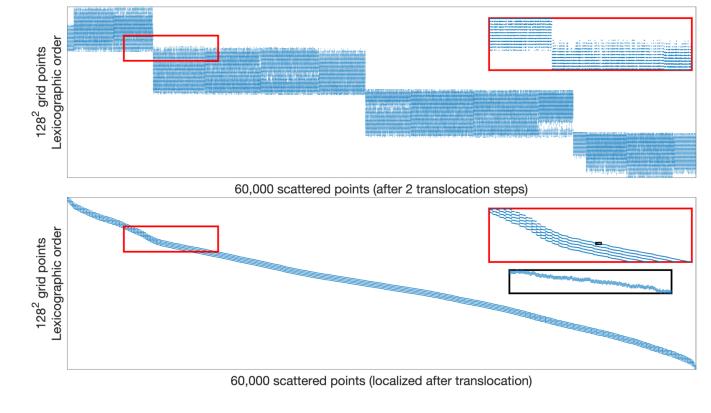
Multi-level data translocation

- Hierarchical binning (coarser to finer grids)
- adhere to memory hierarchy
- utilize memory bandwidth
- explore data & task parallelism
- Matrix view
- block-wise factorization of permutation Π
- blocks not necessarily of equal size
- recursion not necessarily uniform in size and depth

Memory access patterns

- Scattered points are translocated prior to S2G and G2S translations
- points residing in the same grid cell are placed together





Local translation matrix, at each translocation step, with 960,000 coefficients between 128×128 regular grid points (rows) and 60,000 scattered points (columns). The local window support is $w = 4 \times 4$ grid points.

Red-black non-overlapping partition of a 2D grid, with 30×30 grid points (blue) and 60,000 scattered points

Red-black scheduling

Partition grid into non-overlapping regions (red-black)

- Data coherence
- No write conflicts
- No data racing
- Minimal synchronization barriers
- Maximal use of parallel resources

Architecture specification

CPU	Clock (GHz)	Cores	L1 (KiB) per	L2 (KiB)	L3 (MiB) shared	RAM (GiB)	BW (GiB/s)
Intel Core i7-4558U	2.80	2	64	256	4	8	10.5
Intel Core i7-6700	3.40	4	32	256	8	32	19.9
AMD Ryzen 1900X	3.80	8	96	512	16	64	31.6

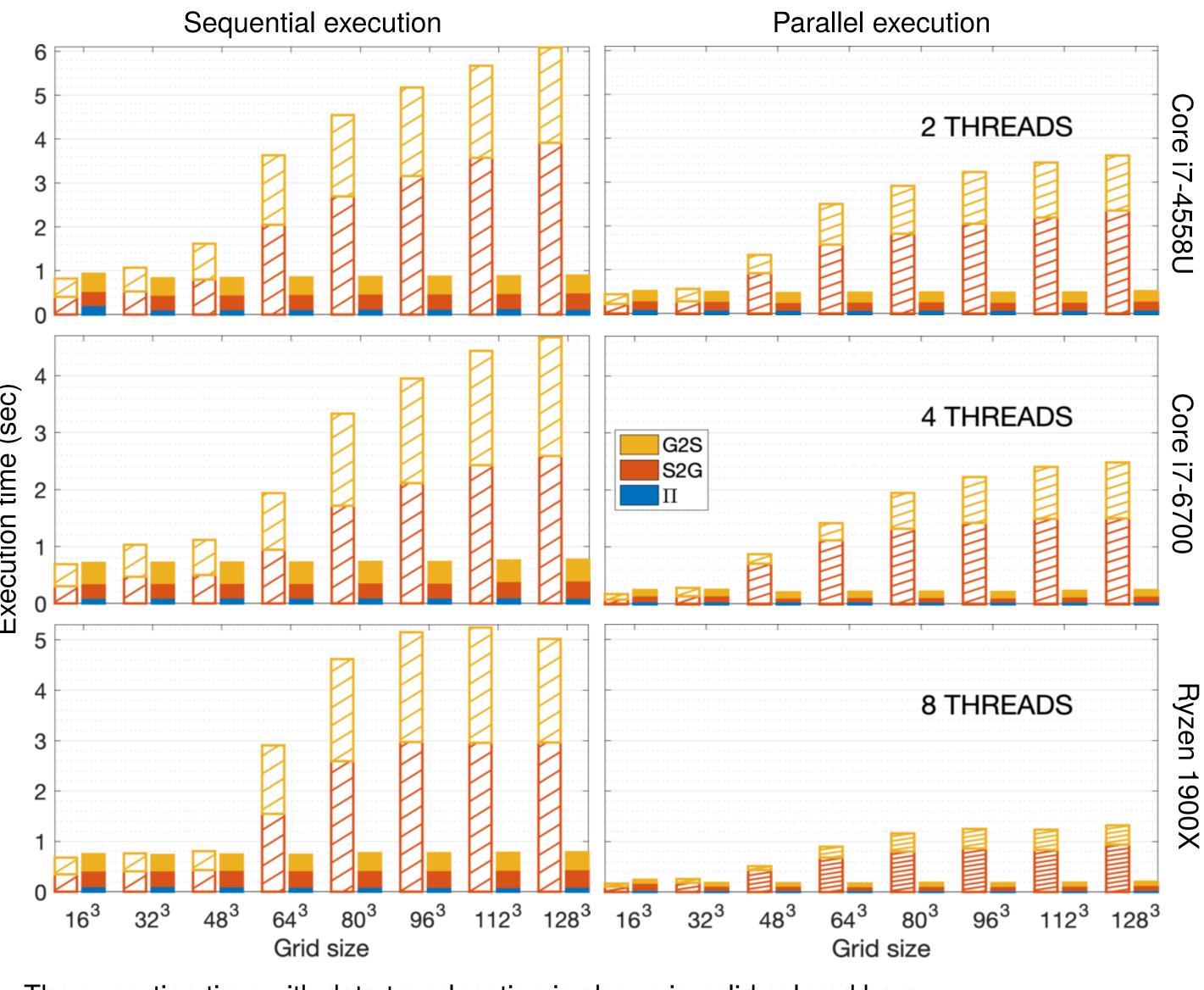
Memory bandwidth measured with the parallel STREAM copy benchmark

Performance results

Performance of S2G & G2S on 3D dataset of n = 2,097,152 scattered points drawn randomly following a uniform distribution over $[0,1)^3$

Speedup at $128 \times 128 \times 128$ grid

Core i	7-4558U	Core i7-6700		Ryzen 1900>	
1	2	1	4	1	8
1.0	1.7	1.0	1.9	1.0	3.8
6.9	11.9	6.1	19.5	6.4	24.6
6.9	7.0	6.1	10.3	6.4	6.5
	1 1.0 6.9	1 2 1.0 1.7 6.9 11.9	1 2 1 1.0 1.7 1.0 6.9 11.9 6.1	1 2 1 4 1.0 1.7 1.0 1.9 6.9 11.9 6.1 19.5	1 2 1 4 1 1.0 1.7 1.0 1.9 1.0 6.9 11.9 6.1 19.5 6.4



The execution time with data translocation is shown in solid colored bars. The data translocation overhead, denoted as Π (blue bar), is well paid-off

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Acknowledgments

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