You are encouraged to type in your answers. Homework must be done individually. Due on April 19th, 2010.

1 Understanding TCP retransmission timers

TCPs computes an average round-trip time (RTT) for a connection using an exponentional weighted moving average (EWMA) estimator:

\[ srtt_n = \alpha \cdot RTT_n + (1 - \alpha) \cdot srtt_{n-1} \]  
\[ rttvar_n = \beta \cdot (|RTT_n - srtt_{n-1}|) + (1 - \beta) \cdot rttvar_{n-1} \]  
\[ RTO_n = srtt_n + \gamma \cdot rttvar_n \]

where \( RTT_n \) is the \( n \)th RTT measurement, \( srtt_n \) is the \( n \)th estimate of RTT, \( rttvar_n \) is the \( n \)th estimate of RTT variance, and \( RTO_n \) is the \( n \)th retransmission timeout value. In a typical TCP implementation, the parameters are set to: \( \alpha = 1/4 \), \( \beta = 1/8 \), and \( \gamma = 4 \).

Ben Bitdiddle considers this design a bit convoluted, and proposes to use a simple arithmetic average to estimate the RTT over a fixed number of past samples. That is,

\[ srtt_n = \sum_{i=n-k}^{n} RTT(i) / k + 1; k \geq 0 \]  
\[ RTO_n = \mu \cdot srtt_n \]

where \( k = 99 \), and \( \mu = 2 \). Suppose that at time before \( n = 0 \), a TCP’s RTT time is a constant 10ms. Both RTT estimates in EQ (1) and (4) converge to 10ms. At \( n = 0 \), the RTT of the TCP connection experiences a sudden change, and goes up to 1 second, and does not change afterwards. Answer the following questions:

1. With TCP’s RTO estimate, what’s the RTO at time when \( n = 1 \)?
2. With Ben’s RTO estimate, what’s the RTO at time when \( n = 1 \)?
3. Will TCP suffer spurious (premature) retransmissions? If no, explain why. If yes, compute when the RTO estimate will be sufficiently accurate to prevent spurious retransmission. (Hint: you may want to write a small script to do this.)
4. Will Ben’s TCP suffer spurious (premature) retransmission? If no, explain why. If yes, compute when the RTO estimate will be sufficiently accurate to prevent spurious retransmission.

2 Go deep with Wireshark

Wireshark [http://www.wireshark.org/](http://www.wireshark.org/) is a useful tool for network traffic analysis. In this problem, you will use Wireshark to look at packet traces from a TCP connection, and answer a few questions. Log onto a linux machine

```
ssh -X linux.cs.duke.edu
```

and download the tcp trace from

```
wget http://www.cs.duke.edu/courses/spring10/cps114/labs/tcp-trace
```

Then start Wireshark using the command “wireshark tcp-trace”. You may also install wireshark on your own machine, and finish the assignment there. (If you are using Ubuntu, you can use “sudo apt-get install wireshark” to install the package.) Answer the following questions:

1. Click on the first packet. The window below will show you the packet details. What protocol headers are included in this packet?
2. What TCP flags are set in the first packet?
3. What TCP options are included in the first packet?
4. Which host does the active open open?
5. Go to packet 91. Why is a Dup ACK sent?

6. Why does the sender send packet 96?

7. Go to packet 119. What might cause the sender to send this packet?

8. Go to packet 425. Which host does an active close?

9. Go to the “Statistics” menu, and follow the sub-menu “TCP Stream Graph.” Click the “Time-Sequence Graph (tcptrace).” Zoom this graph so that you can see 1-second time ticks. From this graph, can you tell how many times the TCP sender does slow start, and when each slow start phase finishes?

10. What (duplicate acks or timeout) is likely to trigger the retransmission around time 2.5s?

11. What is likely to limit the sender’s sending window size at time 2 (cwnd, or receiver advertise window)?

12. How many packets are likely to be lost between time 2s and 7s? Circle the correct answer.
   A. One  B. Two  C. Three  D. > Three

13. Around time 9.5s, another retransmission occurs. What is likely to trigger the retransmission?

14. When does the sender receive the ACK for the packet retransmitted at around time 9.5?

15. What limits the sender’s sending window size at time 12s?

16. Which congestion mode is the sender in at time 13s?

17. Can you estimate the sender’s congestion window size at time 14s?

18. Click the “Round Trip Time Graph.” Can you explain why the measured RTT values fluctuate so much over time from what you learned from the Time-Sequence graph?

19. Similarly, Click the “Throughput Graph.” Can you explain why the TCP throughput fluctuates a lot between 5-10 second, but stabilizes afterwards?

3 Understanding Delay-Bandwidth Product and Router Buffer Size

Figure 1 shows a simple network topology. Node A is sending a large file to Node D using TCP. Suppose each data packet’s size is 1K bytes and the ACK size is 40 bytes. The initial slow start threshold $ssthresh$ is set to a very large value (i.e., infinite). To help you answer the question, we have provided you a simulation script tcp-buffer.tcl. You do not need to run the simulation to answer the following questions, but the simulation will help you understand TCP dynamics and provide hints to the answers. You may follow the instructions in the file http://www.cs.duke.edu/courses/spring10/cps114/labs/ns2-instr.txt to add the simulator to your path and run the script.

1. What is the delay-bandwidth product of this TCP connection?

2. Suppose the router B keeps a buffer size of $Buf$. Let $Buf = 100$. What is the maximum TCP congestion window size before the TCP sender suffers from a packet loss?

3. What is the TCP sender’s congestion window size when it detects a packet loss after three duplicate acknowledgment?

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4. What is the connection’s RTT when router B’s buffer is full?

5. What’s the TCP connection’s maximum throughput?

6. Can the router buffer size be reduced so that the maximum TCP throughput remains the same? If so, what’s the minimal router buffer size to “keep the pipe full”?

7. Is it advantageous to keep a large or small buffer? What will be the “optimal” router buffer size in this case?

4 BGP in the wild

The Route Views project (http://www.routeviews.org/) has an archive that stores periodic BGP table dumps from several routers. Each Route Views router peers with a bunch of other routers located all over the world and receives their BGP updates in real time. So we can discover a lot about Internet routing by examining the data Route Views routers collect, or login to the routers. In this problem, we will telnet to a Route Views router and learn a bunch about its internal state.

1. Login to one of the linux cluster machines: ssh linux.cs.duke.edu. In the command line, type telnet route-views.routeviews.org 23. The first parameter after telnet is the router’s DNS name, and the second is the port we telnet into. Enter the username rviews as prompted.

2. Type ? to see the commands supported by the router. You can also type ? to complete a partial command.

3. Now run the command show ip bgp 152.3.0.0/16. The IP prefix 152.3.0.0/16 is Duke’s network prefix. This command tells you the route advertisements the router receives from its peers. How many routes are available from this router to Duke? (The answer will be displayed at the 2nd line of the command’s output.)

4. Now look at the first route output by the command. What’s the IP address of the next hop router that advertises the route? What’s the AS path from the next hop router to Duke? What’s the local preference of the route?

5. Open another terminal, and login to linux.cs.duke.edu. Run the command traceroute -A nextHopRouter, where nextHopRouter is the IP address of the router obtained from the previous step. The -A option outputs the AS path information. What is the AS path from Duke to the next hop router? Is the path the reverse AS path from the next hop router to Duke? If not, explain why this may happen.

6. Now go back to the router login window, and scroll down until you find the line ending with the word “best”. This is the best route. What’s the next hop router that advertises this best route? What’s the best AS path? Can you explain why this route is chosen as the best path?

5 Digging DNS for fun

Learn how to use dig by typing “man dig” on a Unix machine. Use dig to figure out the Akamai DNS server hierarchy. You may start with the domain name us.i2.yimg.com that is a customer of Akamai, and use dig to discover its Akamai canonical name, and use the +norecurse option to discover Akamai name server hierarchy. Turn in each dig commands you used and their outputs.