

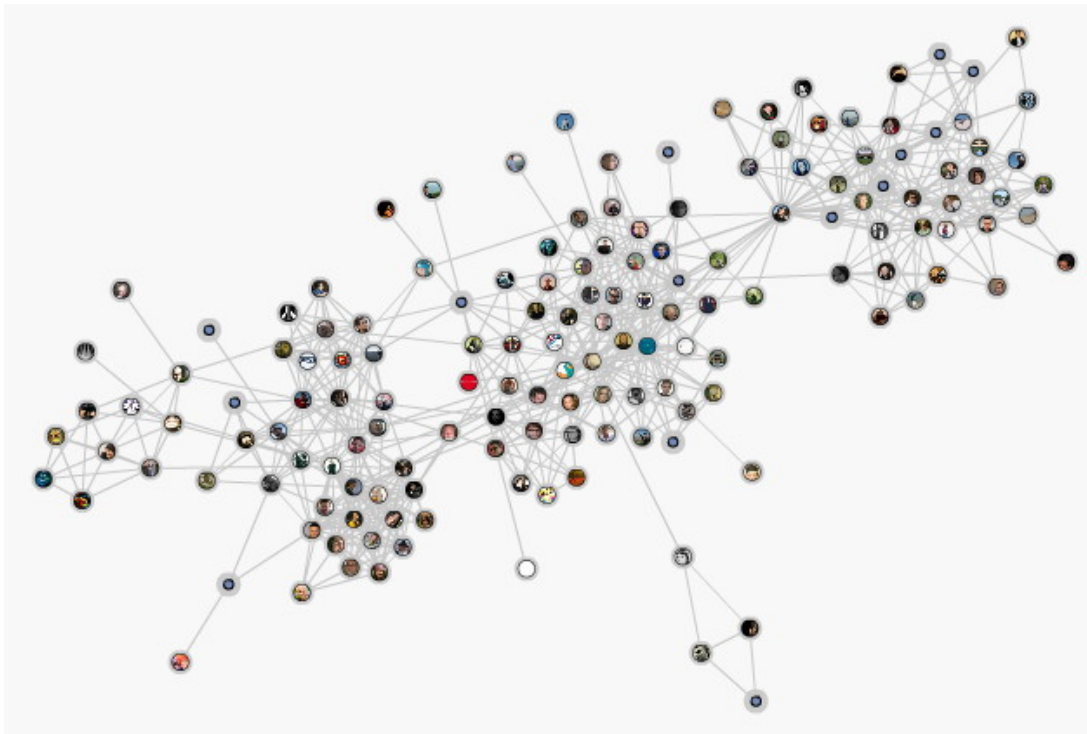
# Matrices and Graphs – Spectral Techniques

**PPT by Brandon Fain**

# Outline

- Motivating Problem: Community Detection
- Spectral Clustering and Conductance
- Graph Laplacian

# Motivating Problem: Community Detection



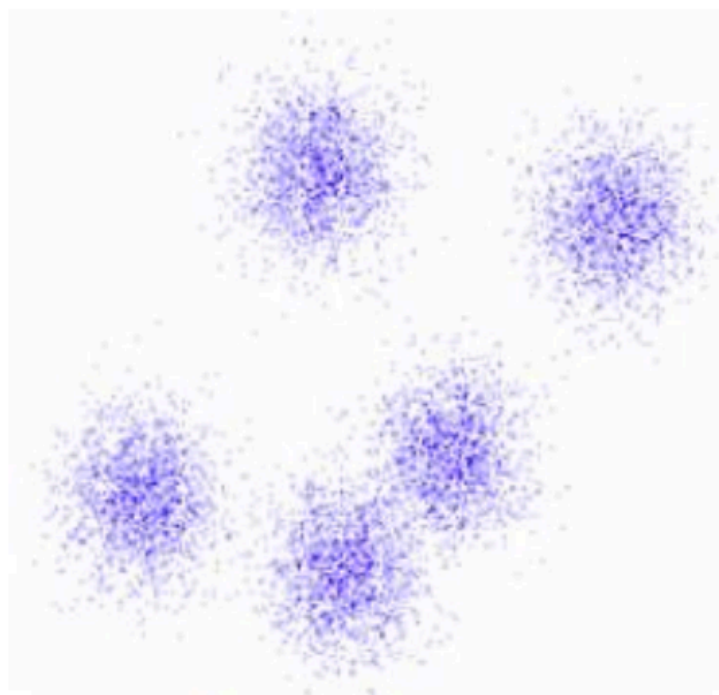
Given a social network, how do you find the strongly connected communities?

Corollary question: How would you suggest friends to a user?

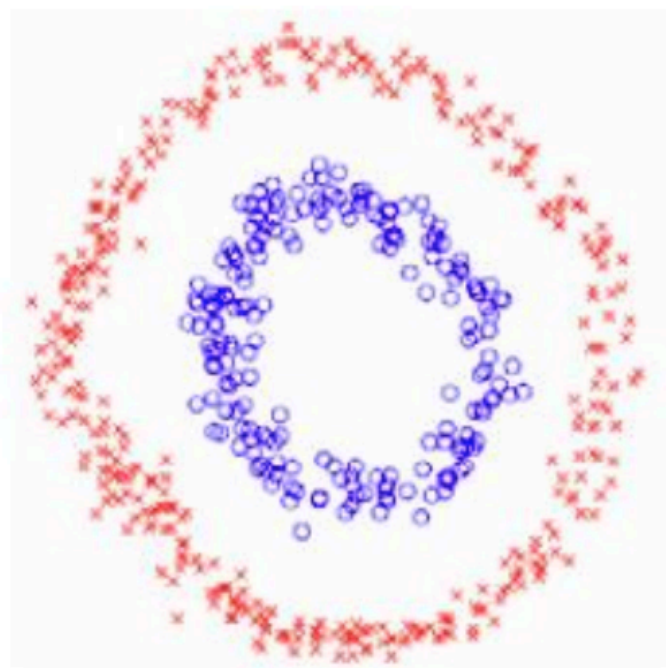
# Motivating Community Detection

- But the question is much broader than this: it's about **clustering**, the fundamental task in unsupervised machine learning.
- Stated loosely, clustering is about creating a *partition* of the data so that points within a partition are more *similar* to one another than points outside of the partition.
- But what is the right notion of “similar?” For geometric data, perhaps it is geometric distance or compactness. What about for a graph?

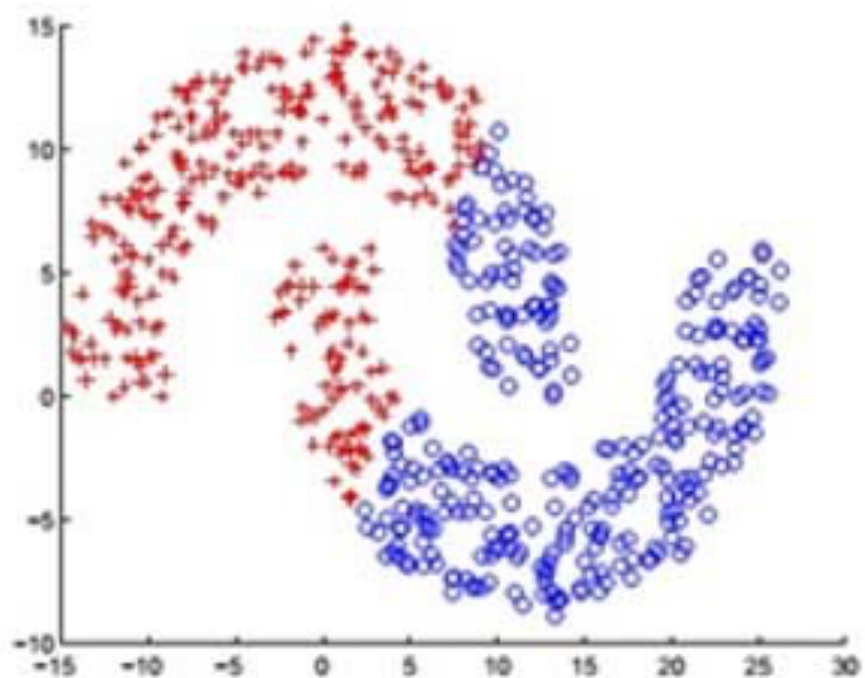
- Compactness, e.g., k-means, mixture models
- Connectivity, e.g., spectral clustering



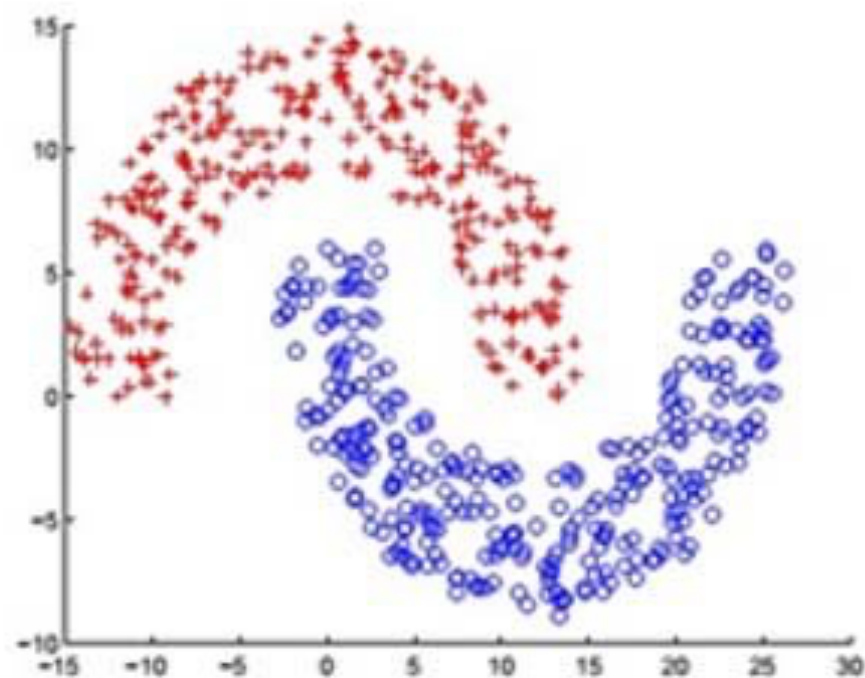
**Compactness**



**Connectivity**



**(a) K-means**



**(b) Spectral Clustering**

# Community Detection – Spectral Clustering

- In fact, as you can probably see, there are other applications where you might like to cluster based on “community structure” in a graph: it captures the idea of similarity by connectivity!
- In such applications, one typically:
  - Given a similarity measure  $S()$ , draw a graph by placing an edge between data points  $x$  and  $y$  with  $S(x,y) < t$ , for some threshold  $t$ .
  - Use **spectral clustering** to partition the graph based on its connectivity (i.e., ignore everything else about the data).

# Community Detection - Spectral Clustering

- Questions:
  - How do we measure the quality of a cluster?
  - How do we compute such a clustering?
- This class and next, we'll try to answer these questions.
- Today, we'll focus on the first question, and reviewing some tools we will need to answer the second next week.



# Outline

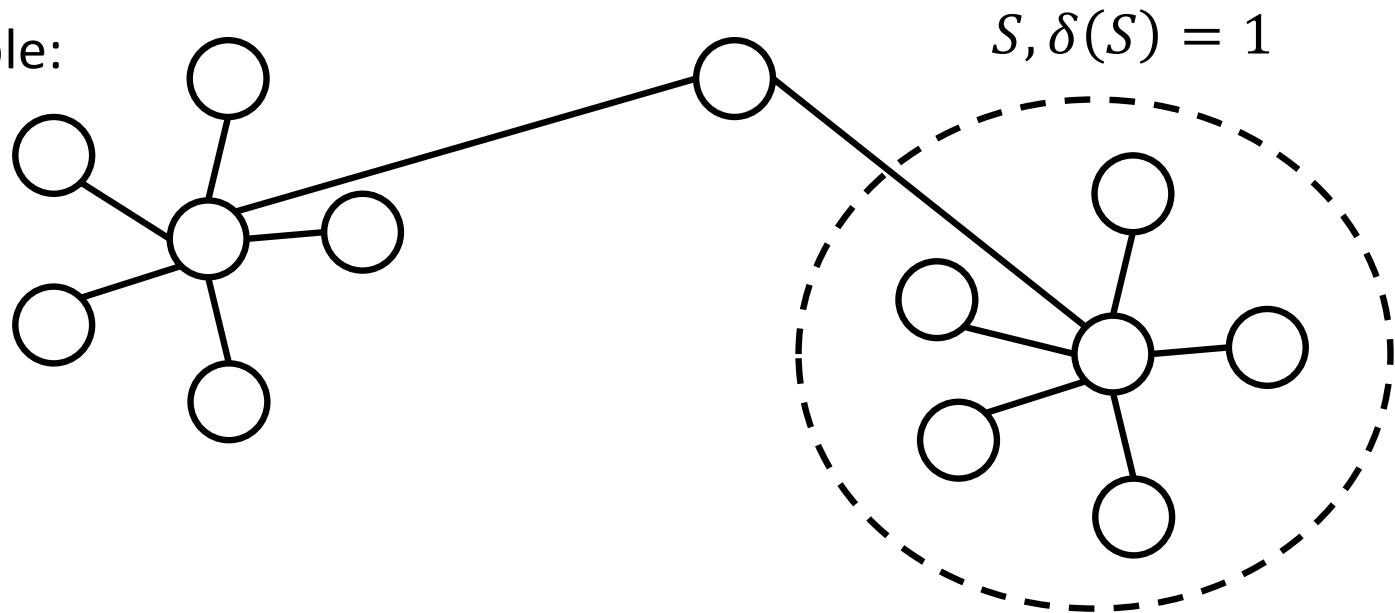
- ~~Motivating Problem: Community Detection~~
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- Graph Laplacian

# Spectral Clustering – Measuring Quality

- In the simplest case, suppose we just want to generate a 2-partition (i.e., split the data into two sets). What objective should we minimize?
- Take 0: Compute the minimum cut in the graph.
  - Let  $G = (V, E)$  be an undirected graph.
  - Let  $S \subseteq V$  denote a cut in the graph.
  - Let  $\delta(S) := |\{(u, v) \in E : u \in S, v \notin S\}|$ .
  - In the minimum cut problem, we want to find a cut  $S$  that minimizes  $\delta(S)$ .

## Take 0 – Minimum Cut

- For example:



- Looks pretty reasonable!
- What is wrong with this notion, and how would you fix it?

# Take 1 – Normalized Minimum Cut

- Somehow, we want to find a *large* partition of vertices with a small number of cut edges.
- Take 1: Compute the minimum *normalized* cut in the graph.
  - Again Let  $G = (V, E)$  be an undirected graph and  $S \subseteq V$  denote a cut in the graph, with  $\delta(S) := |\{(u, v) \in E : u \in S, v \notin S\}|$ .
  - In the minimum *normalized* cut problem, we want to find a cut  $S$  that minimizes

$$\frac{\delta(S)}{|S| \cdot |V - S|}$$

- (Sometimes called the *isoperimetric ratio*)

## Take 2 - Conductance

- The isoperimetric ratio should look very familiar to you.
- Recall that last week we discussed **conductance** as a property of graphs related to how quickly power iteration computes the stationary distribution of a random walk on the graph.
- Let  $S$  be a cut. Let  $Vol(S) = \sum_{i \in S} d_i$ , where  $d_i$  is the degree of node  $i$ . The **conductance** of  $S$  is

$$\phi(S) = \frac{\delta(S)}{\min(Vol(S), Vol(V - S))}.$$

# Conductance

- Last week, we were interested in graphs for which power iteration could quickly converge to find the stationary distribution of a random walk.
- For that, we wanted a graph for which **all** cuts had **high** conductance.
- For the spectral partitioning problem, we will look for a particular cut in the graph with very **low** conductance (which means there are a lot of internal edges, but very few cut edges).

# Outline

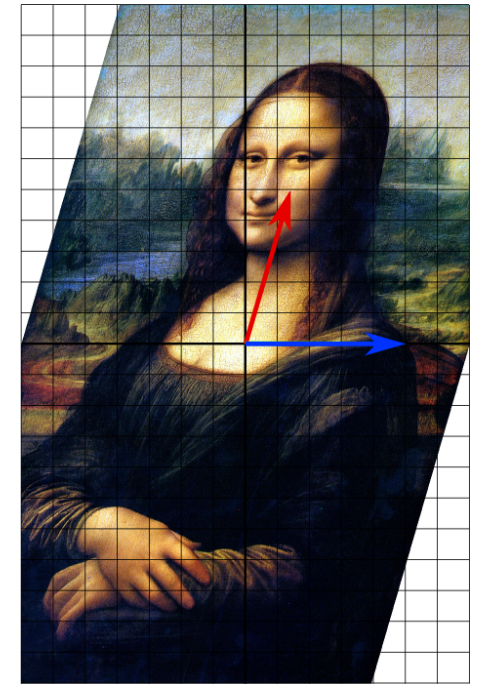
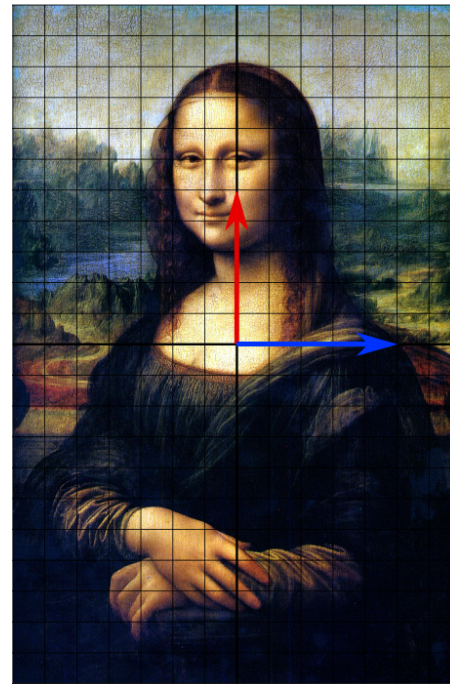
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# Eigenvectors and Eigenvalues

- Consider a matrix  $M$ . We say that  $\lambda$  is an **eigenvalue** of  $M$  with associated **eigenvector**  $\vec{v}$  if

$$M \vec{v} = \lambda \vec{v}.$$

- Think of  $M$  as a linear function or linear transformation. An eigenvector is an input that only changes in scale, *not* direction.





# Laplacian Matrix

- The particular matrix we are interested in is the **graph Laplacian**, defined as

$$L = D - A$$

where  $D$  is the diagonal matrix with  $D_{ii} = d_i$  and  $D_{ij} = 0$  for  $i \neq j$ , and  $A$  is the adjacency matrix.

- There is also a **normalized graph Laplacian** defined as

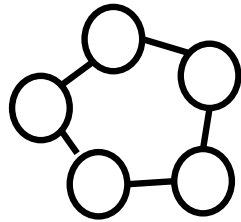
$$\hat{L} = I - D^{-\frac{1}{2}} A D^{-\frac{1}{2}}.$$

It is helpful to get some intuition by trying to write down some graph Laplacians, so let's try a few.

# Laplacian Matrix

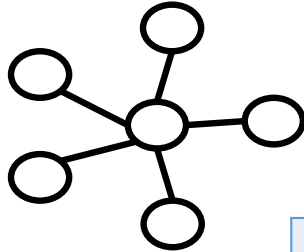
- Write down the Laplacian (un-normalized) of the following:

1. A 5-cycle



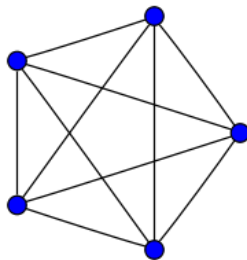
2	-1	0	0	-1
-1	2	-1	0	0
0	-1	2	-1	0
0	0	-1	2	-1
-1	0	0	-1	2

2. A 5-star



5	-1	-1	-1	-1	-1
-1	1	0	0	0	0
-1	0	1	0	0	0
-1	0	0	1	0	0
-1	0	0	0	1	0
-1	0	0	0	0	1

3.  $K_5$



4	-1	-1	-1	-1
-1	4	-1	-1	-1
-1	-1	4	-1	-1
-1	-1	-1	4	-1
-1	-1	-1	-1	4

# Laplacian Matrix

- Why do we care about this matrix? The first intuition is to note that

$$(Lv)_i = d_i v_i - \sum_{j:(i,j) \in E} v_j = \sum_{j:(i,j) \in E} v_i - v_j.$$

Therefore,

$$v^t Lv = \sum_{i \in V} v_i \left( \sum_{j:(i,j) \in E} v_i - v_j \right) = \sum_{(i,j) \in E} v_i (v_i - v_j)$$

$$v^t Lv = \sum_{(i,j) \in E, i < j} (v_i - v_j)^2$$

# Laplacian Matrix

- In other words, given a graph with Laplacian  $L$ , and a vector  $v$  that assigns a value to every vertex in the graph,  $v^t L v$  is the sum of squared differences *of neighbors* in the graph.
- So the Laplacian matrix is definitely recording *something* about connectivity.
- But actually, this simple fact allows us to show a few other connections.

# Laplacian Matrix

- Given  $v^t L v = \sum_{(i,j) \in E, i < j} (v_i - v_j)^2$ .
- Then  $L$  is a positive semidefinite matrix (by definition). It follows that all of the eigenvalues of  $L$  are nonnegative (an excellent exercise).
- Note that it's also clear that  $L \vec{1} = (0) \vec{1}$ , so the smallest eigenvalue is 0, and  $\vec{1}$  is an associated eigenvector.
- In fact, there is an important connection between this eigenvalue and the graph.

# Laplacian Matrix

- **Claim.** The number of connected components in a graph equals the number of orthonormal eigenvectors associated with eigenvalue 0 of the graph Laplacian.
- **Proof.** Suppose there are  $k$  connected components in the graph. For each component  $S \subseteq V$ , define  $\vec{u}^S$  such that  $u_i^S = 1$  if  $i \in S$ , else 0.
- Clearly these vectors are orthogonal, and since  $(Lv)_i = \sum_{j:(i,j) \in E} v_i - v_j$ , we know that  $L \vec{u}^S = (0) \vec{u}^S$ . So there are at least  $k$  such eigenvectors.
- But any other eigenvector also has to take the same value on every vertex of a connected component, so any other such vector is not orthogonal to these. So there are just  $k$  such eigenvectors.

# Laplacian Matrix

- **Corollary.** Given the eigenvectors of the graph Laplacian with eigenvalue 0, one can read off the connected components of the graph.
  - In particular, for each of these 0 eigenvalue eigenvectors, you just have to pick the vertices for which the eigenvector has the same value.
- So far then, we have built a purely algebraic way of computing the connected components of a graph.
- But, of course, you learned how to do this (much more efficiently) with depth first search last week. So who cares?

# Laplacian and Spectral Clustering

- What we really care about for clustering and community detection: finding low conductance cuts in the graph.
- Recall that  $(Lv)_i = \sum_{j:(i,j) \in E} v_i - v_j$ . So if  $v$  is an eigenvector associated with a small (but positive) eigenvalue, it assigns very similar values to neighboring vertices.
- Of course, if the graph is connected, we already know that the eigenvector for the smallest eigenvalue (0) is just  $\vec{1}$ , so that isn't very helpful.



# Laplacians and Spectral Clustering

- What if we take the eigenvalue associated with the *second* smallest eigenvalue (for a connected graph)?
- Then we take a large subset of vertices to which this vector assigns very similar values.
- Such a subset of vertices should be more likely to be neighbors.
- How do we formalize this idea? Can we prove anything about its performance? (Stay tuned for next time!)