Section: Properties of Regular Languages

Example

\[ L = \{a^n b a^n \mid n > 0\} \]

Closure Properties

A set is closed over an operation if

\[ L_1, L_2 \in \text{class} \]
\[ L_1 \text{ op } L_2 = L_3 \]
\[ \Rightarrow L_3 \in \text{class} \]
\( L_1 = \{ x \mid x \text{ is a positive even integer} \} \)

\( L \) is closed under

- addition?
- multiplication?
- subtraction?
- division?

Closure of Regular Languages

Theorem 4.1 If \( L_1 \) and \( L_2 \) are regular languages, then

\[
L_1 \cup L_2 \\
L_1 \cap L_2 \\
L_1 L_2 \\
\overline{L}_1 \\
L_1^* 
\]

are regular languages.
Proof(sketch)

$L_1$ and $L_2$ are regular languages

$\Rightarrow \exists$ reg. expr. $r_1$ and $r_2$ s.t.

$L_1 = L(r_1)$ and $L_2 = L(r_2)$

$r_1 + r_2$ is r.e. denoting $L_1 \cup L_2$

$\Rightarrow$ closed under union

$r_1r_2$ is r.e. denoting $L_1L_2$

$\Rightarrow$ closed under concatenation

$r_1^*$ is r.e. denoting $L_1^*$

$\Rightarrow$ closed under star-closure
complementation:
\[ L_1 \text{ is reg. lang.} \]
\[ \Rightarrow \exists \text{ DFA } M \text{ s.t. } L_1 = L(M) \]
Construct \( M' \) s.t.
final states in \( M \) are
nonfinal states in \( M' \)
nonfinal states in \( M \) are
final states in \( M' \)
\[ \Rightarrow \text{ closed under complementation} \]
intersection:

$L_1$ and $L_2$ are reg. lang.

$\Rightarrow \exists$ DFA $M_1$ and $M_2$ s.t.

$L_1 = L(M_1)$ and $L_2 = L(M_2)$

$M_1 = (Q, \Sigma, \delta_1, q_0, F_1)$

$M_2 = (P, \Sigma, \delta_2, p_0, F_2)$

Construct $M' = (Q', \Sigma, \delta', (q_0, p_0), F')$

$Q' = (Q \times P)$

$\delta'$:

$\delta'((q_i, p_j), a) = (q_k, p_l)$ if

$w \in L(M') \iff w \in L_1 \cap L_2$

$\Rightarrow$ closed under intersection
Regular languages are closed under

reversal $L^R$

difference $L_1 - L_2$

right quotient $L_1 / L_2$

homomorphism $h(L)$

Right quotient

Def: $L_1 / L_2 = \{ x | xy \in L_1 \text{ for some } y \in L_2 \}$

Example:

$L_1 = \{ a^*b^* \cup b^*a^* \}$

$L_2 = \{ b^n | n \text{ is even, } n > 0 \}$

$L_1 / L_2 =$
Theorem If $L_1$ and $L_2$ are regular, then $L_1/L_2$ is regular.

Proof (sketch)

$\exists$ DFA $M=\langle Q, \Sigma, \delta, q_0, F \rangle$ s.t. $L_1 = L(M)$.

Construct DFA $M' = \langle Q, \Sigma, \delta, q_0, F' \rangle$

For each state $i$ do

Make $i$ the start state (representing $L_i'$)

if $L_i' \cap L_2 \neq \emptyset$ then

put $q_i$ in $F'$ in $M'$

QED.
Homomorphism

Def. Let $\Sigma, \Gamma$ be alphabets. A homomorphism is a function

$$h: \Sigma \rightarrow \Gamma^*$$

Example:

$$\Sigma = \{a, b, c\}, \quad \Gamma = \{0, 1\}$$

- $h(a) = 11$
- $h(b) = 00$
- $h(c) = 0$

$$h(b)$$ =

$$h(ab^*)$$ =


Questions about regular languages:
L is a regular language.

- Given L, Σ, w ∈ Σ*, is w ∈ L?
- Is L empty?
- Is L infinite?
- Does L₁ = L₂?
Identifying Nonregular Languages

If a language $L$ is finite, is $L$ regular?

If $L$ is infinite, is $L$ regular?

- $L_1 = \{a^n b^m | n > 0, m > 0\} = a^* b^*$
- $L_2 = \{a^n b^n | n > 0\}$

Prove that $L_2 = \{a^n b^n | n > 0\}$ is ?

- Proof:
Pumping Lemma: Let $L$ be an infinite regular language. $\exists$ a constant $m > 0$ such that any $w \in L$ with $|w| \geq m$ can be decomposed into three parts as $w = xyz$ with

\[
  |xy| \leq m \\
  |y| \geq 1 \\
  x y^i z \in L \text{ for all } i \geq 0
\]
To Use the Pumping Lemma to prove L is not regular:

- Proof by Contradiction.
  Assume L is regular.
  \[ L \text{ satisfies the pumping lemma.} \]
  Choose a long string \( w \) in \( L \),
  \[ |w| \geq m. \]
  Show that there is NO division of \( w \) into \( xyz \) (must consider all possible divisions) such that \( |xy| \leq m, |y| \geq 1 \) and \( xy^iz \in L \ \forall \ i \geq 0 \).
  The pumping lemma does not hold. Contradiction!
  \[ \Rightarrow L \text{ is not regular. QED.} \]
Example $L = \{a^n c b^n | n > 0\}$

L is not regular.

- **Proof:**
  
  Assume $L$ is regular.
  
  $\Rightarrow$ the pumping lemma holds.
  
  Choose $w =$
Example $L = \{a^n b^{n+s} c^s | n, s > 0\}$

$L$ is not regular.

- **Proof:**
  
  Assume $L$ is regular.
  
  $\Rightarrow$ the pumping lemma holds.
  
  Choose $w =$
  
  So the partition is:
Example $\Sigma = \{a, b\}$,
$L = \{w \in \Sigma^* \mid n_a(w) > n_b(w)\}$

$L$ is not regular.

- Proof:
  
  Assume $L$ is regular.
  
  $\Rightarrow$ the pumping lemma holds.
  
  Choose $w =$
  
  So the partition is:
Example $L=\{a^3b^n c^{n-3} | n > 3\}$

$L$ is not regular.
To Use Closure Properties to prove \( L \) is not regular:

**Proof Outline:**
Assume \( L \) is regular.
Apply closure properties to \( L \) and other regular languages, constructing \( L' \) that you know is not regular.
closure properties \( \Rightarrow \) \( L' \) is regular.
Contradiction!
\( L \) is not regular. QED.

Example \( L = \{ a^3 b^n c^{n-3} | n > 3 \} \)
\( L \) is not regular.

**Proof:** (proof by contradiction)
Assume \( L \) is regular.
Define a homomorphism \( h : \Sigma \rightarrow \Sigma^* \)
\( h(a) = a \ h(b) = a \ h(c) = b \)
\( h(L) = \)
Example \( L = \{a^n b^m a^m | m \geq 0, n \geq 0\} \)

\( L \) is not regular.

- **Proof: (proof by contradiction)**
  
  Assume \( L \) is regular.
Example: $L_1 = \{a^n b^n a^n | n > 0\}$

$L_1$ is not regular.