Section: Finite Automata

Deterministic Finite Accepter (or Automata)

A DFA = (Q, Σ, δ, q₀, F)

where

Q is finite set of states
Σ is tape (input) alphabet
q₀ is initial state
F ⊆ Q is set of final states.
δ: Q × Σ → Q
Example: DFA that accepts even binary numbers.

Transition Diagram:

\[ M = (Q, \Sigma, \delta, q_0, F) = \]

Tabular Format

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<td>q1</td>
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<td>q1</td>
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Example of a move: \( \delta(q_0, 1) = \)
Algorithm for DFA:

Start in start state with input on tape
q = current state
s = current symbol on tape
while (s != blank) do
    q = δ(q,s)
    s = next symbol to the right on tape
if q ∈ F then accept

Example of a trace: 11010
Pictorial Example of a trace for 100:

1) 1 0 0
   q0
   q1

2) 1 0 0
   q0
   q1

3) 1 0 0
   q0
   q1

4) 1 0 0
   q0
   q1
Definition:
\[ \delta^*(q, \lambda) = q \]
\[ \delta^*(q, wa) = \delta(\delta^*(q, w), a) \]

Definition: The language accepted by a DFA \( M = (Q, \Sigma, \delta, q_0, F) \) is set of all strings on \( \Sigma \) accepted by \( M \). Formally, \[
L(M) = \{ w \in \Sigma^* \mid \delta^*(q_0, w) \in F \}\]
Trap State

Example: $L(M) =$
Example:

$L = \{ w \in \Sigma^* \mid w \text{ has an even number of } a\text{’s and an even number of } b\text{’s} \}$
Example: DFA that accepts even binary numbers that have an even number of 1’s.
Definition A language is regular iff there exists DFA $M$ s.t. $L = L(M)$. \


Chapter 2.2

Nondeterministic Finite Automata (or Accepter)

Definition

An NFA = $(Q, \Sigma, \delta, q_0, F)$

where

- $Q$ is finite set of states
- $\Sigma$ is tape (input) alphabet
- $q_0$ is initial state
- $F \subseteq Q$ is set of final states.
- $\delta: Q \times (\Sigma \cup \{\lambda\}) \rightarrow 2^Q$
Example

Note: In this example $\delta(q_0, a) =$

$L =$
Example

\[ L = \{(ab)^n \mid n > 0\} \cup \{a^n b \mid n > 0\} \]
Definition: \( q_j \in \delta^*(q_i, w) \) if and only if there is a walk from \( q_i \) to \( q_j \) labeled \( w \).

Example: From previous example:
\[ \delta^*(q_0, ab) = \]
\[ \delta^*(q_0, aba) = \]

Definition: For an NFA \( M \),
\[ L(M) = \{ w \in \Sigma^* \mid \delta^*(q_0, w) \cap F \neq \emptyset \} \]
2.3 NFA vs. DFA: Which is more powerful?

Example:
Theorem Given an NFA \( M_N = (Q_N, \Sigma, \delta_N, q_0, F_N) \), then there exists a DFA \( M_D = (Q_D, \Sigma, \delta_D, q_0, F_D) \) such that \( L(M_N) = L(M_D) \).

Proof:

We need to define \( M_D \) based on \( M_N \).

\[ Q_D = \]

\[ F_D = \]

\[ \delta_D : \]
Algorithm to construct $M_D$

1. start state is $\{q_0\} \cup \text{closure}(q_0)$
2. While can add an edge
   (a) Choose a state $A=\{q_i, q_j, \ldots q_k\}$ with missing edge for $a \in \Sigma$
   (b) Compute $B = \delta^*(q_i, a) \cup \delta^*(q_j, a) \cup \ldots \cup \delta^*(q_k, a)$
   (c) Add state $B$ if it doesn’t exist
   (d) add edge from $A$ to $B$ with label $a$
3. Identify final states
4. if $\lambda \in L(M_N)$ then make the start state final.
Example:
Properties and Proving - Problem 1

Consider the property 
Replace_one_a_with_b or R1awb for short. If L is a regular, prove 
R1awb(L) is regular.

The property R1awb applied to a 
language L replaces one a in each 
string with a b. If a string does not 
have an a, then the string is not in 
R1awb(L).

Example 1: Consider L={aaab, bbaa} 
R1awb(L)=

Example 2: Consider \( \Sigma = \{a, b\} \), \( L = \{w \in \Sigma^* \mid w \text{ has an even number of a’s and an even number of b’s}\} \) 
R1awb(L)=
Proof:
Properties and Proving - Problem 2

Consider the property
Truncate_all_preceeding_b’s or TruncPreb for short. If L is a regular, prove TruncPreb(L) is regular.

The property TruncPreb applied to a language L removes all preceeding b’s in each string. If a string does not have an preceeding b, then the string is the same in TruncPreb(L).

Example 1: Consider L={aaab, bbba}
TruncPreb(L)=

Example 2: Consider L =
{(bba)^n | n > 0}
TruncPreb(L)=

Proof:
Minimizing Number of states in DFA
Why?
Algorithm

• Identify states that are indistinguishable
  These states form a new state

Definition Two states \( p \) and \( q \) are indistinguishable if for all \( w \in \Sigma^* \)

\[
\delta^*(q, w) \in F \Rightarrow \delta^*(p, w) \in F \\
\delta^*(p, w) \not\in F \Rightarrow \delta^*(q, w) \not\in F
\]

Definition Two states \( p \) and \( q \) are distinguishable if \( \exists \ w \in \Sigma^* \) s.t.

\[
\delta^*(q, w) \in F \Rightarrow \delta^*(p, w) \not\in F \text{ OR } \\
\delta^*(q, w) \not\in F \Rightarrow \delta^*(p, w) \in F
\]
Example:
Example: