Announcements (Tue. Sep. 24)

- Guest lecture this Thu. by Bill Adair, Public Policy
  - Time to think about your course project—first milestone in a little more than 3 weeks
- Note changes from tentative syllabus
  - Triggers/views/indexes first; recursion after Adair lecture
- Homework #1 sample solution posted on Sakai
  - Need by Homework #2, Problem 3
- Homework #2 due next Thursday
  - Do Problem 7 (Gradiance) before Problem 3

“Active” data

- Constraint enforcement: When an operation violates a constraint, abort the operation or try to “fix” data
  - Example: enforcing referential integrity constraints
  - Generalize to arbitrary constraints?
- Data monitoring: When something happens to the data, automatically execute some action
  - Example: When price rises above $20 per share, sell
  - Example: When enrollment is at the limit and more students try to register, email the instructor
Triggers

- A trigger is an event-condition-action (ECA) rule
  - When event occurs, test condition; if condition is satisfied, execute action
- Example:
  - Event: whenever there comes a new student…
  - Condition: with GPA higher than 3.0…
  - Action: then make him/her take CPS316!

Trigger example

```sql
CREATE TRIGGER CPS316AutoRecruit
    AFTER INSERT ON Student
    REFERENCING NEW ROW AS newStudent
    FOR EACH ROW
    WHEN (newStudent.GPA > 3.0)
    INSERT INTO Enroll
    VALUES(newStudent.SID, 'CPS316');
```

Trigger options

- Possible events include:
  - INSERT ON table
  - DELETE ON table
  - UPDATE (OF column) ON table
- Granularity—trigger can be activated:
  - FOR EACH ROW modified
  - FOR EACH STATEMENT that performs modification
- Timing—action can be executed:
  - AFTER or BEFORE the triggering event
  - INSTEAD OF the triggering event on views (more later)
Transition variables

- **OLD ROW**: the modified row before the triggering event
- **NEW ROW**: the modified row after the triggering event
- **OLD TABLE**: a hypothetical read-only table containing all modified rows before the triggering event
- **NEW TABLE**: a hypothetical table containing all modified rows after the triggering event

*Not all of them make sense all the time, e.g.*
- **AFTER INSERT** statement-level triggers
  - Can use only **NEW TABLE**
- **BEFORE DELETE** row-level triggers
  - Can use only
  - etc.

Statement-level trigger example

```sql
CREATE TRIGGER CPS316AutoRecruit
AFTER INSERT ON Student
REFERENCING NEW TABLE AS newStudents
FOR EACH STATEMENT
INSERT INTO Enroll
(SELECT SID, 'CPS316'
FROM newStudents
WHERE GPA > 3.0);
```

BEFORE trigger example

- Never give faculty more than 50% raise in one update

```sql
CREATE TRIGGER NotTooGreedy
BEFORE UPDATE OF salary ON Faculty
REFERENCING OLD ROW AS o, NEW ROW AS n
FOR EACH ROW
WHEN (n.salary > 1.5 * o.salary)
SET n.salary = 1.5 * o.salary;
```

*BEFORE triggers are often used to “condition” data*
*Another option is to raise an error in the trigger body to abort the transaction that caused the trigger to fire*
Statement-vs. row-level triggers

Why are both needed?
- Certain triggers are only possible at statement level
- Simple row-level triggers are easier to implement
  - Statement-level triggers require significant amount of state to be maintained in OLD TABLE and NEW TABLE
  - However, a row-level trigger does get fired for each row, so complex row-level triggers may be inefficient for statements that generate lots of modifications

Another statement-level trigger

- Give faculty a raise if GPA's in one update statement are all increasing
  CREATE TRIGGER AutoRaise
  AFTER UPDATE OF GPA ON Student
  REFERENCING OLD TABLE AS o, NEW TABLE AS n
  FOR EACH STATEMENT
  WHEN ( )
  UPDATE Faculty SET salary = salary + 1000;
- A row-level trigger would be difficult to write in this case

System issues

- Recursive firing of triggers
  - Action of one trigger causes another trigger to fire
  - Can get into an infinite loop
    - Some leave it to programmers/database administrators (e.g., PostgreSQL)
    - Some DBMS restrict trigger actions (e.g., Oracle)
    - Many DBMS set a maximum level of recursion (e.g., 16 in DB2)
- Interaction with constraints (very tricky to get right!)
  - When do we check if a triggering event violates constraints?
    - After a BEFORE trigger (so the trigger can fix a potential violation)
    - Before an AFTER trigger
  - AFTER triggers also see the effects of, say, cascaded deletes caused by referential integrity constraint violations
    (Based on DB2; other DBMS may implement a different policy)
Views

- A view is like a “virtual” table
  - Defined by a query, which describes how to compute the view contents on the fly
  - DBMS stores the view definition query instead of view contents
  - Can be used in queries just like a regular table

Creating and dropping views

- Example: CPS316 roster
  - CREATE VIEW CPS316Roster AS
    SELECT SID, name, age, GPA
    FROM Student
    WHERE SID IN (SELECT SID FROM Enroll
      WHERE CID = 'CPS316');
  - To drop a view
    DROP VIEW view_name;

Using views in queries

- Example: find the average GPA of CPS316 students
  - SELECT AVG(GPA) FROM CPS316Roster;
  - To process the query, replace the reference to the view by its definition
  - SELECT AVG(GPA)
    FROM (SELECT SID, name, age, GPA
      FROM Student
      WHERE SID IN (SELECT SID
        FROM Enroll
        WHERE CID = 'CPS316'));
Why use views?

- To hide data from users
- To hide complexity from users
- Logical data independence
  - If applications deal with views, we can change the underlying schema without affecting applications
  - Recall physical data independence: change the physical organization of data without affecting applications
- To provide a uniform interface for different implementations or sources
- Real database applications use tons of views

Modifying views

- Does not seem to make sense since views are virtual
- But does make sense if that is how users see the database
- Goal: modify the base tables such that the modification would appear to have been accomplished on the view

A simple case

```sql
CREATE VIEW StudentGPA AS
  SELECT SID, GPA FROM Student;
DELETE FROM StudentGPA WHERE SID = 123;
```

translates to:

```sql
DELETE FROM Student WHERE SID = 123;
```
An impossible case

CREATE VIEW HighGPAStudent AS
SELECT SID, GPA FROM Student
WHERE GPA > 3.7;
INSERT INTO HighGPAStudent
VALUES(987, 2.5);
∴ No matter what you do on Student, the inserted row
will not be in HighGPAStudent

A case with too many possibilities

CREATE VIEW AverageGPA(GPA) AS
SELECT AVG(GPA) FROM Student;
∴ Note that you can rename columns in view definition
UPDATE AverageGPA SET GPA = 2.5;
∴ Set everybody’s GPA to 2.5?
∴ Adjust everybody’s GPA by the same amount?
∴ Just lower Lisa’s GPA?

SQL92 updateable views

∴ More or less just single-table selection queries
  ∴ No join
  ∴ No aggregation
  ∴ No subqueries
∴ Arguably somewhat restrictive
∴ Still might get it wrong in some cases
  ∴ See the slide titled “An impossible case”
  ∴ Adding WITH CHECK OPTION to the end of the view
definition will make DBMS reject such modifications
INSTEAD OF triggers for views

- CREATE TRIGGER AdjustAverageGPA
  INSTEAD OF UPDATE ON AverageGPA
  REFERENCING OLD ROW AS o, NEW ROW AS n
  FOR EACH ROW
  UPDATE Student
  SET GPA = GPA + (n.GPA - o.GPA);

- What does this trigger do?

Indexes

- An index is an auxiliary persistent data structure
  - Search tree (e.g., B*-tree), lookup table (e.g., hash table), etc.
  - More on indexes later in this course!
- An index on \( R.A \) can speed up accesses of the form
  - \( R.A = \text{value} \)
  - \( R.A > \text{value} \) (sometimes; depending on the index type)
- An index on \( (R.A_1, ..., R.A_n) \) can speed up
  - \( R.A_1 = \text{value}_1 \land \cdots \land R.A_n = \text{value}_n \)
  - \( (R.A_1, ..., R.A_n) > (\text{value}_1, ..., \text{value}_n) \) (again depends)
- Is an index on \( (R.A, R.B) \) equivalent to one on \( (R.B, R.A) \)?
- How about an index on \( R.A \) plus another index on \( R.B \)?

Examples of using indexes

- SELECT * FROM Student WHERE name = 'Bart'
  - Without an index on Student.name: must scan the entire table if we store Student as a flat file of unordered rows
  - With index: go "directly" to rows with name = 'Bart'
- SELECT * FROM Student, Enroll
  WHERE Student.SID = Enroll.SID;
  - Without any index: for each Student row, scan the entire Enroll table for matching SID
  - Sorting could help
  - With an index on Enroll.SID: for each Student row, directly look up Enroll rows with matching SID
Creating and dropping indexes in SQL

- `CREATE [UNIQUE] INDEX index_name ON table_name(column_name1, ..., column_name_n);`
  - With UNIQUE, the DBMS will also enforce that `{column_name1, …, column_name_n}` is a key of `table_name`
- `DROP INDEX index_name;`

- Typically, the DBMS will automatically create indexes for PRIMARY KEY and UNIQUE constraint declarations

Choosing indexes to create

More indexes = better performance?

- Optimal index selection depends on both query and update workload and the size of tables
  - Automatic index selection is still an area of active research

Summary of SQL features covered

- Query
- Modification
- Constraints
- Triggers
- Views
- Indexes